

Plan for today

Nominal (re)centering: From Kalaallisut to UC₁

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(W1: July 30, 2009)

- Outline of CCG+UC₁ theory
- Definition of UC₁ (see handout)
- drt-notation (DRT-like, see handout)
- From Kalaallisut text, via UC₁, to empirically testable prediction

Outline of CCG + UC₁

• language-specific **lexicon** + universal **CCG-rules** (both to be given on **M2**)
translate NL-syntax into UC₁-syntax

• A UC₁-model $M = \langle \{D_a \mid a \in \Theta\}, \|\cdot\| \rangle$
assigns to each constant of type a , $A \in \text{Con}_a$ a value of type a , $\|A\| \in D_a$.

• Given model $M = \langle \{D_a \mid a \in \Theta\}, \|\cdot\| \rangle$
and assignment g , UC₁-semantics

i. extends $\|\cdot\|$ to $\|\cdot\|^g$, which assigns
value $\|A\|^g$ to each UC₁-term A

ii. defines *default state* of infotention:

$$c_0 = \{ \langle \langle \rangle, \langle \rangle \rangle \}$$

(no objects available for anaphora)

• Truth definition:

A UC₁-drt K is *true* in M , $M \models K$,

iff $\forall g: c_0 \|K\|^g \neq \emptyset$

(c_0 updated with $\|K\|^g$ is not the
absurd state).

- e.g.

English

(1) *A man saw an enemy.*

to UC₁

(1') $[\mathbf{x} \mid \text{man } \mathbf{x}]; ([y \mid \text{enm}^{of} \langle y, \mathbf{T} \rangle];$

$[\text{see} \langle \mathbf{T}, \perp \rangle])$

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absurd state).

- e.g. here is a possible model
(call it M_0)

$$D_e = \{M, M'\}$$

$$\|man\| = \lambda \{M, M'\}$$

$$\|enm^{of}\| = \lambda \{ \langle M, M' \rangle, \langle M', M \rangle \}$$

$$\|sell\| = \lambda \{ \langle M, M' \rangle \}$$

(Notation: $\lambda \{ \dots \}$ stands for the
characteristic function of $\{ \dots \}$)

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(no objects available for anaphora)

• Truth definition:

A UC₁-drs K is *true* in M , $M \models K$,

iff $\forall g: c_0 \|K\|^g \neq \emptyset$

(c_0 updated with $\|K\|^g$ is not the
absurd state).

• e.g. in M_0 , for any g , we derive:

• **T-list update:** add a man

$$c_0 \| [\mathbf{x} \mid \text{man } \mathbf{x}] \|^g \\ = \{ \langle \langle M \rangle, \langle \rangle \rangle, \\ \langle \langle M' \rangle, \langle \rangle \rangle \}$$

$=: c_1$

• **⊥-list update:** add an enemy of the T-man

$$c_1 \| [y \mid \text{enm}^{of}(y, \mathbf{T})] \|^g \\ = \{ \langle \langle M \rangle, \langle M' \rangle \rangle, \\ \langle \langle M' \rangle, \langle M \rangle \rangle \}$$

$=: c_2$

• **test:** the T-man saw the ⊥-enemy

$$c_2 \| [\text{see}(\mathbf{T}, \perp)] \|^g \\ = \langle \langle \langle M \rangle, \langle M' \rangle \rangle \rangle$$

• $c_0 \| [\mathbf{x} \mid \text{man } \mathbf{x}]; \dots ; [\text{see}(\mathbf{T}, \perp)] \|^g$

$$= \langle \langle \langle M \rangle, \langle M' \rangle \rangle \rangle$$

Once upon a time in the N.

Kal (1) → UC₁ syn, ...

Model M_1 with $D_e = \{M, M', K, K'\}$

$\| \text{man} \| = \{M, M'\}$

$\| \text{enm}^{of} \| = \{ \langle M, M' \rangle, \langle M', M \rangle \}$

$\| \text{kayak} \| = \{K, K'\}$

(*pos* for 'inalienable-possessor-of')

$\| \text{pos} \| (K) = M, \| \text{pos} \| (K') = M'$

$\| \text{use} \| = \{ \langle M, K \rangle, \langle M', K' \rangle \}$

$\| \text{lie.in.wait} \| = \{M, M'\}$

$\| \text{approach} \| = \{ \langle M, M' \rangle, \langle M, K' \rangle \}$

$\| \text{see} \| = \{ \langle M, M' \rangle, \langle M', M \rangle, \\ \langle M, K' \rangle, \langle M', K \rangle \}$

$\| \text{int.to.harpoon} \| = \{ \langle M, M' \rangle, \langle M', M \rangle \}$

$\| \text{forest} \| = \{ \langle M', M \rangle \}$

$\| \text{harpoon} \| = \{ \langle M', M \rangle \}$

$\| \text{capsize} \| = \{M\}$

$\| \text{die} \| = \{M\}$

$\| \text{kill} \| = \{ \langle M', M \rangle \}$

$\| \text{go.hml} \| = \{M'\}$

$\| \text{tell.story} \| = \{M'\}$

(1) Long ago there was a man who had (an)
enemy(ies).

K. angut-qar-pu-q akiraq-lik-mik.

man-have-DEC_{iv}-3S enemy-with-MOD

UC₁ drs

$[y \mid \text{man } y]; -\downarrow [y \mid \text{enm}^{of}(y, \perp)]$

UC₁ sem:

Step 1:

$$c_0 \| [y \mid \text{man } y]; \downarrow [y \mid \text{enm}^{of}(y, \perp)] \|^g \\ = c_0 \| [y \mid \text{man } y] \|^g \| \downarrow [y \mid \text{enm}^{of}(y, \perp)] \|^g \quad \text{D6.x.};$$

Step 2:

Compute each update in turn.

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assigns to each constant of type a , $A \in \text{Con}_a$ a value of type a , $\|A\| \in D_a$.

• Given model $M = \langle \{D_a \mid a \in \Theta\}, \|\cdot\| \rangle$
and assignment g , UC₁-semantics

i. extends $\|\cdot\|$ to $\|\cdot\|^g$, which assigns
value $\|A\|^g$ to each UC₁-term A

ii. defines *default state* of infotention:

$$c_0 = \langle \langle \rangle, \langle \rangle \rangle$$

(no objects available for anaphora)

• A UC₁-drs K is UC₁-true in M ,

written $M \models K$, iff $\forall g: c_0 \|K\|^g \neq \emptyset$

(i.e. iff the default state c_0 updated
with $\|K\|^g$ is not the absurd state).

• e.g. given that in M_0 , for any g ,

$$c_0 \| [\mathbf{x} \mid \text{man } \mathbf{x}]; \dots ; [\text{see}(\mathbf{T}, \perp)] \|^g$$

$$= \{ \langle \langle M \rangle, \langle M' \rangle \rangle \}$$

$\neq \emptyset$

UC₁ (1') is UC₁-true in the model M_0 ,

i.e. $M_0 \models (1')$.

• Empirically testable prediction:

English

(1) *A man saw an enemy.*

(represented by UC₁ drs

(1') $[\mathbf{x} \mid \text{man } \mathbf{x}]; ([y \mid \text{enm}^{of}(y, \mathbf{T})];$

$[\text{see}(\mathbf{T}, \perp)])$)

will be judged *intuitively true* in (the
situation represented by) the model M_0 .

Once upon a time in the N.

(1): sample computation

Model M_1 with $D_e = \{M, M', K, K'\}$

$\| \text{man} \| = \{M, M'\}$

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$\| \text{kayak} \| = \{K, K'\}$

(*pos* for 'inalienable-possessor-of')

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$\| \text{approach} \| = \{ \langle M, M' \rangle, \langle M, K' \rangle \}$

$\| \text{see} \| = \{ \langle M, M' \rangle, \langle M', M \rangle, \\ \langle M, K' \rangle, \langle M', K \rangle \}$

$\| \text{int.to.harpoon} \| = \{ \langle M, M' \rangle, \langle M', M \rangle \}$

$\| \text{forest} \| = \{ \langle M', M \rangle \}$

$\| \text{harpoon} \| = \{ \langle M', M \rangle \}$

$\| \text{capsize} \| = \{M\}$

$\| \text{die} \| = \{M\}$

$\| \text{kill} \| = \{ \langle M', M \rangle \}$

$\| \text{go.hml} \| = \{M'\}$

$\| \text{tell.story} \| = \{M'\}$

(1) man-have-DEC_{iv}-3S enemy-with-MOD

$[y \mid \text{man } y]; -\downarrow [y \mid \text{enm}^{of}(y, \perp)]$

1. $c_0 \| [y \mid \text{man } y] \|^g$

2. $c_0 \| \lambda \lambda j. \exists y \exists i (j = (y \cdot i) \wedge \text{li} \wedge \text{man } y) \|^g \quad \text{T2.iii}$

3. $\| \lambda \lambda j. \exists y \exists i (j = (y \cdot i) \wedge \text{li} \wedge \text{man } y) \|^g (c_0) \quad \text{D6.ln3}$

4. $\ast \langle \langle \mathbf{T}i, d \cdot \perp i \rangle \mid i \in c_0 \ \& \ d \in \emptyset \| \text{man} \| \rangle \quad \text{D6.sem}$

5. $\ast \langle \langle \mathbf{T}i, d \cdot \perp i \rangle \mid i \in \langle \langle \rangle, \langle \rangle \rangle \ \& \ d \in \emptyset \| \text{man} \| \rangle \quad \text{D7.c}_0$

6. $\ast \langle \langle \mathbf{T}i, d \cdot \perp i \rangle \mid i = \langle \langle \rangle, \langle \rangle \rangle \ \& \ d \in \emptyset \| \text{man} \| \rangle \quad \text{set thr}$

7. $\ast \langle \langle \rangle, \langle d \rangle \mid d \in \emptyset \| \text{man} \| \rangle \quad \text{elim. i}$

8. $\langle \langle \rangle, \langle M \rangle \rangle, \quad \text{df. } M_1, \| \text{man} \|$

$\langle \langle \rangle, \langle M' \rangle \rangle$

$=: c_1$

df. c_1

Once upon a time in the N.
Kal (1) → UC₁ syn, sem

Model M_1 with $D_e = \{M, M', K, K'\}$
 $llmanll = \lambda\{M, M'\}$
 $llenm^o|| = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $llkayakll = \lambda\{K, K'\}$
 (pos for 'inalienable-possessor-of')
 $llposll(K) = M, llposll(K') = M'$
 $llusel = \lambda\{\langle M, K \rangle, \langle M', K' \rangle\}$
 $llie.in.waill = \lambda\{M, M'\}$
 $llapproachll = \lambda\{\langle M, M' \rangle, \langle M, K' \rangle\}$
 $llseell = \lambda\{\langle M, M' \rangle, \langle M', M \rangle, \langle M, K' \rangle, \langle M', K \rangle\}$
 $llint.to.harpoonll = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
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 $lldiell = \lambda\{M\}$
 $llkillll = \lambda\{\langle M', M \rangle\}$
 $llgo.hmll = \lambda\{M'\}$
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(1) ...there was a man who had (an) enemy(ies)

angut-qar-pu-q akiraq-lik-mik
 man-have-DEC_{iv}-3S enemy-with-MOD
 [yl *man* y]; ↓ [yl *enm^{of}*(y, ⊥)]

⊥-upd: add a man test: ⊥-man has an enemy
 $C_0||[yl \textit{man} y]||^{\#}$ $C_1||\downarrow[yl \textit{enm}^{\textit{of}}(y, \perp)]||^{\#}$
 $= \lambda\{\langle \langle \rangle, \langle M \rangle \rangle, \langle \langle \rangle, \langle M' \rangle \rangle\}$ $= \lambda\{\langle \langle \rangle, \langle M \rangle \rangle, \langle \langle \rangle, \langle M' \rangle \rangle\}$
 $\equiv: C_1$ $\equiv: C_2$

Once upon a time in the N.
Kal (2²) → UC₁ (part 1)

Model M_1 with $D_e = \{M, M', K, K'\}$
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(2¹) (Once) when he^T was hunting in a kayak ...

$\lambda\{\langle \langle M \rangle, \langle K, M \rangle \rangle, \langle \langle M' \rangle, \langle K', M' \rangle \rangle\}$
 $\equiv: C_4$

2 he_T saw another kayak ...;

alla-mik qajaq-si-ga-mi
 other-MOD kayak-see-FCT_T-3S_T
 [yl *kayak* y]; [⊥₂ ∈ ⊥||]; [⊥ ≠ ⊥₂]; [see(T, ⊥)]; ...

⊥-upd: add a kayak global test: ⊥₂ is in ⊥-col
 $C_4||[yl \textit{kayak} y]||^{\#}$ $C_5||[\perp_2 \in \perp||]||^{\#}$
 $= \lambda\{\langle \langle M \rangle, \langle K, K, \dots \rangle \rangle, \langle \langle M \rangle, \langle K', K, \dots \rangle \rangle, \langle \langle M' \rangle, \langle K, K', \dots \rangle \rangle, \langle \langle M' \rangle, \langle K', K', \dots \rangle \rangle\}$
 $\equiv: C_5$ $\equiv: C_6$
 (i.e. the antecedent, ⊥₂, is a kayak)

Once upon a time in the N.
Kal (2¹) → UC₁ syn, sem

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(2¹) (Once) when he^T was hunting in a kayak ...

qajaq-tur-llu-ni
 kayak-use-ELA_T-3S_T
 [xl *x* = ⊥]; [yl *kayak* y, *use*(T, y)]; ...

T-upd: add ⊥ (i.e. promote ⊥ to T)
 $C_2||[xl \textit{x} = \perp]||^{\#}$
 $= \lambda\{\langle \langle M \rangle, \langle M \rangle \rangle, \langle \langle M' \rangle, \langle M' \rangle \rangle\}$
 $\equiv: C_3$

⊥-upd: add a kayak used by T
 $C_3||[yl \textit{kayak} y, \textit{use}(T, y)]||^{\#}$
 $= \lambda\{\langle \langle M \rangle, \langle K, M \rangle \rangle, \langle \langle M' \rangle, \langle K', M' \rangle \rangle\}$
 $\equiv: C_4$

Once upon a time in the N.
Kal (2²) → UC₁ (part 2)

Model M_1 with $D_e = \{M, M', K, K'\}$
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 $llkayakll = \lambda\{K, K'\}$
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 $llkillll = \lambda\{\langle M', M \rangle\}$
 $llgo.hmll = \lambda\{M'\}$
 $lltell.storyll = \lambda\{M'\}$

(2²) he_T saw another kayak ...;

alla-mik qajaq-si-ga-mi
 other-MOD kayak-see-FCT_T-3S_T
 [yl *kayak* y]; [⊥₂ ∈ ⊥||]; [⊥ ≠ ⊥₂]; [see(T, ⊥)]; ...

⊥-upd: add a kayak global test: ⊥₂ is in ⊥-col
 $C_4||[yl \textit{kayak} y]||^{\#}$ $C_5||[\perp_2 \in \perp||]||^{\#}$
 $= \lambda\{\langle \langle M \rangle, \langle K, K, \dots \rangle \rangle, \langle \langle M \rangle, \langle K', K, \dots \rangle \rangle, \langle \langle M' \rangle, \langle K, K', \dots \rangle \rangle, \langle \langle M' \rangle, \langle K', K', \dots \rangle \rangle\}$
 $\equiv: C_5$ $\equiv: C_6$

local test: ⊥ is not ⊥₂ local test: T saw ⊥
 $C_6||[\perp \neq \perp_2]||^{\#}$ $C_7||[\textit{see}(T, \perp)]||^{\#}$
 $= \lambda\{\langle \langle M \rangle, \langle K', K, \dots \rangle \rangle, \langle \langle M' \rangle, \langle K, K', \dots \rangle \rangle\}$ $\langle \langle M \rangle, \langle K', K, \dots \rangle \rangle, \langle \langle M' \rangle, \langle K, K, \dots \rangle \rangle\}$
 $\equiv: C_7$ $\equiv: C_8$

Once upon a time in the N.
Kal (2³) → UC₁ syn, sem

Model M_1 with $D_e = \{M, M', K, K'\}$
 $\llbracket man \rrbracket = \lambda\{M, M'\}$
 $\llbracket enm^{of} \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $\llbracket kayak \rrbracket = \lambda\{K, K'\}$
 (pos for 'inalienable-possessor-of')
 $\llbracket pos \rrbracket(K) = M, \llbracket pos \rrbracket(K') = M'$
 $\llbracket use \rrbracket = \lambda\{\langle M, K \rangle, \langle M', K' \rangle\}$
 $\llbracket lie.in.wait \rrbracket = \lambda\{M, M'\}$
 $\llbracket approach \rrbracket = \lambda\{\langle M, M' \rangle, \langle M, K' \rangle\}$
 $\llbracket see \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle, \langle M, K' \rangle, \langle M', K \rangle\}$
 $\llbracket int.to.harpoon \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $\llbracket forest \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket harpoon \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket capsizell \rrbracket = \lambda\{M\}$
 $\llbracket die \rrbracket = \lambda\{M\}$
 $\llbracket kill \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket go.hm \rrbracket = \lambda\{M'\}$
 $\llbracket tell.story \rrbracket = \lambda\{M'\}$

(2²) he_T saw another kayak ...;
alla-mik qajaq-si-ga-mi
 other-MOD kayak-see-FCT_T-3S_T
 $[y \mid kayak \ y]; [\perp_2 \in \perp \rrbracket]; [\perp \neq \perp_2]; [see\langle T, \perp \rangle]; \dots$
 $C_7 \llbracket see\langle T, \perp \rangle \rrbracket \rrbracket^g$
 $= \lambda\{\langle\langle M \rangle, \langle K', K, \dots \rangle\rangle, \langle\langle M' \rangle, \langle K, K, \dots \rangle\rangle\}$
 $\therefore C_8$
 3 approaching it[⊥] ...
umig-llu-gu
 approach-ELA_T-3S_⊥
 $[approach\langle T, \perp \rangle, T \neq \perp]$
local test: T approached ⊥ & T is not ⊥
 $C_8 \llbracket approach\langle T, \perp \rangle, T \neq \perp \rrbracket \rrbracket^g$
 $= \lambda\{\langle\langle M \rangle, \langle K', K, \dots \rangle\rangle, \langle\langle M' \rangle, \langle K, K, \dots \rangle\rangle\}$
 $\therefore C_9$

Once upon a time in the N.
Kal (6) → UC₁ → truth

Model M_1 with $D_e = \{M, M', K, K'\}$
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 $\llbracket enm^{of} \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $\llbracket kayak \rrbracket = \lambda\{K, K'\}$
 (pos for 'inalienable-possessor-of')
 $\llbracket pos \rrbracket(K) = M, \llbracket pos \rrbracket(K') = M'$
 $\llbracket use \rrbracket = \lambda\{\langle M, K \rangle, \langle M', K' \rangle\}$
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 $\llbracket forest \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket harpoon \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket capsizell \rrbracket = \lambda\{M\}$
 $\llbracket die \rrbracket = \lambda\{M\}$
 $\llbracket kill \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket go.hm \rrbracket = \lambda\{M'\}$
 $\llbracket tell.story \rrbracket = \lambda\{M'\}$

(see handout for (4)-(5))
 (6) ...the killer^T went home and told the story.
tuqut-si-tuq-Ø angirlar-ga-mi uqaluttuar-pu-q.
 kill-antip-cn\iv-T go.hm-FCT_T-3S_T tell.story-DEC_{iv}-3S_⊥
 $[kill\langle T_2, \perp \rangle]; [x \mid x = T_2]; [go.hm \ T]; [tell.story \ T]$
 $C_{25} \llbracket kill\langle T_2, \perp \rangle; [x \mid x = T_2] \rrbracket \rrbracket^g$
 $= \{\langle\langle M', M, \dots \rangle, \langle M, M', \dots \rangle\rangle\}$
 $\therefore C_{27}$
 $C_{27} \llbracket go.hm \ T; [tell.story \ T] \rrbracket \rrbracket^g$
 $= \{\langle\langle M', M, \dots \rangle, \langle M, M', \dots \rangle\rangle\}$
 $\therefore C_{29}$

- CONCLUSION: Since $C_{29} \neq \emptyset$ this UC₁ drs is UC₁-true in the given model
- PREDICTION: The Kalaallisut text (represented by this UC₁ drs) will be judged *intuitively true* in (the situation represented by) the given model.

Once upon a time in the N.
Kal (2^{4,5}) → UC₁ syn, sem

Model M_1 with $D_e = \{M, M', K, K'\}$
 $\llbracket man \rrbracket = \lambda\{M, M'\}$
 $\llbracket enm^{of} \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $\llbracket kayak \rrbracket = \lambda\{K, K'\}$
 (pos for 'inalienable-possessor-of')
 $\llbracket pos \rrbracket(K) = M, \llbracket pos \rrbracket(K') = M'$
 $\llbracket use \rrbracket = \lambda\{\langle M, K \rangle, \langle M', K' \rangle\}$
 $\llbracket lie.in.wait \rrbracket = \lambda\{M, M'\}$
 $\llbracket approach \rrbracket = \lambda\{\langle M, M' \rangle, \langle M, K' \rangle\}$
 $\llbracket see \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle, \langle M, K' \rangle, \langle M', K \rangle\}$
 $\llbracket int.to.harpoon \rrbracket = \lambda\{\langle M, M' \rangle, \langle M', M \rangle\}$
 $\llbracket forest \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket harpoon \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket capsizell \rrbracket = \lambda\{M\}$
 $\llbracket die \rrbracket = \lambda\{M\}$
 $\llbracket kill \rrbracket = \lambda\{\langle M', M \rangle\}$
 $\llbracket go.hm \rrbracket = \lambda\{M'\}$
 $\llbracket tell.story \rrbracket = \lambda\{M'\}$

(2^{4,5}) he_T saw (it was) his_T enemy[⊥] lying in wait.
taku-pa-a-Ø akiraq-ni-Ø qama-tu-q
 see-DEC_{iv}-3S-3S enm-3S_T-[⊥] lie.in.wait-ELA_{iv}-3S_⊥
 $[enm^{of}\langle pos \ \perp, T \rangle]; [y \mid y = pos \ \perp]; [see\langle T, \perp \rangle, T \neq \perp]$
 $[\text{lie.in.wait} \ \perp]$
 $C_9 \llbracket enm^{of}\langle pos \ \perp, T \rangle \rrbracket \rrbracket^g$
 $= \{\langle\langle M \rangle, \langle K', K, \dots \rangle\rangle\}$
 $\therefore C_{10}$
 $C_{10} \llbracket [y \mid y = pos \ \perp] \rrbracket \rrbracket^g$
 $= \{\langle\langle M \rangle, \langle M', K', \dots \rangle\rangle\}$
 $\therefore C_{11}$
 $C_{11} \llbracket see\langle T, \perp \rangle, T \neq \perp \rrbracket \rrbracket^g$
 $= \{\langle\langle M \rangle, \langle M', K', \dots \rangle\rangle\}$
 $\therefore C_{12}$
 $C_{12} \llbracket [\text{lie.in.wait} \ \perp] \rrbracket \rrbracket^g$
 $= \{\langle\langle M \rangle, \langle M', K', \dots \rangle\rangle\}$
 $\therefore C_{13}$