

**Encapsulated quantification:  
Stone 1997, 1999**

## 1. DATA &amp; BASIC IDEAS

• Referential parallels*Deictic reference*

- (1) *She*<sub>u1</sub> left me. *nominal* (Partee 1973)  
[uttered by a man sitting alone with his head in his hands]
- (2) My neighbours *would*<sub>ω1</sub> kill me. *modal* (Stone 1997)  
[looking at a high-end stereo in an electronics store]

*Anaphoric reference within a clause*

- (3<sup>1</sup>) Pedro owns *a*<sup>u2</sup> [donkey]<sub>u2</sub> *nominal*
- (4<sup>1</sup>) I *must*<sub>ω</sub><sup>ω1</sup> [buy a car]<sub>ω1</sub> *modal* (Stone 1997)

*Anaphoric reference in discourse*

- (3) Pedro<sup>u1</sup> owns *a*<sup>u2</sup> donkey. *He*<sub>u1</sub> beats *it*<sub>u2</sub>. *nominal*
- (4) I *must*<sub>ω</sub><sup>ω1</sup> buy a car. But I *can't*<sub>ω, ω1</sub>. *modal* (Stone 1997)

*Anaphoric reference in conditionals*

- (5) *If*<sub>ω</sub><sup>ω1</sup> *a*<sup>u1</sup> man commits a murder, ***weak rdg. only!***  
*he*<sub>u1</sub> *should*<sub>ω, ω1</sub><sup>ω2</sup> go to prison.

*Nominal anaphora across modalities*

- (6) *If*<sub>ω</sub><sup>ω1</sup> there are bears in this area, *a*<sup>u2</sup> bear *might*<sub>ω, ω1</sub><sup>ω2</sup> come by. (Stone 1999)  
But *he*<sub>u2</sub> *wouldn't*<sub>ω, ω2</sub><sup>ω3</sup> hurt you. I **have**<sub>ω</sub> *a*<sup>u3</sup> gun.  
*If*<sub>ω, ω2</sub><sup>ω5</sup> *the*<sub>u2</sub> bear attacked you, I **would**<sub>ω, ω5</sub><sup>ω6</sup> use *it*<sub>u3</sub> to shoot *him*<sub>u2</sub>.

• Stone's (1997, 1999) theory

Adapting and extending Muskens (1995), Stone (1997, 1999) proposes:

- (a) parallel analysis of nominal and modal reference in terms of drefs for *individual concepts* (type *swe*) and *modal concepts* (type *sw(wt)*).
- (b) nouns and verbs relate their nominal arguments to a modal domain (*encapsulated quant.*  $\forall w$ )
- (c) modals (*may*, *would*, etc) do not quantify over worlds; instead they relate modal drefs.
- (d) representation of reality:

“Assume that the initial environment *i* assigns a value to  $\omega$  corresponding to the information implicitly available to the participants in the conversation, which we can represent by a proposition or set of worlds *E*. Formally, then, for each world  $w \in E$ ,  $\omega(i, w) = E$ , and for any other world  $\omega(i, w)$  is empty.” (Stone 1999:20)

## NOTATION:

<p><math>a^*</math> for the speaker (type <i>e</i>)</p> <p><math>p^*</math> for <i>E</i> of Stone 1999:20 (type <i>wt</i>)</p> <p><i>M</i> for modal base functions (type <i>w(wt)</i>)</p> <p><math>w' &lt;_w w'' := (w' \leq_w w'' \wedge \neg w'' \leq_w w')</math></p> <p><math>\min_{&lt;_w} p := \{w' \in p \mid \neg \exists w'' \in p (w'' &lt;_w w')\}</math></p> <p><math>\lambda w \in p. f_w := \{\langle w, f_w \rangle : w \in p\} \cup \{\langle w, \emptyset \rangle : w \notin p\}</math></p>	<p><math>u_n</math> for nominal drefs (type <i>swe</i>)</p> <p><math>u := \lambda i. (\lambda w. a^*)</math></p> <p><math>\omega_n</math> for modal drefs (type <i>sw(wt)</i>)</p> <p><math>\omega := \lambda i. (\lambda w \in p^*. p^*)</math></p> <p><math>\text{Dom } M := \{w \mid Mw \neq \emptyset\}</math></p> <p><math>\text{Ran } M := \{w' \mid \exists w (w' \in Mw)\}</math></p>
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## 2. DEICTIC REFERENCE

(1) *She* left me.(1<sup>0</sup>) Start-up

a la Bittner 2007

$$i_0$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$
(1<sup>1</sup>) Accommodate presupposition of  $\text{she}_{u_1}$ 

$$i_1$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$u_1 \quad \lambda w. a_1$$
(1<sup>2</sup>)  $\text{She}_{u_1}$ 

$$[\text{female}_{\omega}\{u_1\}]i_1i_1$$

$$\text{iff } \forall w \in \text{Ran } \omega i_1(\text{female}_w(u_1i_1w))$$

$$\text{iff } \forall w \in p^*(\text{female}_w(u_1i_1w))$$

[In every  $p^*$ -world,  $a_1$  is female.]

The past tense *left* further adds the presupposition that  $a_1$  used to be with Bill, but this cannot be represented in this tenseless system.

(1<sup>3</sup>)  $\text{left}_{\omega}$  me.
$$[\text{leave}_{\omega}\{u_1, u\}]i_1i_1$$

$$\text{iff } \forall w \in p^*(\text{leave}_w(a_1, a^*))$$

[In every  $p^*$ -world,  $a_1$ -female leaves the speaker  $a^*$ .]

(2) My neighbours *would* kill me.

[uttered while looking at high-end stereo,  $c$ , in electronics store]

(2<sup>0</sup>) Start-up

$i_0$   
 $u$   $\lambda w. a^*$   
 $u_1$   $\lambda w. c$   
 $\omega$   $\lambda w \in p^*. p^*$

(2<sup>1</sup>) Accommodate presup. of *would* <sub>$\omega, \omega_1$</sub>  <sup>$\omega_2$</sup>  by pragmatic inference (*pfb* = play for full blast)

$\max_{\omega}^{\omega_1}([\text{buy}_{\omega_1}\{u, u_1\}, \text{pfb}_{\omega_1}\{u, u_1\}])i_0i_1$   
iff  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in p^*(\omega_1i_1w = \{w' \mid \text{buy}_w(a^*, c) \wedge \text{pfb}_w(a^*, c)\})$   
 $i_1$   
 $u$   $\lambda w. a^*$   
 $u_1$   $\lambda w. c$   
 $\omega$   $\lambda w \in p^*. p^*$   
 $\omega_1$   $\lambda w \in p^*. \mathbf{B}$  s.t.  $\bullet \mathbf{B} := \{w' \mid \text{buy}_w(a^*, c) \wedge \text{pfb}_w(a^*, c)\}$

(2<sup>2</sup>) My neighbours

$[\omega_1: u_2 \mid \text{neighbours.of}_{\omega_1}\{u_2, u\}]i_1i_2$   
iff  $i_1[\omega_1: u_2]i_2 \wedge \text{Dom } u_2i_2 = \mathbf{B} \wedge \forall w \in \mathbf{B}(\text{neighbours.of}_w(u_2i_2w, ui_2w))$   
 $i_2$   
 $u$   $(a^*)^\circ$   
 $u_1$   $\lambda w. c$   
 $u_2$   $\lambda w \in \mathbf{B}. b_w$  s.t.  $\bullet \forall w \in \mathbf{B}(\text{neighbours.of}_w(b_w, a^*))$   
 $\omega$   $\lambda w \in p^*. p^*$   
 $\omega_1$   $\lambda w \in p^*. \mathbf{B}$  s.t.  $\bullet \mathbf{B} := \{w' \mid \text{buy}_w(a^*, c) \wedge \text{pfb}_w(a^*, c)\}$

(2<sup>3</sup>) *would* <sub>$\omega, \omega_1$</sub>  <sup>$\omega_2$</sup>  kill me.

*Notation:* ' $<$ ' for ranking by similarity to  $w$ , ' $<_1$ ' for ranking  $a^*$ 's expectations in  $w$  ( $[\mathbf{remote}_{<_1}(\omega, \omega_1)]$ );

$\max_{\omega}^{\omega_2}([\omega_2 \subseteq \omega_1, \text{kill}_{\omega_2}\{u_2, u\}]; [\mathbf{necc}_{<}(\omega_1, \omega_2)])i_2i_3$   
iff  $\forall w \in p^*((\mathbf{min}_{<_1, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $i_2[\omega: \omega_2]i_3 \wedge \text{Dom } \omega_2i_3 = p^* \wedge \forall w \in p^*(\omega_2i_3w = \{w' \in \mathbf{B} \mid \text{kill}_w(b_{w'}, a^*)\})$   
 $\wedge \forall w \in p^*((\mathbf{min}_{<, w} \mathbf{B}) \subseteq \mathbf{K})$   
 $i_2$   
 $u$   $(a^*)^\circ$   
 $u_1$   $\lambda w. c$   
 $u_2$   $\lambda w \in \mathbf{B}. b_w$  s.t.  $\bullet \forall w \in \mathbf{B}(\text{neighbours.of}_w(b_w, a^*))$   
 $\omega$   $\lambda w \in p^*. p^*$   
 $\omega_1$   $\lambda w \in p^*. \mathbf{B}$  s.t.  $\bullet \mathbf{B} := \{w' \mid \text{buy}_w(a^*, c) \wedge \text{pfb}_w(a^*, c)\}$   
 $\bullet \forall w \in p^*((\mathbf{min}_{<_1, w} p^*) \cap \mathbf{B} \neq \emptyset)$   
 $\omega_2$   $\lambda w \in p^*. \mathbf{K}$  s.t.  $\bullet \mathbf{K} := \{w' \in \mathbf{B} \mid \text{kill}_w(b_{w'}, a^*)\}$   
 $\bullet \forall w \in p^*((\mathbf{min}_{<, w} \mathbf{B}) \subseteq \mathbf{K})$

[For every  $p^*$ -world  $w$ , the most  $<_{1, w}$ -likely  $p^*$ -worlds are not  $\mathbf{B}$ -worlds, and in every  $w$ -closest  $\mathbf{B}$ -world  $w'$ , where the speaker  $a^*$  buys the stereo  $c$  and plays it for full blast, he is killed by his  $w'$ -neighbours  $b_{w'}$ .]

## 3. ANAPHORIC REFERENCE

(3) Pedro owns a donkey. *He beats it.**nominal anaphora*(3<sup>0</sup>) Start-up
$$i_0$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$
(3<sup>1</sup>) Pedro
$$[\omega: u_1 | u_1 = \text{Pedro}, \text{male}_\omega \{u_1\}] i_0 i_1$$

$$\text{iff } i_0[\omega: u_1] i_1 \wedge \text{Dom } u_1 i_1 = p^* \wedge \forall w \in p^* (u_1 i_1 w = \text{pedro} \wedge \text{male}_w u_1 i_1 w)$$

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in p^*. \text{pedro} \quad \text{s.t. } \forall w \in p^* (\text{male}_w \text{pedro})$$

$$\omega \quad \lambda w \in p^*. p^*$$
(3<sup>2</sup>) owns
$$[\omega: u_2 | \text{own}_\omega \{u_1, u_2\}] i_1 i_2$$

$$\text{iff } i_1[\omega: u_2] i_2 \wedge \text{Dom } u_2 i_2 = p^* \wedge \forall w \in p^* (\text{own}_w (\text{pedro}, u_2 i_2 w))$$

$$i_2$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in p^*. \text{pedro} \quad \text{s.t. } \forall w \in p^* (\text{male}_w \text{pedro})$$

$$u_2 \quad \lambda w \in p^*. b_w \quad \text{s.t. } \forall w \in p^* (\text{own}_w (\text{pedro}, b_w))$$

$$\omega \quad \lambda w \in p^*. p^*$$
(3<sup>3</sup>) a donkey
$$[\text{donkey}_\omega \{u_2\}] i_2 i_2$$

$$\text{iff } \forall w \in p^* (\text{donkey}_w (b_w))$$

$$i_2$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in p^*. \text{pedro} \quad \text{s.t. } \forall w \in p^* (\text{male}_w \text{pedro})$$

$$u_2 \quad \lambda w \in p^*. b_w \quad \text{s.t. } \forall w \in p^* (\text{own}_w (\text{pedro}, b_w) \wedge \text{donkey}_w b_w)$$

$$\omega \quad \lambda w \in p^*. p^*$$
(3<sup>4</sup>) He beats it
$$[\text{male}_\omega \{u_1\}, \text{beat}_\omega \{u_1, u_2\}, u_2 \neq_\omega u_1] i_2 i_2$$

$$\text{iff } \forall w \in p^* (\text{male}_w (\text{pedro}) \wedge \text{beat}_w (\text{pedro}, b_w) \wedge b_w \neq \text{pedro})$$

$$i_2$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in p^*. \text{pedro} \quad \text{s.t. } \forall w \in p^* (\text{male}_w \text{pedro})$$

$$u_2 \quad \lambda w \in p^*. b_w \quad \text{s.t. } \forall w \in p^* (\text{own}_w (\text{pedro}, b_w) \wedge \text{donkey}_w b_w \wedge \text{beat}_w (\text{pedro}, b_w) \wedge b_w \neq \text{pedro})$$

$$\omega \quad \lambda w \in p^*. p^*$$

(4) I must buy a car. But I can't.

*modal anaphora*(4<sup>0</sup>) Start-up
$$i_0$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$
(4<sup>1</sup>) [I \_ buy a car] <sub>$\omega$</sub>  <sup>$\omega_1$</sup> 

$$\mathbf{max}_{\omega}^{\omega_1}([\omega_1 \subseteq \omega]; [\omega: u_1 | \text{car}_{\omega_1} \{u_1\}, \text{buy}_{\omega_1} \{u, u_1\}])i_0i_1$$

iff  $\exists k(i_0[\omega: \omega_1]k \wedge k[\omega: u_1]i_1$   
 $\wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in p^*(\omega_1i_1w = \{w' \in p^* | \exists y(\text{car}_{w'}y \wedge \text{buy}_{w'}(a^*, y))\}$   
 $\wedge \text{Dom } u_1i_1 = \text{Ran } \omega_1i_1 \wedge \forall w' \in \text{Ran } \omega_1i_1(\text{car}_{w'}u_1i_1w' \wedge \text{buy}_{w'}(a^*, u_1i_1w'))$

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{B}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(\text{car}_{w'}c_w \wedge \text{buy}_{w'}(a^*, c_w))$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \in p^* | \exists y(\text{car}_{w'}y \wedge \text{buy}_{w'}(a^*, y))\}$$
(4<sup>2</sup>) must <sub>$\omega, \omega_1$</sub> 

*Notation:* ' $\prec_2$ ' for ranking by  $a^*$ 's needs in  $w$   
 ([**vid** <sub>$\prec_2$</sub> ( $\omega, \omega_1$ ); [**necc** <sub>$\prec_2$</sub> ( $\omega, \omega_1$ ))]  $i_1i_1$

iff  $\forall w \in p^*((\mathbf{min}_{\prec_2, w} p^*) \cap \mathbf{B} \neq \emptyset)$ iff  $\forall w \in p^*((\mathbf{min}_{\prec_2, w} p^*) \subseteq \mathbf{B})$ iff  $\forall w \in p^*(\emptyset \subset (\mathbf{min}_{\prec_2, w} p^*) \subseteq \mathbf{B})$ 

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{B}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(\text{car}_{w'}c_w \wedge \text{buy}_{w'}(a^*, c_w))$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \in p^* | \exists y(\text{car}_{w'}y \wedge \text{buy}_{w'}(a^*, y))\}$$

$$\bullet \forall w \in p^*(\emptyset \subset (\mathbf{min}_{\prec_2, w} p^*) \subseteq \mathbf{B})$$
(4<sup>3</sup>) But I can't <sub>$\omega, \omega_1$</sub> 

*Notation:* ' $\prec_1$ ' for ranking by  $a^*$ 's expectations in  $w$   
 ([**vid** <sub>$\prec_2$</sub> ( $\omega, \omega_1$ ); [**not** <sub>$\prec_1$</sub> ( $\omega, \omega_1$ ))]  $i_1i_1$

iff  $\forall w \in p^*((\mathbf{min}_{\prec_2, w} p^*) \cap \omega_1i_1w \neq \emptyset)$ iff  $\forall w \in p^*((\mathbf{min}_{\prec_1, w} p^*) \cap \omega_1i_1w = \emptyset)$ 

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{B}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(\text{car}_{w'}c_w \wedge \text{buy}_{w'}(a^*, c_w))$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \in p^* | \exists y(\text{car}_{w'}y \wedge \text{buy}_{w'}(a^*, y))\}$$

$$\bullet \forall w \in p^*(\emptyset \subset (\mathbf{min}_{\prec_2} p^*) \subseteq \mathbf{B})$$

$$\bullet \forall w \in p^*((\mathbf{min}_{\prec_1} p^*) \cap \mathbf{B} = \emptyset)$$

## 4. DONKEY CONDITIONALS: WEAK READING ONLY

(5) If a man commits a murder, he should go to prison.

(5<sup>0</sup>) Start-up
$$i_0$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$
(5<sup>1</sup>) If <sub>$\omega$</sub>  <sup>$\omega_1$</sup>  a <sup>$u_1$</sup>  man commits a murderNotation: *mrd* = commit a murder
$$\mathbf{max}_{\omega}^{\omega_1}([\omega_1: u_1 | \mathbf{man}_{\omega_1}\{u_1\}, \mathbf{mrd}_{\omega_1}\{u_1\}])i_0i_1$$

iff  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in p^*(\omega_1i_1w = \{w' | \exists x(\mathbf{man}_{w'}x \wedge \mathbf{mrd}_{w'}x)\}$   
 $\wedge \text{Dom } u_1i_1 = \text{Ran } \omega_1i_1 \wedge \forall w \in \text{Ran } \omega_1i_1(\mathbf{man}_{w'}u_1i_1w' \wedge \mathbf{mrd}_{w'}u_1i_1w')$

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{M}. b_w \quad \text{s.t. } \forall w' \in \mathbf{M}(\mathbf{man}_{w'}b_{w'} \wedge \mathbf{mrd}_{w'}b_{w'})$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{M} \quad \text{s.t. } \mathbf{M} = \{w' | \exists x(\mathbf{man}_{w'}x \wedge \mathbf{mrd}_{w'}x)\}$$
(5<sup>2</sup>) he <sub>$u_1$</sub>  should <sub>$\omega, \omega_1, \omega_2$</sub>  go to prisonNotation: *gtp* = go to prison'<<sub>3</sub>' for ranking by conformity to the laws of  $w$ 

$$([\mathbf{remote}_{<3}(\omega, \omega_1)];$$

$$\mathbf{max}_{\omega}^{\omega_2}([\omega_2 \subseteq \omega_1, \mathbf{gtp}_{\omega_2}\{u_1\}]);$$

$$[\mathbf{necc}_{<3}(\omega_1, \omega_2)]i_1i_2$$
*presup. test on  $i_1$* *update to  $i_2$* *test  $i_2$* 

iff  $\forall w \in p^*((\mathbf{min}_{<3, w} p^*) \cap \mathbf{M} = \emptyset)$   
 $\wedge i_1[\omega: \omega_2]i_2 \wedge \text{Dom } \omega_2i_2 = p^* \wedge \forall w \in p^*(\omega_2i_2w = \{w' \in \mathbf{B} | \mathbf{gtp}_{w'}b_{w'}\})$   
 $\wedge \forall w \in p^*((\mathbf{min}_{<3, w} \mathbf{M}) \subseteq \omega_2i_2w)$

$$i_2$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{M}. b_w \quad \text{s.t. } \bullet \forall w' \in \mathbf{M}(\mathbf{man}_{w'}b_{w'} \wedge \mathbf{mrd}_{w'}b_{w'})$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{M} \quad \text{s.t. } \bullet \mathbf{M} = \{w' | \exists x(\mathbf{man}_{w'}x \wedge \mathbf{mrd}_{w'}x)\}$$

$$\bullet \forall w \in p^*((\mathbf{min}_{<3, w} p^*) \cap \mathbf{M} = \emptyset)$$

[UNREALISTIC: in the most law-abiding worlds, among the speaker's doxastic alternatives  $p^*$ , there is no murder]

$$\omega_2 \quad \lambda w \in p^*. \mathbf{P} \quad \text{s.t. } \bullet \mathbf{P} = \{w' \in \mathbf{M} | \mathbf{gtp}_{w'}b_{w'}\}$$

$$\bullet \forall w \in p^*((\mathbf{min}_{<3, w} \mathbf{M}) \subseteq \mathbf{P})$$

[BAD: in  $\mathbf{M}$ -worlds  $w'$ , with one or more murderers, the laws of  $w$  require that at least one murderer,  $b_{w'}$ , go to prison—oops!!...]

## 5. HOW TO SHOOT A HYPOTHETICAL BEAR WITH A REAL GUN

(6a) If there are bears in this area, a bear might come by.

(6<sup>0</sup>) Start-up
$$i_0$$

$$u \quad \lambda w. a^*$$

$$\omega \quad \lambda w \in p^*. p^*$$
(6<sup>1</sup>) If <sub>$\omega$</sub>  <sup>$\omega_1$</sup>  there are bears in this area
$$\mathbf{max}_{\omega}^{\omega_1}([\omega_1: u_1 | bear_{\omega_1}\{u_1\}, near_{\omega_1}\{u_1, u\}])i_0i_1$$

iff  $i_0[\omega: \omega_1]i_1$   
 $\wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in p^*(\omega_1i_1w = \{w' | \exists x(bear_{w'}x \wedge near_{w'}(x, a^*))\})$   
 $\wedge \text{Dom } u_1i_1 = \text{Ran } \omega_1i_1 \wedge \forall w' \in \text{Ran } \omega_1i_1(bear_{w'}u_1i_1w' \wedge near_{w'}(u_1i_1w', a^*))$

$$i_1$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t. } \forall w' \in \mathbf{B}(bear_{w'}b_w \wedge near_{w'}b_w)$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t. } \mathbf{B} = \{w' | \exists x(bear_{w'}x \wedge near_{w'}(x, a^*))\}$$
(6<sup>2</sup>) a<sup>u2</sup> bear might <sub>$\omega, \omega_1, \omega_2$</sub>  come byNotation: ' $<_4$ ' for ranking by  $a^*$ 's hopes in  $w$ , ' $<_1$ ' for ranking by  $a^*$ 's expectations in  $w$ 

$$([\mathbf{remote}_{<_4}(\omega, \omega_1)]); \quad \text{presup. test on } i_1$$

$$\mathbf{max}_{\omega}^{\omega_2}([\omega_2 \subseteq \omega_1]; [\omega_2: u_2 | bear_{\omega_2}\{u_2\}, come_{\omega_2}\{u_2\}]); \quad \text{update to } i_2$$

$$[\mathbf{poss}_{<_1}(\omega_1, \omega_2)]i_1i_2 \quad \text{test } i_2$$

iff  $\forall w \in p^*((\mathbf{min}_{<_4, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $\wedge i_1[\omega: \omega_2]i_2 \wedge \text{Dom } \omega_2i_2 = p^* \wedge \forall w \in p^*(\omega_2i_2w = \{w' \in \mathbf{B} | \exists y(bear_{w'}y \wedge come_{w'}y)\})$   
 $\wedge \text{Dom } u_2i_2 = \text{Ran } \omega_2i_2 \wedge \forall w \in \text{Ran } \omega_2i_2(bear_w u_2i_2w \wedge come_w u_2i_2w)$   
 $\wedge \forall w \in p^*((\mathbf{min}_{<_1, w} \mathbf{B}) \cap \omega_2i_2w \neq \emptyset)$

$$i_2$$

$$u \quad \lambda w. a^*$$

$$u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t. } \bullet \forall w' \in \mathbf{B}(bear_{w'}b_w \wedge near_{w'}b_w)$$

$$u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t. } \bullet \forall w' \in \mathbf{C}(bear_{w'}c_w \wedge come_{w'}c_w)$$

$$\omega \quad \lambda w \in p^*. p^*$$

$$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t. } \bullet \mathbf{B} = \{w' | \exists x(bear_{w'}x \wedge near_{w'}(x, a^*))\}$$

$$\bullet \forall w \in p^*((\mathbf{min}_{<_4, w} p^*) \cap \mathbf{B} = \emptyset)$$

$$\text{[in } w, a^* \text{ hopes there are no bears nearby]}$$

$$\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t. } \bullet \mathbf{C} = \{w' \in \mathbf{B} | \exists y(bear_{w'}y \wedge come_{w'}y)\}$$

$$\bullet \forall w \in p^*((\mathbf{min}_{<_1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset)$$

$$\text{[in } w, a^* \text{'s expected } \mathbf{B}\text{-worlds include some } \mathbf{C}\text{-worlds.]}$$

(6b) But he wouldn't hurt you.

(6<sup>3</sup>)  $he_{u_2}$  wouldn't <sub>$\omega, \omega_2, \omega_3$</sub>  hurt you

*Notation:* ' $<_4$ ' for ranking by  $a^*$ 's hopes in  $w$ ,

' $<_1$ ' for ranking by  $a^*$ 's expectations in  $w$

([remote <sub>$<_4$</sub> ( $\omega, \omega_2$ )];

$\max_{\omega}^{\omega_3}([\omega_3 \subseteq \omega_2]; [hurt_{\omega_3}\{u_3, \mathbf{adr}\{u\}\}])$

[not <sub>$<_1$</sub> ( $\omega_2, \omega_3$ )] $i_2i_3$

*presup. test on  $i_2$*

*update to  $i_3$*

*test  $i_3$*

iff  $\forall w \in p^*((\min_{<_4, w} p^*) \cap C = \emptyset$

$\wedge i_2[\omega: \omega_3]i_3 \wedge \text{Dom } \omega_3i_3 = p^* \wedge \forall w \in p^*(\omega_3i_3w = \{w' \in C \mid hurt_w(c_{w'}, \mathbf{adr } a^*)\})$

$\wedge \forall w \in p^*((\min_{<_1, w} C) \cap \omega_3i_3w = \emptyset)$

$i_3$

$u \quad \lambda w. a^*$

$u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(bear_{w'} b_{w'} \wedge near_{w'} b_{w'})$

$u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{C}(bear_{w'} c_{w'} \wedge come_{w'} c_{w'})$

$\omega \quad \lambda w \in p^*. p^*$

$\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \mid \exists x(bear_{w'} x \wedge near_{w'}(x, a^*))\}$

$\bullet \forall w \in p^*((\min_{<_4, w} p^*) \cap \mathbf{B} = \emptyset$

[in  $w$ ,  $a^*$  hopes there are no bears nearby]

$\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t.} \quad \bullet \mathbf{C} = \{w' \in \mathbf{B} \mid \exists y(bear_{w'} y \wedge come_{w'} y)\}$

$\bullet \forall w \in p^*((\min_{<_1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset$

[in  $w$ ,  $a^*$ 's expected  $\mathbf{B}$ -worlds include some  $\mathbf{C}$ -worlds.]

$\bullet \forall w \in p^*((\min_{<_4, w} \mathbf{B}) \cap \mathbf{C} = \emptyset$

[in  $w$ ,  $a^*$ 's hoped for  $\mathbf{B}$ -worlds do not include any  $\mathbf{C}$ -worlds.]

$\omega_3 \quad \lambda w \in p^*. \mathbf{H} \quad \text{s.t.} \quad \bullet \mathbf{H} = \{w' \in \mathbf{C} \mid hurt_w(c_{w'}, \mathbf{adr } a^*)\}$

$\bullet \forall w \in p^*((\min_{<_1, w} \mathbf{C}) \cap \mathbf{H} = \emptyset$

[in  $w$ ,  $a^*$ 's expected  $\mathbf{C}$ -worlds do not include any  $\mathbf{H}$ -worlds.]

(6c) I (do) have a gun.

(6<sup>4</sup>) [I \_ have a<sup>u3</sup> gun]<sub>ω</sub><sup>ω4</sup>

$\mathbf{max}_{\omega}^{\omega4}([\omega_4: u_3 | \mathit{gun}_{\omega4}\{u_3\}, \mathit{have}_{\omega4}\{u, u_3\}])i_3i_4$

iff  $i_3[\omega: \omega_4]i_4 \wedge \text{Dom } \omega_4i_4 = p^* \wedge \forall w \in p^*(\omega_4i_4w = \{w' | \exists y(\mathit{gun}_{w'}y \wedge \mathit{have}_w(a^*, y))\})$   
 $\wedge \text{Dom } i_3i_4 = \text{Ran } \omega_4i_4 \wedge \forall w' \in \text{Ran } \omega_4i_4(\mathit{gun}_{w'}u_3i_4w' \wedge \mathit{have}_w(a^*, u_3i_4w'))$

$i_4$   
 $u \quad \lambda w. a^*$   
 $u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(\mathit{bear}_{w'}b_w \wedge \mathit{near}_{w'}b_w)$   
 $u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{C}(\mathit{bear}_{w'}c_w \wedge \mathit{come}_{w'}c_w)$   
 $u_3 \quad \lambda w \in \mathbf{G}. d_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{G}(\mathit{gun}_{w'}d_w \wedge \mathit{have}_w(a^*, d_w))$   
 $\omega \quad \lambda w \in p^*. p^*$   
 $\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' | \exists x(\mathit{bear}_{w'}x \wedge \mathit{near}_{w'}(x, a^*))\}$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<4, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t.} \quad \bullet \mathbf{C} = \{w' \in \mathbf{B} | \exists y(\mathit{bear}_{w'}y \wedge \mathit{come}_{w'}y)\}$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset)$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<4, w} \mathbf{B}) \cap \mathbf{C} = \emptyset)$   
 $\omega_3 \quad \lambda w \in p^*. \mathbf{H} \quad \text{s.t.} \quad \bullet \mathbf{H} = \{w' \in \mathbf{C} | \mathit{hurt}_{w'}(c_w, \mathbf{adr } a^*)\}$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<1, w} \mathbf{C}) \cap \mathbf{H} = \emptyset)$   
 $\omega_4 \quad \lambda w \in p^*. \mathbf{G} \quad \text{s.t.} \quad \bullet \mathbf{G} = \{w' | \exists z(\mathit{gun}_{w'}z \wedge \mathit{have}_w(a^*, z))\}$

(6<sup>5</sup>)  $\text{do}_{\omega, \omega4}$   
 $[\mathbf{real}(\omega, \omega_4)]i_4i_4$

iff  $\omega i_4 \subseteq \omega_4 i_4$

iff  $p^* \subseteq \mathbf{G}$

(6<sup>6</sup>) Pragmatic strengthening

$[\mathbf{real}(\omega_2, \omega_4)]$

iff  $\omega_2 i_4 \subseteq \omega_4 i_4$

iff  $\mathbf{C} \subseteq \mathbf{G}$

$i_4$   
 $u \quad \lambda w. a^*$   
 $u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(\mathit{bear}_{w'}b_w \wedge \mathit{near}_{w'}b_w)$   
 $u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{C}(\mathit{bear}_{w'}c_w \wedge \mathit{come}_{w'}c_w)$   
 $u_3 \quad \lambda w \in \mathbf{G}. d_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{G}(\mathit{gun}_{w'}d_w \wedge \mathit{have}_w(a^*, d_w))$   
 $\omega \quad \lambda w \in p^*. p^* \quad \text{s.t.} \quad \bullet p^* \subseteq \mathbf{G} \quad (6^5)$   
 $\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' | \exists x(\mathit{bear}_{w'}x \wedge \mathit{near}_{w'}(x, a^*))\}$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<4, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t.} \quad \bullet \mathbf{C} = \{w' \in \mathbf{B} | \exists y(\mathit{bear}_{w'}y \wedge \mathit{come}_{w'}y)\} \subseteq \mathbf{G} \quad (6^6)$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset)$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<4, w} \mathbf{B}) \cap \mathbf{C} = \emptyset)$   
 $\omega_3 \quad \lambda w \in p^*. \mathbf{H} \quad \text{s.t.} \quad \bullet \mathbf{H} = \{w' \in \mathbf{C} | \mathit{hurt}_{w'}(c_w, \mathbf{adr } a^*)\}$   
 $\quad \bullet \forall w \in p^*((\mathbf{min}_{<1, w} \mathbf{C}) \cap \mathbf{H} = \emptyset)$   
 $\omega_4 \quad \lambda w \in p^*. \mathbf{G} \quad \text{s.t.} \quad \bullet \mathbf{G} = \{w' | \exists z(\mathit{gun}_{w'}z \wedge \mathit{have}_w(a^*, z))\}$

(6<sup>7</sup>) If <sub>$\omega, \omega_2$</sub>  <sup>$\omega_5$</sup>  the <sub>$u_2$</sub>  bear attacked you  
 ( $\max_{\omega}^{\omega_5}([\omega_5 \subseteq \omega_2]; [bear_{\omega_5}\{u_2\}, attack_{\omega_5}\{u_2, \mathbf{adr} u\}]); [\mathbf{remote}_{<1}(\omega_2, \omega_5)]i_4i_5$ )  
 iff  $i_4[\omega: \omega_5]i_5 \wedge \text{Dom } \omega_5i_5 = p^* \wedge \forall w \in p^*(\omega_5i_5w = \{w' \in \mathbf{C} \mid attack_{w'}(c_{w'}, \mathbf{adr} a^*)\})$   
 $\wedge \forall w \in p^*((\min_{<1, w} \mathbf{C}) \cap \omega_5i_4w = \emptyset)$

$i_5$   
 $u \quad \lambda w. a^*$   
 $u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(bear_{w'} b_{w'} \wedge near_{w'} b_{w'})$   
 $u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{C}(bear_{w'} c_{w'} \wedge come_{w'} c_{w'})$   
 $u_3 \quad \lambda w \in \mathbf{G}. d_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{G}(gun_{w'} d_{w'} \wedge have_w(a^*, d_{w'}))$   
 $\omega \quad \lambda w \in p^*. p^* \quad \text{s.t.} \quad \bullet p^* \subseteq \mathbf{G}$   
 $\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \mid \exists x(bear_{w'} x \wedge near_w(x, a^*))\}$   
 $\quad \bullet \forall w \in p^*((\min_{<4, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t.} \quad \bullet \mathbf{C} = \{w' \in \mathbf{B} \mid \exists y(bear_{w'} y \wedge come_{w'} y)\} \subseteq \mathbf{G}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset)$   
 $\quad \bullet \forall w \in p^*((\min_{<4, w} \mathbf{B}) \cap \mathbf{C} = \emptyset)$   
 $\omega_3 \quad \lambda w \in p^*. \mathbf{H} \quad \text{s.t.} \quad \bullet \mathbf{H} = \{w' \in \mathbf{C} \mid hurt_{w'}(c_{w'}, \mathbf{adr} a^*)\}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{C}) \cap \mathbf{H} = \emptyset)$   
 $\omega_4 \quad \lambda w \in p^*. \mathbf{G} \quad \text{s.t.} \quad \bullet \mathbf{G} = \{w' \mid \exists z(gun_{w'} z \wedge have_w(a^*, z))\}$   
 $\omega_5 \quad \lambda w \in p^*. \mathbf{A} \quad \text{s.t.} \quad \bullet \mathbf{A} = \{w' \in \mathbf{C} \mid attack_{w'}(c_{w'}, \mathbf{adr} a^*)\}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{C}) \cap \mathbf{A} = \emptyset)$

(6<sup>8</sup>) I would <sub>$\omega, \omega_5$</sub>  <sup>$\omega_6$</sup>  use it <sub>$u_3$</sub>  to shoot him <sub>$u_2$</sub>   $uts = \text{use to shoot}$   
 ( $[\mathbf{remote}_{<1}(\omega, \omega_5)]; \max_{\omega}^{\omega_6}([\omega_6 \subseteq \omega_5]; [uts_{\omega_6}\{u, u_3, u_4\}]); [\mathbf{necc}_{<1}(\omega_5, \omega_6)]i_5i_6$ )  
 iff  $\forall w \in p^*((\min_{<1, w} p^*) \cap \mathbf{A} = \emptyset \wedge i_5[\omega: \omega_6]i_6 \wedge \text{Dom } \omega_6i_6 = p^*$   
 $\wedge \forall w \in p^*(\omega_6i_6w = \{w' \in \mathbf{A} \mid uts_{w'}(a^*, d_{w'}, c_{w'})\}) \wedge \forall w \in p^*(\min_{<1, w} \mathbf{A}) \subseteq \omega_6i_6w)$

$i_6$   
 $u \quad \lambda w. a^*$   
 $u_1 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{B}(bear_{w'} b_{w'} \wedge near_{w'} b_{w'})$   
 $u_2 \quad \lambda w \in \mathbf{C}. c_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{C}(bear_{w'} c_{w'} \wedge come_{w'} c_{w'})$   
 $u_3 \quad \lambda w \in \mathbf{G}. d_w \quad \text{s.t.} \quad \bullet \forall w' \in \mathbf{G}(gun_{w'} d_{w'} \wedge have_w(a^*, d_{w'}))$   
 $\omega \quad \lambda w \in p^*. p^* \quad \text{s.t.} \quad \bullet p^* \subseteq \mathbf{G}$   
 $\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t.} \quad \bullet \mathbf{B} = \{w' \mid \exists x(bear_{w'} x \wedge near_w(x, a^*))\}$   
 $\quad \bullet \forall w \in p^*((\min_{<4, w} p^*) \cap \mathbf{B} = \emptyset)$   
 $\omega_2 \quad \lambda w \in p^*. \mathbf{C} \quad \text{s.t.} \quad \bullet \mathbf{C} = \{w' \in \mathbf{B} \mid \exists y(bear_{w'} y \wedge come_{w'} y)\} \subseteq \mathbf{G}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{B}) \cap \mathbf{C} \neq \emptyset)$   
 $\quad \bullet \forall w \in p^*((\min_{<4, w} \mathbf{B}) \cap \mathbf{C} = \emptyset)$   
 $\omega_3 \quad \lambda w \in p^*. \mathbf{H} \quad \text{s.t.} \quad \bullet \mathbf{H} = \{w' \in \mathbf{C} \mid hurt_{w'}(c_{w'}, \mathbf{adr} a^*)\}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{C}) \cap \mathbf{H} = \emptyset)$   
 $\omega_4 \quad \lambda w \in p^*. \mathbf{G} \quad \text{s.t.} \quad \bullet \mathbf{G} = \{w' \mid \exists z(gun_{w'} z \wedge have_w(a^*, z))\}$   
 $\omega_5 \quad \lambda w \in p^*. \mathbf{A} \quad \text{s.t.} \quad \bullet \mathbf{A} = \{w' \in \mathbf{C} \mid attack_{w'}(c_{w'}, \mathbf{adr} a^*)\}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{C}) \cap \mathbf{A} = \emptyset)$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} p^*) \cap \mathbf{A} = \emptyset)$   
 $\omega_6 \quad \lambda w \in p^*. \mathbf{S} \quad \text{s.t.} \quad \bullet \mathbf{S} = \{w' \in \mathbf{A} \mid uts_{w'}(a^*, d_{w'}, c_{w'})\}$   
 $\quad \bullet \forall w \in p^*((\min_{<1, w} \mathbf{A}) \subseteq \mathbf{S})$

## APPENDIX 1: THREE VERSIONS OF KEY DEFINITIONS

Version 1: Stone 1997 (our notation)

- *Discourse referents* as concepts restricted to modal domain  $\omega$

$$\text{D1. } i[\omega: u]j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq u \rightarrow vi = vj) \quad \textit{nominal dref}$$

$$\quad \quad \quad \wedge \forall w(w \in \text{Ran } \omega j \rightarrow ujw \text{ in } w) \quad \text{p. 29}$$

$${}^1\text{D2. } i[\omega: \omega']j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq \omega' \rightarrow vi = vj) \quad \textit{modal dref}$$

$$\quad \quad \quad \wedge \text{Dom } \omega'j \subseteq \text{Ran } \omega j \quad \text{p. 29}$$

- *Lexical relations* (relativized to modal domain  $\omega$ )

$$\text{D3. } R_{\omega}\{\delta_1, \dots, \delta_n\} \quad := \quad \lambda i. \forall w(w \in \text{Ran } \omega i \rightarrow R_w(\delta_1 i w, \dots, \delta_n i w)) \quad \text{p. 28}$$

- *Modal relations* (parametrized by modal base  $M$  and weak optimality ranking  $\leq$ )

$${}^1\text{D4. "}\omega' \text{ elaborates on } \omega\text{" (cf. Brasoveanu's 2007 "}\sqsubseteq\text{"})$$

$$\omega \rightsquigarrow \omega' \quad := \quad \lambda i. \forall w \in \text{Dom } \omega i (\omega' i w \subseteq \omega i w) \quad \text{cf. p. 23, 30}$$

$${}^1\text{D5. } \mathbf{necc}_{M, \leq}(\omega, \omega') \quad := \quad \lambda i. \forall w \in \text{Dom } \omega i ((\mathbf{min}_{\leq}(\omega i w \cap M)) \subseteq \omega' i w) \quad \text{p. 31}$$

(For all  $w \in \text{Dom } \omega i$ , every  $M$ -accessible  $\leq_w$ -ideal world in  $\omega i w$  is also in  $\omega' i w$ .)

$${}^1\text{D6. } \mathbf{poss}_{M, \leq}(\omega, \omega') \quad := \quad \lambda i. \forall w \in \text{Dom } \omega i ((\mathbf{min}_{\leq}(\omega i w \cap M)) \cap \omega' i w \neq \emptyset) \quad \text{p. 31}$$

(For every  $w$ , some  $M$ -accessible  $\leq_w$ -ideal world in  $\omega i w$  is also in  $\omega' i w$ .)

- *Modal presuppositions*

"Modals are classified not only by the relationship they impose between  $\omega_r$  [MB: reference modality] and  $\omega_e$  [MB: event modality], but also by the relationship they presuppose between  $\omega_r$  and  $\omega_s$  [MB: speech modality]. These are given by further conditions: **real**( $\omega_s, \omega_r, *$ ), **vivid**( $\omega_s, \omega_r, *$ ) and **remote**( $\omega_s, \omega_r, *$ ), describe the presuppositions of factual, vivid and remote modals, respectively." p. 25

$${}^1\text{D8. } \mathbf{real}_{M, \leq}(\omega, \omega') \quad := \quad \mathbf{necc}_{M, \leq}(\omega, \omega') \quad \text{cf. p. 31}$$

$$\mathbf{vivid}_{M, \leq}(\omega, \omega') \quad := \quad \mathbf{poss}_{M, \leq}(\omega, \omega') \quad \text{cf. p. 31}$$

$$\mathbf{remote}_{M, \leq}(\omega, \omega') \quad := \quad \lambda i. \neg \mathbf{poss}_{M, \leq}(\omega, \omega')i \quad \text{(idea 1) cf. p. 31}$$

$$\quad \quad \quad := \quad \lambda i. i = i \quad \text{(idea 2)}$$

- *Conditional update* ( $\underline{\leq}$  for weak similarity ranking) and "default inference" (implicature?)

$${}^1\text{D9. } \mathbf{if}_{\underline{\leq}}(\omega, \omega', D) \quad := \quad \text{p. 30}$$

$$\lambda ij. \exists k(i[\omega: \omega']k \wedge Dkj)$$

$$\quad \wedge \forall h(\exists k(i[\omega: \omega']k \wedge Dkh)$$

$$\quad \rightarrow \forall w(w \in \text{Ran } \omega i$$

$$\quad \rightarrow \forall w'w''(w' \in \omega' jw \wedge w'' \in \omega' hw \wedge w'' \underline{\leq}_w w' \rightarrow w' \underline{\leq}_w w''))$$

$$\quad \wedge \forall w''(w'' \in \omega' hw \rightarrow \exists w'(w' \in \omega' jw \wedge w' \underline{\leq}_w w''))))$$

Version 2: Stone and Hardt 1999 (our notation)

- *Discourse referents* as concepts restricted to modal domain  $\omega$

$$\text{D1. } i[\omega: u]j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq u \rightarrow vi = vj) \quad \textit{nominal dref}$$

$$\wedge \forall w(w \in \text{Ran } \omega j \rightarrow u j w \text{ in } w) \quad \text{in (6): 1st clause}$$

$${}^1\text{D2. } i[\omega: \omega']j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq \omega' \rightarrow vi = vj) \quad \textit{modal dref}$$

$$\wedge \text{Dom } \omega' j \subseteq \text{Ran } \omega j \quad \text{in (6): 2nd clause}$$

- *Lexical relations* (relativized to modal domain  $\omega$ )

$$\text{D3. } R_{\omega}\{\delta_1, \dots, \delta_n\} := \lambda i. \forall w(w \in \text{Ran } \omega i \rightarrow R_w(\delta_1 i w, \dots, \delta_n i w)) \quad \text{in (7)}$$

- *Modal relations* (no parameters)

$${}^2\text{D6. } \mathbf{poss}(\omega, \omega') \quad := \quad \lambda i. \exists w(w \in \text{Ran } \omega i \wedge w \in \omega' i w) \quad \text{in (10)}$$

(Some world  $w$  in  $\text{Ran } \omega i$  is in  $\omega' i w$ .)

$${}^2\text{D7. } \mathbf{not}(\omega, \omega') \quad := \quad \lambda i. \neg \exists w(w \in \text{Ran } \omega i \wedge w \in \omega' i w) \quad \text{in (10)}$$

(No world  $w$  in  $\text{Ran } \omega i$  is in  $\omega' i w$ .)

- *Conditional update* ( $\underline{\leq}$  for weak similarity ranking)

$${}^1\text{D9. } \mathbf{if}_{\underline{\leq}}(\omega, \omega', D) \quad := \quad (9)$$

$$\lambda i j. \exists k(i[\omega: \omega']k \wedge Dkj)$$

$$\wedge \forall h(\exists k(i[\omega: \omega']k \wedge Dkh)$$

$$\rightarrow \forall w(w \in \text{Ran } \omega i$$

$$\rightarrow \forall w' w''(w' \in \omega' j w \wedge w'' \in \omega' h w \wedge w'' \underline{\leq}_w w' \rightarrow w' \underline{\leq}_w w''))$$

$$\wedge \forall w''(w'' \in \omega' h w \rightarrow \exists w'(w' \in \omega' j w \wedge w' \underline{\leq}_w w''))))$$

Version 3: Stone 1999 (our notation)

- *Discourse referents* as concepts restricted to modal domain  $\omega$

$$\text{D1. } i[\omega: u]j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq u \rightarrow vi = vj) \quad \textit{nominal dref} \\ \wedge \forall w(w \in \text{Ran } \omega j \rightarrow u j w \text{ in } w) \quad (28)$$

$${}^3\text{D2. } i[\omega: \omega']j \quad := \quad \forall v(\mathbf{mk}(v) \wedge v \neq \omega' \rightarrow vi = vj) \quad \textit{modal dref} \\ \wedge \text{Dom } \omega' j \subseteq \text{Dom } \omega i \quad (27)$$

- *Lexical relations* (relativized to modal domain  $\omega$ )

$$\text{D3. } R_{\omega}\{\delta_1, \dots, \delta_n\} \quad := \quad \lambda i. \forall w(w \in \text{Ran } \omega i \rightarrow R_w(\delta_1 i w, \dots, \delta_n i w)) \quad (26): \text{1st line}$$

- *Modal relations* (parametrized by modal base *relation*  $M$  and weak optimality ranking  $\leq$ )

$${}^3\text{D5. } \mathbf{necc}_{M, \leq}(\omega, \omega') \quad := \\ \lambda i. \forall w(w \in \text{Dom } \omega i \\ \rightarrow \forall w'(w' \in \omega i w \wedge w' \in M w \\ \wedge \exists w''(w'' \in \omega i w \vee w'' \in \omega' i w) \wedge w'' \in M w \wedge w'' \leq_w w' \\ \wedge \forall w'''(w''' \in \omega i w \wedge w''' \in M w \wedge w''' \leq_w w' \rightarrow w''' \in \omega' i w))))$$

(For all  $w \in \text{Dom } \omega i$ , for every  $w' \in (M w \cap \omega i w)$  there is some approximation  $w''$  to the  $\leq_w$ -ideal s.t. every  $(M w \cap \omega i w)$ -world  $w'''$  as good as or better than  $w''$  is in  $\omega' i w$ .)

This is a “close enough” theory, which as we’ve seen gets into trouble with comparatives.)

$${}^3\text{D6. } \mathbf{poss}_{M, \leq}(\omega, \omega') \quad := \\ \lambda i. \forall w(w \in \text{Dom } \omega i \\ \rightarrow \exists w'(w' \in \omega i w \wedge w' \in M w \\ \wedge \forall w''(w'' \in \omega i w \vee w'' \in \omega' i w) \wedge w'' \in M w \wedge w'' \leq_w w' \\ \rightarrow \exists w'''(w''' \in \omega i w \wedge w''' \in M w \wedge w''' \leq_w w' \rightarrow w''' \in \omega' i w))))$$

(Dual of “close enough” necessity. Hard to paraphrase and, in my view, not worth the trouble.)

- *Conditional update*

$${}^3\text{D9. } \mathbf{if}_{\leq}(\omega, \omega', D) \quad := \\ \lambda i j. \exists k(i[\omega: \omega']k \wedge D k j) \\ \wedge \forall h(\exists k(i[\omega: \omega']k \wedge D k h) \\ \rightarrow \forall w w'(w' \in \omega i w \\ \rightarrow \forall w_j w_h(w_j \in \omega' j w \wedge w_h \in \omega' h w \wedge w_h \leq_{w'} w_j \rightarrow w_j \leq_w w_h) \\ \wedge \forall w_h(w_h \in \omega' h w \rightarrow \exists w_j(w_j \in \omega' j w \wedge w_j \leq_{w'} w_h))))$$

APPENDIX 2: DYNAMIC  $Ty_2$  WITH ENCAPSULATED QUANTIFICATION ( $EDy_2$ )

## D01. Basic terms

<u>Variables</u>	<u>Constants</u>	<u>Type</u>	<u>Name of objects</u>
$i, j, k, h$		$s$	atomic (info) states ( <i>aka</i> indices)
$x, y, z$	$a^*, john, \dots$	$e$	entities
$w$	$w^*, \dots$	$w$	worlds
	$b, c, \dots$	$we$	entity-concepts
$p$	$p^*, \mathbf{A}, \mathbf{B}, \dots$	$wt$	propositions
	$M, \dots$	$wwt$	accessibility relations
	$\leq_n$	$wwwt$	ordering relations
	$u_n$	$s(we)$	$e$ -stores ( $u_0$ for the speaker)
	$\omega_n$	$s(wwt)$	$wwt$ -stores ( $\omega_0$ for the reality)

## D02. Auxiliary definitions

- $i[\omega: u]j$  for (for every dref  $\delta$  other than  $u$ ,  $\delta i = \delta j$ ) &  $\text{Dom } uj = \text{Ran } \omega i$
- $i[\omega: \omega']j$  for (for every dref  $\delta$  other than  $\omega$ ,  $\delta i = \delta j$ ) &  $\text{Dom } \omega'j = \text{Dom } \omega i$
- $\mathbf{min}_{<_w}(p)$  :=  $\{w' \mid w' \in p \wedge \neg \exists w''(w'' \in p \wedge w'' <_w w')\}$
- $\lambda w \in p. f_w$  :=  $\{\langle w, f_w \rangle : w \in p\} \cup \{\langle w, \emptyset \rangle : w \notin p\}$
- $M \subseteq M'$  :=  $\forall w w'(w' \in Mw \rightarrow w' \in M'w)$
- $\text{Dom } M$  :=  $\{w \mid Mw \neq \emptyset\}$
- $\text{Ran } M$  :=  $\{w' \mid \exists w(w' \in Mw)\}$

## D03. Defined dref's

- $John$  :=  $\lambda i. (\lambda w. john)$
- $u_0$  :=  $\lambda i. (\lambda w. a^*)$
- $\omega_0$  :=  $\lambda i. (\lambda w \in p^*. p^*)$

D1. Conditions (type  $st$ )

- $R_\omega\{\delta_1, \dots, \delta_n\}$  :=  $\lambda i. \forall w \in \text{Ran } \omega i(R_w(\delta_1 i w, \dots, \delta_n i w))$
- $\delta = \delta'$  :=  $\lambda i. \delta i = \delta' i$
- $\omega \subseteq \omega'$  :=  $\lambda i. \omega i \subseteq \omega' i$  cf. S97:  $\rightsquigarrow$
- $\mathbf{necc}_{<_n}(\omega, \omega')$  :=  $\lambda i. \forall w \in \text{Dom } \omega i((\mathbf{min}_{<_n, w} \omega i w) \subseteq \omega' i w)$  S97:31
- $\mathbf{poss}_{<_n}(\omega, \omega')$  :=  $\lambda i. \forall w \in \text{Dom } \omega i((\mathbf{min}_{<_n, w} \omega i w) \cap \omega' i w \neq \emptyset)$  S97:31
- $\mathbf{not}_{<_n}(\omega, \omega')$  :=  $\lambda i. \forall w \in \text{Dom } \omega i((\mathbf{min}_{<_n, w} \omega i w) \cap \omega' i w = \emptyset)$  S97:31
- $\mathbf{real}(\omega, \omega')$  :=  $(\omega \subseteq \omega')$  cf. S97:25
- $\mathbf{vivid}_{<_n}(\omega, \omega')$  :=  $\mathbf{poss}_{<_n}(\omega, \omega')$  cf. S97:25
- $\mathbf{remote}_{<_n}(\omega, \omega')$  :=  $\mathbf{not}_{<_n}(\omega, \omega')$  cf. S97:25
- $(C, C')$  :=  $\lambda i. Ci \wedge C' i$

D2. Updates (type  $s(st)$ )

- $[C]$  :=  $\lambda ij. i = j \wedge Cj$
- $[\omega: \delta \mid C]$  :=  $\lambda ij. i[\omega: \delta]j \wedge Cj$
- $(D_1; D_2)$  :=  $\lambda ij. \exists k(D_1 i j \wedge D_2 k j)$
- $\mathbf{max}_\omega^{\omega'}(D)$  :=  $\lambda ij. (\exists k(i[\omega: \omega']k \wedge Dk j) \wedge \forall h(\exists k([\omega: \omega']ik \wedge Dkh) \rightarrow \omega' h \subseteq \omega' j))$

## D3. Truth (relative to default state)

- $\models D$  :=  $\exists ij D ij$

## APPENDIX 3: SAMPLE DERIVATIONS

(2) My neighbours *would* kill me.[uttered while looking at high-end stereo, *c*, in electronics store](2<sup>0</sup>) Start-up
$$\begin{array}{l}
i_0 \\
u \quad \lambda w. a^* \\
u_1 \quad \lambda w. c \\
\omega \quad \lambda w \in p^*. p^*
\end{array}$$
(2<sup>1</sup>) Accommodate presup. of *would* <sub>$\omega, \omega_1$</sub>  <sup>$\omega_2$</sup>  by pragmatic inference (*pfb* = play for full blast) $\mathbf{max}_\omega^{\omega_1}([buy_{\omega_1}\{u, u_1\}, pfb_{\omega_1}\{u, u_1\}])i_0i_1$ 

(1) iff (9)

1.  $\mathbf{max}_\omega^{\omega_1}([buy_{\omega_1}\{u, u_1\}, pfb_{\omega_1}\{u, u_1\}])i_0i_1$
2.  $\exists k(i_0[\omega: \omega_1]k \wedge [buy_{\omega_1}\{u, u_1\}, pfb_{\omega_1}\{u, u_1\}]ki_1)$  D2.max  
 $\wedge \forall h(\exists k(i_0[\omega: \omega_1]k \wedge [buy_{\omega_1}\{u, u_1\}, pfb_{\omega_1}\{u, u_1\}]kh) \rightarrow \omega_1h \subseteq \omega_1i_1)$
3.  $\exists k(i_0[\omega: \omega_1]k \wedge k = i_1 \wedge \forall w \in \text{Ran } \omega_1i_1(buy_w(ui_1w, u_1i_1w) \wedge pfb_w(ui_1w, u_1i_1w)))$  D2.C, D1  
 $\wedge \forall h(\exists k(i_0[\omega: \omega_1]k \wedge [buy_{\omega_1}\{u, u_1\}, pfb_{\omega_1}\{u, u_1\}]kh) \rightarrow \omega_1h \subseteq \omega_1i_1)$
4.  $i_0[\omega: \omega_1]i_1 \wedge \forall w \in \text{Ran } \omega_1i_1(buy_w(ui_1w, u_1i_1w) \wedge pfb_w(ui_1w, u_1i_1w))$  simplify  
 $\wedge \forall h(i_0[\omega: \omega_1]i_1 \wedge \forall w \in \text{Ran } \omega_1h(buy_w(uhw, u_1hw) \wedge pfb_w(uhw, u_1hw)))$  D2, 1, smp  
 $\rightarrow \omega_1h \subseteq \omega_1i_1)$
5.  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = \text{Dom } \omega_1i_0 \wedge \forall w \in \text{Ran } \omega_1i_1(buy_w(a^*, c) \wedge pfb_w(a^*, c))$  D02.[ $\omega: \omega'$ ]  
 $\wedge \forall h(i_0[\omega: \omega_1]h \wedge \text{Dom } \omega_1h = \text{Dom } \omega_1i_0 \wedge \forall w \in \text{Ran } \omega_1h(buy_w(a^*, c) \wedge pfb_w(a^*, c)))$   
 $\rightarrow \omega_1h \subseteq \omega_1i_1)$
6.  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in \text{Ran } \omega_1i_1(buy_w(a^*, c) \wedge pfb_w(a^*, c))$  D02, df.  $i_0$   
 $\wedge \forall h(i_0[\omega: \omega_1]h \wedge \text{Dom } \omega_1h = p^* \wedge \forall w \in \text{Ran } \omega_1h(buy_w(a^*, c) \wedge pfb_w(a^*, c)))$   
 $\rightarrow \forall ww'(w' \in \omega_1hw \rightarrow w' \in \omega_1i_1w)$  D03. $\subseteq$
7.  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = p^* \wedge \forall w \in \text{Ran } \omega_1i_1(buy_w(a^*, c) \wedge pfb_w(a^*, c))$   
 $\wedge \forall h(i_0[\omega: \omega_1]h \wedge \text{Dom } \omega_1h = p^* \wedge \forall w \in \text{Ran } \omega_1h(buy_w(a^*, c) \wedge pfb_w(a^*, c)))$   
 $\rightarrow \forall w \in p^*(\omega_1hw \subseteq \omega_1i_1w)$  simplify
8.  $i_0[\omega: \omega_1]i_1 \wedge \text{Dom } \omega_1i_1 = p^*$   
 $\wedge \forall w \in p^*(\omega_1i_1w = \{w' \mid buy_w(a^*, c) \wedge pfb_w(a^*, c)\})$  simplify

$$\begin{array}{l}
i_1 \\
u \quad \lambda w. a^* \\
u_1 \quad \lambda w. c \\
\omega \quad \lambda w \in p^*. p^* \\
\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t. } \bullet \mathbf{B} := \{w' \mid buy_w(a^*, c) \wedge pfb_w(a^*, c)\}
\end{array}$$

(2<sup>2</sup>) My neighbours $[\omega_1: u_2 \mid \text{neighbours.of}_{\omega_1}\{u_2, u\}]i_1i_2$ 

(1) iff (4)

1.  $[\omega_1: u_2 \mid \text{neighbours.of}_{\omega_1}\{u_2, u\}]i_1i_2$ 2.  $i_1[\omega_1: u_2]i_2 \wedge (\text{neighbours.of}_{\omega_1}\{u_2, u\})i_2$ 

D2

3.  $i_1[\omega_1: u_2]i_2 \wedge \text{Dom } u_2i_2 = \text{Ran } \omega_1i_1$   
 $\wedge \forall w \in \text{Ran } \omega_1i_2(\text{neighbours.of}_w(u_2i_2w, ui_2w))$ D02.[ $\omega: u$ ]

D1.R

4.  $i_1[\omega_1: u_2]i_2 \wedge \text{Dom } u_2i_2 = \mathbf{B} \wedge \forall w \in \mathbf{B}(\text{neighbours.of}_w(u_2i_2w, a^*))$ D02, df.  $i_1$  $i_2$   
 $u \quad \lambda w. a^*$  $u_1 \quad \lambda w. c$  $u_2 \quad \lambda w \in \mathbf{B}. b_w \quad \text{s.t. } \bullet \forall w \in \mathbf{B}(\text{neighbours.of}_w(b_w, a^*))$  $\omega \quad \lambda w \in p^*. p^*$  $\omega_1 \quad \lambda w \in p^*. \mathbf{B} \quad \text{s.t. } \bullet \mathbf{B} := \{w' \mid \text{buy}_w(a^*, c) \wedge \text{pfb}_w(a^*, c)\}$ 

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### Strong readings in UTy<sub>3</sub>

#### 0. OUTLINE OF ANALYSIS

(!) Start-up update a la Bittner, 2007  
 $[ul\ u = AI]; [rl\ speaking_r\langle u \rangle]; [p]; [p \sim r];$   $u := u_0, r := r_0, p := p_0$

(V) Strong readings as maximizing anaphora

(1) 0. presupposition accommodation  
 $[u_1\mid u_1 = country.of_r\langle u \rangle]; [p_1\mid law.of_r\langle p_1, u_1 \rangle];$

i. [According to<sup>S<sub>1</sub>, p<sub>2</sub></sup> the law]  
 $[S_1]; [S_1 \sim_r p_1]; [p_2\mid p_2 = \cup S_1];$

Output if the GB-law in  $w_1$  is  $\pi_0$ , the laws in  $w_2$  are  $\pi_0, \pi_1$ , and  $\pi_2$ , and in  $w_3, \pi_0$  and  $\pi_1$ , where  $\pi_0 := \{w_0\}$ ,  $\pi_1 := \{w_0, w_1\}$  and  $\pi_2 := \{w_0, w_1, w_2\}$ .

$J_1$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	
$i_1$	<i>al</i>	$w_1$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0\}$	$\pi_0$	$J_{1, r=w_1}$
$i_2$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$J_{1, r=w_2}$
$i_3$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	
$i_4$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	
$i_5$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$J_{1, r=w_3}$
$i_6$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	

ii. if<sub>p<sub>2</sub></sub><sup>r<sub>2</sub></sup> [**a**<sup>u<sub>2</sub>, A<sub>2</sub></sup> man commits a murder]<sup>p<sub>3</sub></sup>  
 $[r_2\mid r_2 \in p_2]; [u_2\mid man_{r_2}\langle u_2 \rangle, mrd_{r_2}\langle u_2 \rangle]; [A_2]; [A_2 \sim_{r_2} u_2]; [p_3]; [p_3 \sim_r r_2]$

Output if no man commits murder in  $w_0$ , men  $a$  and  $b$  commit murders in  $w_1$ , men  $a$  and  $c$  in  $w_2$  (and men  $b$  and  $c$  in  $w_3$ ):

$J_2$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	$p_3$	
$i_{21}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$J_{2, r=w_2}$
$i_{22}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	
$i_{31}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	
$i_{32}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	
$i_{41}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	
$i_{42}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	
$i_{23}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{24}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{33}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{34}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{43}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{44}$	<i>al</i>	$w_2$	...	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$	
$i_{51}$	<i>al</i>	$w_3$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	$J_{2, r=w_3}$
$i_{52}$	<i>al</i>	$w_3$	...	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$	
$i_{61}$	<i>al</i>	$w_3$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	
$i_{62}$	<i>al</i>	$w_3$	...	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$	

iii. [**he**<sub>*u*2</sub> should<sub>*S*1, *p*2, *p*3</sub><sup>*p*4</sup> go to prison].

[**min**<sub>*S*1</sub>⟨*p*2⟩ ∅ *p*3]; [*p*4 | *p*4 = **min**<sub>*S*1</sub>⟨*p*3⟩]; [*gtp*<sub>*r*2</sub>⟨*u*2⟩]; [**A**<sub>2</sub> ~<sub>*r*2</sub> **u**<sub>2</sub>]; [*p*4 ~<sub>*r*</sub> *r*2]; [*p*]; [*p* ~ *r*]

Output if *a* and *b* go to prison (gtp) in *w*<sub>1</sub>; only *a* gtp in *w*<sub>2</sub>; (and nobody gtp in *w*<sub>3</sub>).

<i>J</i> <sub>3</sub>	<i>u</i>	<i>r</i>	<i>p</i>	<i>u</i> <sub>1</sub>	<i>p</i> <sub>1</sub>	<i>S</i> <sub>1</sub>	<i>p</i> <sub>2</sub>	<i>r</i> <sub>2</sub>	<i>u</i> <sub>2</sub>	<i>A</i> <sub>2</sub>	<i>p</i> <sub>3</sub>	<i>p</i> <sub>4</sub>
<i>i</i> <sub>21</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>22</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>31</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>32</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>41</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>42</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0$ , $\pi_1$ , $\pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>51</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0$ , $\pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>52</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0$ , $\pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>61</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0$ , $\pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>62</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0$ , $\pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }

## 1. START-UP UPDATE

Default info state (dummy)

	<i>u</i>	...	<i>r</i>	...	<i>p</i>	...	<i>S</i> <sub>1</sub>	...	<i>A</i> <sub>2</sub>
<i>i</i> <sub>@</sub>	@ <sub><i>e</i></sub>		@ <sub><i>w</i></sub>		{}		{}		{}

$$\begin{aligned} (-^2) \quad & [\mathbf{u} | \mathbf{u} = \mathbf{A}I] \\ & = \lambda I \lambda j. \exists i \in I(i[u]j \wedge uj = al) \end{aligned}$$

<i>J</i> <sub>-2</sub>	<i>u</i>	...
<i>i</i> <sub>-2</sub>	<i>al</i>	

$$\begin{aligned} (-^1) \quad & [\mathbf{r} | \mathbf{speaking}_r \langle \mathbf{u} \rangle] \\ & = \lambda I \lambda j. \exists i \in I(i[r]j \wedge \mathbf{speaking}_{r_j}(uj)) \end{aligned}$$

Output if *A*<sub>1</sub> is speaking in *w*<sub>0</sub>, *w*<sub>1</sub>, *w*<sub>2</sub> and *w*<sub>4</sub>:

<i>J</i> <sub>-1</sub>	<i>u</i>	<i>r</i>	...
<i>i</i> <sub>-11</sub>	<i>al</i>	<i>w</i> <sub>1</sub>	
<i>i</i> <sub>-12</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	
<i>i</i> <sub>-13</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	

$$\begin{aligned} (0) \quad & [\mathbf{p}]; [\mathbf{p} \sim \mathbf{r}] \\ & = \lambda I \lambda j. \exists i \in I(i[p]j \wedge @_{\mathbf{w}} \notin pj \wedge pj = rI) \end{aligned}$$

<i>J</i> <sub>0</sub>	<i>u</i>	<i>r</i>	<i>p</i>	...
<i>i</i> <sub>1</sub>	<i>al</i>	<i>w</i> <sub>1</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	
<i>i</i> <sub>2</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	
<i>i</i> <sub>3</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	

## 2. STRONG READING AS MAXIMIZING ANAPHORA

(1<sup>0</sup>) **presupposition accommodation** (*the law presupposes a salient region with laws*)  
 $[u_1 | \text{live.in}_r \langle u, u_1 \rangle]; [p_1 | \text{law.of}_r \langle p_1, u_1 \rangle]$

Output  $K_{11} := [u_1 | \text{live.in}_r \langle u, u_1 \rangle] J_0$

if the speaker (Al) lives in Great Britain (GB):

$K_{11}$	$u$	$r$	$p$	$u_1$
$k_1$	$al$	$w_1$	$\{w_1, w_2, w_3\}$	$gb$
$k_2$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$
$k_3$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$

Output  $K_{12} := [p_1 | \text{law.of}_r \langle p_1, u_1 \rangle] K_{11}$

$= \lambda j. \exists i \in K_{11} (i[p_1]j \wedge \text{law.of}_{r_j} \langle p_1 j, u_1 j \rangle)$

D2, D1.R, D02

if the laws of GB are the propositions  $\pi_0$  in  $w_1$ ,  $\pi_0$ ,  $\pi_1$ , and  $\pi_2$  in  $w_2$ , and  $\pi_0$  and  $\pi_1$  in  $w_3$ :

$K_{12}$	$u$	$r$	$p$	$u_1$	$p_1$
$k_{10}$	$al$	$w_1$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$
$k_{20}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$
$k_{21}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$
$k_{22}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$
$k_{30}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$
$k_{31}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$

(1<sup>1</sup>) **According to**<sup>S1, p2</sup>

$[S_1]; [S_1 \sim_r p_1]; [p_2 | p_2 = \cup S_1]$

Output  $I_{11} := ([S_1]; [S_1 \sim_r p_1]) K_{12}$

$= \lambda j. \exists i \in K_{12} (i[S_1]j \wedge S_1 j = p_1 K_{12, r=r_j})$

D2, D3.~<sub>r</sub>

$I_{11}$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$
$i_1$	$al$	$w_1$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0\}$
$i_2$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$
$i_3$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$
$i_4$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$
$i_5$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$
$i_5$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$

Output  $I_{12} := [p_2 | p_2 = \cup S_1] I_{11}$

$= \lambda j. \exists i \in I_{11} (i[p_2]j \wedge p_2 j = \cup(S_1 j))$

D2, D05.US

if  $\pi_0 := \{w_0\}$ ,  $\pi_1 := \{w_0, w_1\}$  and  $\pi_2 := \{w_0, w_1, w_2\}$ :

$I_{12}$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$
$j_1$	$al$	$w_1$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0\}$	$\pi_0$
$j_2$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_3$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_4$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_5$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$
$j_5$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$

(1<sup>2</sup>) **the law**

$$[S_1 \sim_r p_1]; [law.of_r \langle p_1, u_1 \rangle]$$

$$\begin{aligned} \text{Output } I_{13} &:= ([S_1 \sim_r p_1]; [law.of_r \langle p_1, u_1 \rangle])I_{12} \\ &= I_{12} \end{aligned}$$

D2, D3. $\sim_r$ , D1.R

$I_{13}$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$
$j_1$	<i>al</i>	$w_1$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0\}$	$\pi_0$
$j_2$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_3$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_4$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$
$j_5$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$
$j_5$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$

$$(2^1) \quad \mathbf{if}_{p_2}^{r_2} [ \\ [r_2 | r_2 \in p_2]$$

$$\begin{aligned} \text{Output } H_1 &:= [r_2 | r_2 \in p_2]I_{13} \\ &= \lambda j. \exists i \in I_{13} (i[r_2]j \wedge r_2 j \in p_2 j) \end{aligned}$$

D2, D1. $\in$ given  $\pi_0 := \{w_0\}$ ,  $\pi_1 := \{w_0, w_1\}$  and  $\pi_2 := \{w_0, w_1, w_2\}$ :

$H_1$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$
$h_1$	<i>al</i>	$w_1$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0\}$	$\pi_0$	$w_0$
$h_2$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$
$h_3$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$
$h_4$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$
$h_5$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$
$h_6$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$
$h_7$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$
$h_8$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$
$h_9$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$
$h_{10}$	<i>al</i>	$w_2$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$
$h_{11}$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_0$
$h_{12}$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$
$h_{13}$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_0$
$h_{14}$	<i>al</i>	$w_3$	$\{w_1, w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$

(2<sup>2</sup>) **[a<sup>u2</sup> man**  
**[u<sub>2</sub>| man<sub>r<sub>2</sub></sub>(u<sub>2</sub>)]**

Output  $H_2 := [u_2 | \text{man}_{r_2}(u_2)]H_1$   
 $= \lambda j. \exists i \in H_1(i[u_2]j \wedge \text{man}_{r_2j} u_2j)$

D2, D1.R

if the men are  $a$  in  $w_0$ ,  $a$  and  $b$  and  $w_1$ , and  $a$ ,  $b$ , and  $c$  in  $w_2$  (and  $w_3$ ):

$H_2$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$
$i_1$	$al$	$w_1$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0\}$	$\pi_0$	$w_0$	$a$
$i_2$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$	$a$
$i_3$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{3'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_4$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{4'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$b$
$i_{4''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_5$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$	$a$
$i_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{6'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{7'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$b$
$i_{7''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_8$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_0$	$a$
$i_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{9'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{10'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$b$
$i_{10''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_{11}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_0$	$a$
$i_{12}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{12'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$
$i_{13}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_0$	$a$
$i_{14}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{14'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$

(2<sup>3</sup>) **commits a murder**<sup>p3</sup>  
 $[mrd_{r_2}\langle u_2 \rangle]$

Output  $H_3 := [mrd_{r_2}\langle u_2 \rangle]H_2$

$$= \lambda j. j \in H_2 \wedge mrd_{r_2j} u_2j$$

D2, D1.R

if no men commits murder in  $w_0$ ,  $a$  and  $b$  commit murders and  $w_1$ , and  $a$  and  $c$  in  $w_2$   
 (and  $b$  and  $c$  in  $w_3$ )

$H_3$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$
$i_3$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{3'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_4$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{4''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{6'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{7''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{9'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{10''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_{12}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{12'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$
$i_{14}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{14'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$

i.e., rearranging the rows to highlight the  $r_2$ -substates:

$H_3$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$
$i_{12}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{12'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$
$i_{14}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$
$i_{14'}$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$
$i_3$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{3'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{6'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$
$i_{9'}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$
$i_4$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{4''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{7''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$
$i_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$
$i_{10''}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$

$H_{3, r_2 = w_1}$

$H_{3, r_2 = w_2}$

(2<sup>4</sup>)  $\int_{r, u^3}^{A_3}$  scope of  $\mathbf{a}^{u^3}$   
 $[A_2]; [A_2 \sim_{r_2} u_2]$

Output  $H_4 := ([A_2]; [A_2 \sim_{r_2} u_2])H_3$

$$= \lambda j. \exists i \in H_3(i[A_2]j \wedge A_2j = u_2 H_{3, r_2=r_2j})$$

D2, D3.  $\sim_r$

$H_4$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	
$j_1$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$H_{4, w=w_3}$
$j_2$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	
$j_3$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	
$j_4$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	
$j_5$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$H_{4, w=w_2}$
$j_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	
$j_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	
$j_8$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	
$j_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	
$j_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	
$j_{11}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	
$j_{12}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	
$j_{13}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	
$j_{14}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	
$j_{15}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	
$j_{16}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	

(2<sup>5</sup>)  $\int_{r, r_2}^{p^3}$  scope of  $\mathbf{if}_{p_1}^{p^2}$   
 $[p_3]; [p_3 \sim_r r_2]$

Output  $H_5 := ([p_3]; [p_3 \sim_r r_2])H_4$

$$= \lambda j. \exists i \in H_4(i[p_3]j \wedge p_3j = r_2 H_{4, r=rj})$$

D2, D3.  $\sim_r$

$H_5$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	$p_3$
$k_1$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$
$k_2$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$
$k_3$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$
$k_4$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$
$k_5$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$
$k_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$
$k_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$
$k_8$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$
$k_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$
$k_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$
$k_{11}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$
$k_{12}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$
$k_{13}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$
$k_{14}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$
$k_{15}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$
$k_{16}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$c$	$\{a, c\}$	$\{w_1, w_2\}$

(3<sup>1</sup>) **he**<sub>*u*<sub>2</sub>, *A*<sub>2</sub></sub> [ *male*<sub>*r*<sub>2</sub></sub>(*u*<sub>2</sub>); [*A*<sub>2</sub> ~<sub>*r*<sub>2</sub></sub> *u*<sub>2</sub>]; maximizing **he**

Output  $J_1 := ([male_{r_2}(u_2)]; [A_2 \sim_{r_2} u_2])H_5$   
 $= \lambda j. j \in H_5(male_{r_2j} u_2j \wedge A_2j = u_2H_{5, r_2=r_2j})$  D2, D1.R, D02, D3.~<sub>r</sub>  
 $= H_5$

(3<sup>2</sup>) [**should**<sub>*S*<sub>1</sub>, *p*<sub>2</sub>, *p*<sub>3</sub></sub><sup>*p*<sub>4</sub></sup>  
**min**<sub>*S*<sub>1</sub></sub>(*p*<sub>2</sub>) ∅ *p*<sub>3</sub>]; [*p*<sub>4</sub> | *p*<sub>4</sub> = **min**<sub>*S*<sub>1</sub></sub>(*p*<sub>3</sub>)]

Output  $J_{21} := [\mathbf{min}_{S_1}(p_2) \emptyset p_3]J_1$  remoteness-presup.  
 $= \lambda j. j \in J_1 \wedge (\mathbf{min}_{S_{1j}} p_2j) \cap p_3j = \emptyset$  D2, D1.∅  
 $= J_1$

for, given  $\pi_0 := \{w_0\}$ ,  $\pi_1 := \{w_0, w_1\}$ , and  $\pi_2 := \{w_0, w_1, w_2\}$ , we have:

- (**min**<sub>{ $\pi_0, \pi_1$ }</sub>  $\pi_1$ ) ∩ {*w*<sub>1</sub>}  
 $= \{w_0\} \cap \{w_1\}$  D03  
 $= \emptyset$
- (**min**<sub>{ $\pi_0, \pi_1, \pi_2$ }</sub>  $\pi_2$ ) ∩ {*w*<sub>1</sub>, *w*<sub>2</sub>}  
 $= \{w_0\} \cap \{w_1, w_2\}$  D03  
 $= \emptyset$

The *p*<sub>3</sub>-antecedent modality is remote from the *S*<sub>1</sub>-ideal of the *p*<sub>2</sub>-modal base, i.e. at each index, the *p*<sub>3</sub>-modality is disjoint from the *S*<sub>1</sub>-ideal of *p*<sub>2</sub>—{*w*<sub>0</sub>}, where nobody commits any murder.

Output  $J_{22} := [p_4 | p_4 = \mathbf{min}_{S_1}(p_3)]J_{21}$  necessity-update  
 $= \lambda j. \exists i \in J_{21}(i[p_4]j \wedge p_4j = \mathbf{min}_{S_{1j}} p_3j)$  D2, D1.∅

given  $\pi_0 := \{w_0\}$ ,  $\pi_1 := \{w_0, w_1\}$ , and  $\pi_2 := \{w_0, w_1, w_2\}$ :

<i>J</i> <sub>22</sub>	<i>u</i>	<i>r</i>	<i>p</i>	<i>u</i> <sub>1</sub>	<i>p</i> <sub>1</sub>	<i>S</i> <sub>1</sub>	<i>p</i> <sub>2</sub>	<i>r</i> <sub>2</sub>	<i>u</i> <sub>2</sub>	<i>A</i> <sub>2</sub>	<i>p</i> <sub>3</sub>	<i>p</i> <sub>4</sub>
<i>i</i> <sub>1</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>2</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>3</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>4</sub>	<i>al</i>	<i>w</i> <sub>3</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1$ }	$\pi_1$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>5</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>6</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>7</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>8</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>9</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>a</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>10</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>1</sub>	<i>b</i>	{ <i>a</i> , <i>b</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>11</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>a</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>12</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_0$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>c</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>13</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>a</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>14</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_1$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>c</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>15</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>a</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }
<i>i</i> <sub>16</sub>	<i>al</i>	<i>w</i> <sub>2</sub>	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> , <i>w</i> <sub>3</sub> }	<i>gb</i>	$\pi_2$	{ $\pi_0, \pi_1, \pi_2$ }	$\pi_2$	<i>w</i> <sub>2</sub>	<i>c</i>	{ <i>a</i> , <i>c</i> }	{ <i>w</i> <sub>1</sub> , <i>w</i> <sub>2</sub> }	{ <i>w</i> <sub>1</sub> }

(3<sup>3</sup>) **go to prison**  
 $[gtp_{r_2}\langle u_2 \rangle]$

Output  $J_3 := [gtp_{r_2}\langle u_2 \rangle]J_{22}$  necessity-update  
 $= \lambda j. j \in J_{22} \wedge gtp_{r_2} u_2 j$  D2, D1.R

if all murderers go to prison (gtp) in  $w_1$ , only  $a$  does in  $w_2$  (and no murderer does in  $w_3$ ):

$J_3$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	$p_3$	$p_4$	
$i_1$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	$J_{3, r_2 = w_1}$
$i_2$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	
$i_3$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	
$i_4$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	
$i_5$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_8$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_{10}$	<del><math>al</math></del>	<del><math>w_2</math></del>	<del><math>\{w_1, w_2, w_3\}</math></del>	<del><math>gb</math></del>	<del><math>\pi_2</math></del>	<del><math>\{\pi_0, \pi_1, \pi_2\}</math></del>	<del><math>\pi_2</math></del>	<del><math>w_1</math></del>	<del><math>b</math></del>	<del><math>\{a, b\}</math></del>	<del><math>\{w_1, w_2\}</math></del>	<del><math>\{w_1\}</math></del>	
$i_{11}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	$\{w_1\}$	$J_{3, r_2 = w_2}$
$i_{13}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_{15}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_2$	$a$	$\{a, c\}$	$\{w_1, w_2\}$	$\{w_1\}$	

(3<sup>4</sup>)  $I_{A_2}$  scope of maximizing **he**  
 $[A_2 \sim_{r_2} u_2]$

Output  $J_4 := [A_2 \sim_{r_2} u_2]J_3$  D2, D3.~<sub>r</sub>  
 $= \lambda j. j \in J_3 \wedge A_2 j = u_2 j_{3, r_2 = r_2 j}$

$J_4$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	$p_3$	$p_4$	
$i_1$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	$J_{4, r = w_3}$
$i_2$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$b$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	
$i_3$	$al$	$w_3$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	$a$	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$	
$i_4$	<del><math>al</math></del>	<del><math>w_3</math></del>	<del><math>\{w_1, w_2, w_3\}</math></del>	<del><math>gb</math></del>	<del><math>\pi_1</math></del>	<del><math>\{\pi_0, \pi_1\}</math></del>	<del><math>\pi_1</math></del>	<del><math>w_1</math></del>	<del><math>b</math></del>	<del><math>\{a, b\}</math></del>	<del><math>\{w_1\}</math></del>	<del><math>\{w_1\}</math></del>	
$i_5$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	$J_{4, r = w_2}$
$i_6$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_7$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_8$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_9$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$a$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	
$i_{10}$	$al$	$w_2$	$\{w_1, w_2, w_3\}$	$gb$	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	$b$	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$	

$$(3^5) \quad \mathbb{I}_{p_4} \\ [p_4 \sim_r r_2]$$

scope of **[should]**

$$\begin{aligned} \text{Output } J_5 &:= [p_4 \sim_r r_2]J_4 \\ &= \lambda j. j \in J_4 \wedge p_4 j = r_2 J_{4, r=rj} \\ &= J_4 \end{aligned}$$

D2, D3. $\sim_r$ 

$$(3^6) \quad \mathbb{I}^p \\ [p]; [p \sim r]$$

common ground update

$$\begin{aligned} \text{Output } J_6 &:= ([p]; [p \sim r])J_5 \\ &= \lambda j. \exists i \in J_5 (i[p]j \wedge pj = rJ_5) \end{aligned}$$

D2, D3. $\sim$ 

$J_6$	$u$	$r$	$p$	$u_1$	$p_1$	$S_1$	$p_2$	$r_2$	$u_2$	$A_2$	$p_3$	$p_4$
$j_1$	<i>al</i>	$w_3$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	<i>a</i>	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$
$j_2$	<i>al</i>	$w_3$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	<i>b</i>	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$
$j_3$	<i>al</i>	$w_3$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	<i>a</i>	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$
$j_4$	<i>al</i>	$w_3$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1\}$	$\pi_1$	$w_1$	<i>b</i>	$\{a, b\}$	$\{w_1\}$	$\{w_1\}$
$j_5$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>a</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$
$j_6$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_0$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>b</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$
$j_7$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>a</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$
$j_8$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_1$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>b</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$
$j_9$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>a</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$
$j_{10}$	<i>al</i>	$w_2$	$\{w_2, w_3\}$	<i>gb</i>	$\pi_2$	$\{\pi_0, \pi_1, \pi_2\}$	$\pi_2$	$w_1$	<i>b</i>	$\{a, b\}$	$\{w_1, w_2\}$	$\{w_1\}$

APPENDIX: UPDATE SEMANTICS + Ty<sub>3</sub> (UTy<sub>3</sub>)D01. Basic terms of UTy<sub>3</sub>

<u>Variables</u>	<u>Constants</u>	<u>Type</u>	<u>Name of objects</u>
$i, j, k, h$	$i_@$	$s$	atomic (info) states ( <i>aka</i> indices)
$x, y, z$	$@_e, john, \dots$	$e$	entities
$w$	$@_w, w_0, \dots$	$w$	worlds
$I, J, K, H$	$\{i_@\}$	$st$	(plural) (info) states
$\pi$	$\{w_0\}, \dots$	$wt$	propositions
$\Sigma$	$\{\{w_0\}\}, \dots$	$(wt)t$	sets of propositions
	<b>fut</b>	$w(wt)$	(possible) futures
	<b>adr</b>	$w(ee)$	
	$u_n$	$se$	$e$ -stores ( $u_0$ for the speaker)
	$A_n$	$s(et)$	$et$ -stores
	$r_n$	$sw$	$w$ -stores ( $r_0$ for the reality)
	$p_n$	$s(wt)$	$wt$ -stores ( $p_0$ for the common ground)
	$S_n$	$s((wt)t)$	$(wt)t$ -stores

## D02. Dummies

- $@_w \neq w_n \wedge @_w \notin \alpha_{w(wt)} w_n$   
 $\wedge @_e \neq c_e \wedge @_e \neq \beta_{ee} c_e \wedge @_e \neq \gamma_{w(ee)} w_n c_e$
- $R_w(x_1, \dots, x_n) \rightarrow w \neq @_w \wedge @_e \notin \{x_1, \dots, x_n\}$  for any basic relation  $R_{e\dots ew}$   
 $R_w(x, \pi) \rightarrow w \neq @_w \wedge x \neq @_e \wedge \pi \neq \{\} \wedge @_w \notin \pi$  for any basic relation  $R_{(wt)ew}$
- $\delta i_@ = @_a \wedge \Delta i_@ = \{\}$  for  $\delta_{sa}, \Delta_{s(at)}, a \in \{e, w, wt\}$

## D03. Order &amp; minimal elements

- $w \leq_\Sigma w' := \{\pi / \pi \in \Sigma \wedge w' \in \pi\} \subseteq \{\pi / \pi \in \Sigma \wedge w \in \pi\}$
- $w <_\Sigma w' := w \leq_\Sigma w' \wedge \neg w' \leq_\Sigma w$
- $\min_\Sigma(\pi) := \{w \mid w \in \pi \wedge \neg \exists w'(w' \in \pi \wedge w' <_\Sigma w)\}$

D04. Defined  $sa$ -dref's ( $a \in \{e, w\}$ )

- $John := \lambda i_s. john_e$
- $\alpha_r \langle \delta \rangle := \lambda i_s. \alpha_{ri} \delta i$  for  $\alpha_{wba}, \delta_{sb}$
- $\beta \langle \delta \rangle := \lambda i_s. \beta \delta i$  for  $\beta_{ba}, \delta_{sb}$

D05. Defined  $s(wt)$ -dref's

- $\min_s \langle p \rangle := \lambda i_s. \min_{Si} \langle pi \rangle$
- $\mathbf{dox}_s \langle u \rangle := \lambda i_s. \cap \{\pi \mid believe_{ri}(ui, \pi)\} := \lambda i_s. \{w \mid \forall \pi (believe_{ri}(ui, \pi) \rightarrow w \in \pi)\}$
- $(p - p') := \lambda i_s. (pi - p'i) := \lambda i_s. \{w \mid w \in pi \wedge w \notin p'i\}$
- $\cup S := \lambda i_s. \cup(Si) := \lambda i_s. \{w \mid \exists \pi (\pi \in Si \wedge w \in \pi)\}$

## D06. Global values, substates

- $\delta I := \{\delta i : i_s \in I_s\}$  for  $\delta_{sa}, a \in \{e, w, et, wt, (wt)t\}$
- $I_{\delta=d} := \{i_s \in I \mid \delta i = d\}$  for  $\delta_{sa}, d_a, a \in \{e, w, et, wt, (wt)t\}$

D1. Conditions (type  $st$ )

- $PL\langle A \rangle$  :=  $\lambda i_s. |A| > 1$
- $SG\langle A \rangle$  :=  $\lambda i_s. |A| = 1$
- $\delta = \delta'$  :=  $\lambda i_s. \delta i = \delta' i$  for  $\delta_{sa}, \delta'_{sa}, a \in \{e, w, et, wt, (wt)t\}$
- $\delta \in \Delta$  :=  $\lambda i_s. \delta i \in \Delta i$  for  $\delta_{sa}, \Delta_{s(st)}, a \in \{e, w, wt\}$
- $SM\langle \Delta, \Delta' \rangle$  :=  $\lambda i_s. \emptyset \subset \Delta' i \subseteq \Delta i$
- $\Delta \emptyset \Delta'$  :=  $\lambda i_s. \Delta' i \cap \Delta i = \emptyset$
- $\Delta \circ \Delta'$  :=  $\lambda i_s. \Delta' i \cap \Delta i \neq \emptyset$
- $R_r\langle \delta_1, \dots, \delta_n \rangle$  :=  $\lambda i_s. R_{r_i}(\delta_1 i, \dots, \delta_n i)$  for  $R_{a_1 \dots a_1 w}, \delta_{1, sa_1}, \dots, \delta_{n, san}$
- $R_r\langle u, r' \rangle$  :=  $\lambda i_s. \exists \pi (R_{r_i}(u i, \pi) \wedge r' i \in \pi)$  for  $R_{(wt)ew}$

D2. Local updates (type  $(st)(st)$ )

- $[\delta]$  :=  $\lambda I_{st} \lambda j_s. \exists i_s \in I(i[\delta]j \wedge @_a \neq \delta j)$  for  $\delta_{sa}, a \in \{e, w\}$
- $[\Delta]$  :=  $\lambda I_{st} \lambda j_s. \exists i_s \in I(i[\Delta]j \wedge @_a \notin \Delta j)$  for  $\Delta_{sa}, a \in \{et, wt, (wt)t\}$
- $[C_1, \dots, C_n]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge C_1 j \wedge \dots \wedge C_n j$
- $(D_1; D_2)$  :=  $\lambda I_{st}. D_2(D_1 I)$
- $(D_1 \vee D_2)$  :=  $\lambda I_{st}. (D_1 I \cup D_2 I)$
- $[\delta | C_1, \dots, C_n]$  :=  $[\delta]; [C_1, \dots, C_n]$  for  $\delta_{sa}, a \in \{e, w, et, wt, (wt)t\}$

D3. Non-local updates (type  $(st)(st)$ )

- $[\Delta \sim \delta]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge \Delta j = \delta I$  for  $\Delta_{s(at)}, \delta_{sa}, a \in \{e, w, wt\}$
- $[\Delta \sim_r \delta]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge \Delta j = \delta I_{r=rj}$  for  $\Delta_{s(at)}, \delta_{sa}, a \in \{e, w, wt\}$
- $[S \sim_r r']$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge S j = \{r' I_{r=rj}\}$
- $[SG\{\delta\}]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge |\delta I| = 1$
- $[PL_r\{\delta\}]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge |\delta I_{r=rj}| > 1$
- $[SG_r\{\delta\}]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge |\delta I_{r=rj}| = 1$
- $[_u(SG_r\{\delta\})]$  :=  $\lambda I_{st} \lambda j_s. j \in I \wedge @ \notin u I_{r=rj} \wedge \forall x_e \in u I_{r=rj} (|\delta I_{r=rj, u=x}| = 1)$

## D4. Truth (relative to default state)

- $\models D$  :=  $\exists j_s (j \in D\{i_{@}\})$