

ENGLISH-BASED THEORY:
Kalaallisut as English minus tense (Shaer 2003)

1 TEMPORAL ANAPHORA IN ENGLISH

S speaking to H in w_0 at t_0 (about Ann and John): $s_0 = \{\langle e, \langle \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, \dots \rangle \rangle\}$ entity-drefs
world-drefs
time-drefs
evt-drefs

(1) ¹Ann has gone out. ²John is asleep.
Ann₉₇^T have-PRS go.out-PRF John₉₆^T be-PRS asleep

¹Ann₉₇^T ~ Msk 96:151–3 + ctr

[$! u_{97} = ann, u_{97} \neq u_0$]; [$u_0 | u_0 = u_{97}$]
 $s_{11} = \{\langle e, \langle A, \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, \dots \rangle \rangle\}$

have-PRS ~ Msk 95:169, 173 + ctr

[$! (\exists E_0 = T), (E_0 \text{ in } W)$]; [$E_2 | E_2 = E_0$]; [$E_0 | E_0 < E_2$];
 $s_{12} = \{\langle e, \langle A, \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_1, ?, e_0, \dots \rangle \rangle\}$

go.out-PRF ~ Msk 95:172–3 + ctr

[$E_0: go.out_{\text{et}} u_0$]; [$E_1 | E_1 = E_0$]; [$E_0 | E_1 \bullet < E_0$]; [$E_0 | E_0 = E_2$];
 $s_{13} = \{\langle e, \langle A, \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, e_1, e_0, \dots \rangle \rangle\}$

²John₉₆^T ~ Msk 96:151–3 + ctr

[$! u_{96} = john, u_{96} \neq u_0$]; [$u_0 | u_0 = u_{96}$]
 $s_{21} = \{\langle e, \langle J, \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, e_1, e_0, \dots \rangle \rangle\}$

be-PRS ~ Msk 95:172–3 + ctr

[$! (\exists E_0 = T), (E_0 \text{ in } W)$]; [$E_1 | E_0 \subseteq E_1$];
 $s_{22} = \{\langle e, \langle J, \dots, J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, e_3, e_0, \dots \rangle \rangle\}$

asleep
[$E_1: asleep_{\text{et}} u_0$];

speech world w_0

|

t_0 : speech time

•

$e_0: \exists e_0 = t_0$

•

^{w0?}($e_1: \exists e_1 < \exists e_0 = t_0$, Ann goes out)

•

^{w0?}($e_2: \exists e_1 \bullet < \exists e_2$)

•

^{w0?}($e_3: t_0 = \exists e_0 \subseteq \exists e_2$, John is asleep)

(2) ¹When I came home, ²Ann went out. ³John was asleep.
 when $1s_T$ come.home-PST Ann₉₇^T go.out-PST John₉₆^T be-PST asleep

¹when ~ Msk 95:178 + ctr
 $[E_1 | E_1 = E_0]; [E_0 | E_1 < E_0]$ $s_{11} = \{\langle e, \langle \dots, J, A, H, S \rangle \rangle$

$\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_1, e_0, \dots \rangle \rangle\}$

$1s_T$ ~ Msk 95:166 + ctr
 $[u_0 | u_0 = u]$ $s_{12} = \{\langle e, \langle R, \dots, J, A, H, S \rangle \rangle$

$\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_1, e_0, \dots \rangle \rangle\}$

come.home-PST ~ Msk 95:172-3 + ctr

$[\exists E_0 < T, (E_0 \text{ in } W)];$
 $[E_0 : cmh_{\varepsilon et} u_0]; [E_1 | E_0 = E_1]; [E_0 | E_1 \bullet < E_0]$

$s_{13} = \{\langle e, \langle R, \dots, J, A, H, S \rangle \rangle,$
 $\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_2, e_1, \dots \rangle \rangle\}$

Ann₉₇^T
 $[u_{97} = ann, u_{97} \neq u_0]; [u_0 | u_0 = u_{97}]$

~ Msk 96:151-3 + ctr
 $s_{21} = \{\langle e, \langle A, \dots, J, A, H, S \rangle \rangle$
 $\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_2, e_1, \dots \rangle \rangle\}$

go.out-PST ~ Msk 95:172-3 + ctr
 $[\exists E_0 < T, (E_0 \text{ in } W)];$
 $[E_0 : go.out_{\varepsilon et} u_0]; [E_1 | E_0 = E_1]; [E_0 | E_1 \bullet < E_0]$

$s_{22} = \{\langle e, \langle A, \dots, J, A, H, S \rangle \rangle,$
 $\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_3, e_2, \dots \rangle \rangle\}$

John₉₆^T
 $[u_{96} = john, u_{96} \neq u_0]; [u_0 | u_0 = u_{96}]$

~ Msk 96:151-3 + ctr
 $s_{31} = \{\langle e, \langle J, \dots, J, A, H, S \rangle \rangle$
 $\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_3, e_2, \dots \rangle \rangle\}$

be-PST ~ Msk 95:172-3 + ctr
 $[\exists E_0 < T, (E_0 \text{ in } W)];$
 $[E_1 | E_0 \subseteq E_1];$

$s_{33} = \{\langle e, \langle J, \dots, J, A, H, S \rangle \rangle,$
 $\langle \omega, \langle \dots, w_0 \rangle \rangle,$
 $\langle \tau, \langle \dots, t_0 \rangle \rangle,$
 $\langle \varepsilon, \langle e_3, e_4, \dots \rangle \rangle\}$

asleep
 $[E_1 : asleep_{\varepsilon et} u_0]$

speech world w_0

|

t_0 : speech time

e_0 :
 e_1 : $\exists e_0 < \exists e_1 < t_0$, s_0 -speaker S comes home
 e_2 : $\exists e_1 \bullet < \exists e_2 < t_0$, Ann goes out
 e_3 : $\exists e_2 \bullet < \exists e_3 < t_0$
^{w0?} e_4 : $\exists e_3 \subseteq \exists e_4$, John is asleep)



2 SORTED LOGIC OF CHANGE (LCh_5): MODELS INSTEAD OF AXIOMS

D1.1 (LCh_5 -types)

- $t, e, \varepsilon, \omega, \tau, s \in \mathbf{Typ}$
- $(ab) \in \mathbf{Typ}$ if $a, b \in \mathbf{Typ}$

D1.2 (Basic LCh_5 -terms)

- $\mathbf{Con}_e = \{john_e, mary_e, \dots\}$
 $\mathbf{Con}_{sa} = \{u_{0,a}, \dots, u_{99,a}\}$, for all $a \in \{e, \varepsilon, \tau, \omega\}$
 $\mathbf{Con}_{\varepsilon\tau} = \{\emptyset, \dots\}$
 $\mathbf{Con}_{\tau t} = \{<, \dots\}$
 $\mathbf{Con}_{\omega\varepsilon t} = \{in, \dots\}$
 $\mathbf{Con}_{\varepsilon et} = \{agt_{\varepsilon et}, exp_{\varepsilon et}, enter_{\varepsilon et}, \dots, man_{\varepsilon et}, \dots\}$
 $\mathbf{Con}_{\varepsilon eet} = \{hit_{\varepsilon eet}, \dots, frn.of_{\varepsilon eet}, \dots\}$
 \vdots
- $\mathbf{Var}_e = \{x, y, z, x', y', z', \dots\}$ $\mathbf{Var}_{(se)} = \{\underline{x}, \underline{y}, \underline{z}, \dots\}$
 $\mathbf{Var}_\varepsilon = \{e, e', e'', \dots\}$ $\mathbf{Var}_{(sa)} = \{r_a, r_a', \dots\}$, for all $a \in \{e, \varepsilon, \omega, \tau\}$
 $\mathbf{Var}_\omega = \{w, w', w'', \dots\}$ $\mathbf{Var}_{sst} = \{p, q, \dots\}$
 $\mathbf{Var}_\tau = \{t, t', t'', \dots\}$ $\mathbf{Var}_{(se)(sst)} = \{P, P', \dots\}$
 $\mathbf{Var}_s = \{i, j, k, i', j', k', \dots\}$ $\mathbf{Var}_{(se)(se)(sst)} = \{R, R', \dots\}$

D1.3 (LCh_5 -syntax)

- $\alpha \in \mathbf{Term}_a$ if $\alpha \in \mathbf{Con}_a$ or $\alpha \in \mathbf{Var}_a$
- $\alpha\beta \in \mathbf{Term}_a$ if $\alpha \in \mathbf{Term}_{(ba)}$ and $\beta \in \mathbf{Term}_b$
- $(\lambda v\alpha) \in \mathbf{Term}_{(ba)}$ if $v \in \mathbf{Var}_b$ and $\alpha \in \mathbf{Term}_a$
- $(\alpha = \beta) \in \mathbf{Term}_t$ if $\alpha \in \mathbf{Term}_a$ and $\beta \in \mathbf{Term}_a$
- $(\alpha \wedge \beta) \in \mathbf{Term}_t$ if $\alpha \in \mathbf{Term}_t$ and $\beta \in \mathbf{Term}_t$
- $\neg\alpha \in \mathbf{Term}_t$ if $\alpha \in \mathbf{Term}_t$

D2.1 (LCh_5 -frames)

- $D_t = \{T, F\}$
 D_e, D_ε and D_ω are non-empty sets pairwise disjoint from each other and from D_t and D_τ
 $D_\tau = \{t \mid \emptyset \subset t \subseteq \{\dots, -1, 0, 1, \dots\}$
 $\quad \& \forall i, i', i'' \in \{\dots, -1, 0, 1, \dots\}: i \in t \& i'' \in t \& i < i' < i'' \rightarrow i' \in t\}$
 $D_s = \{s \mid \text{Dom } s = \{e, \varepsilon, \tau, \omega\} \& \forall a \in \text{Dom } s: s_a \in (D_a)^{100}\}$
- $D_{(ab)} = \{f \mid \text{Dom } f = D_a \& \text{Ran } f \subseteq D_b\}$

D2.2 (LCh_5 -models and assignments)

- An LCh_5 -model is a structure $M = \langle D^M, \llbracket \cdot \rrbracket^M \rangle$ such that:
 - $D^M = \{D_a^M: a \in \mathbf{Typ}\}$ is an LCh_5 -frame.
 - $\llbracket \cdot \rrbracket^M$ is a function that assigns to any $\alpha \in \mathbf{Con}_a$ a denotation $\llbracket \alpha \rrbracket^M \in D_a^M$.
 - $\forall s \in D_s: \llbracket u_{n,a} \rrbracket^M(s) = {}^{n+1}(s_a)$
 - $\forall t, t' \in D_\tau$:
 - $t <_\tau t'$ iff $\forall i \in t \forall i' \in t': i < i'$
 - $\llbracket < \rrbracket^M(t')(t) = T$ iff $t <_\tau t'$
- An M -assignment θ assigns to any variable $v \in \mathbf{Var}_a$ a value $\theta(v) \in D_a^M$. If $d \in D_a^M$, then $\theta[v/d]$ is the M -assignment s.t. (i) $\theta[v/d](v) = d$, and (ii) $\theta[v/d](v') = \theta(v')$ if $v' \neq v$.

D2.3 (LCh_5 -semantics).

- $\llbracket \alpha \rrbracket^{M, \theta} = \llbracket \alpha \rrbracket^M$ if $\alpha \in \mathbf{Con}_a$
- $\llbracket \alpha \rrbracket^{M, \theta} = \theta(\alpha)$ if $\alpha \in \mathbf{Var}_a$
- $\llbracket \alpha \beta \rrbracket^{M, \theta} = \llbracket \alpha \rrbracket^{M, \theta} (\llbracket \beta \rrbracket^{M, \theta})$
- $\llbracket (\lambda v \alpha) \rrbracket^{M, \theta} = \langle \llbracket \alpha \rrbracket^{M, \theta[v/d]} : d \in D_b^M \rangle$
- $\llbracket (\alpha = \beta) \rrbracket^{M, \theta} = \mathbf{T}$, iff $\llbracket \alpha \rrbracket^{M, \theta} = \llbracket \beta \rrbracket^{M, \theta}$
- $\llbracket (\alpha \wedge \beta) \rrbracket^{M, \theta} = \mathbf{T}$, iff $\llbracket \alpha \rrbracket^{M, \theta} = \llbracket \beta \rrbracket^{M, \theta} = 1$
- $\llbracket \neg \alpha \rrbracket^{M, \theta} = \mathbf{T}$, iff $\llbracket \alpha \rrbracket^{M, \theta} = \mathbf{F}$
- $\llbracket \alpha \rrbracket^{M, \theta} \in D_a^M$ if $\alpha \in \mathbf{Term}_a$

A1 (mostly standard abbreviations). We write:

- $(\phi \rightarrow \psi)$ for $\neg(\phi \wedge \neg\psi)$
- $(\phi \vee \psi)$ for $\neg(\neg\phi \wedge \neg\psi)$
- $\forall v \phi$ for $(\lambda v[\phi] = \lambda v[v = v])$
- $\exists v \phi$ for $\neg \forall v \neg \phi$
- $\forall v_1 \dots v_n \phi$ for $\forall v_1 \dots \forall v_n \phi$
- $\exists v_1 \dots v_n \phi$ for $\exists v_1 \dots \exists v_n \phi$
- $\lambda v_1 \dots v_n \alpha$ for $(\lambda v_1 \dots (\lambda v_n \alpha) \dots)$
- $\langle a_1, \dots, a_n \rangle$ for a_n , for any $1 \leq n \leq n'$ [nth coordinate]
- $(\alpha \beta \alpha')$ for $\beta \alpha' \alpha$
- $(t < t' < t'')$ for $(t < t' \wedge t' < t'')$
- $(t \bullet < t')$ for $(t < t' \wedge \neg \exists t''(t < t'' < t'))$
- $(t \subseteq t')$ for $\forall t''((t'' < t' \rightarrow t'' < t) \wedge (t' < t'' \rightarrow t < t''))$

Msk 95:167

A2 (DRT notation). We write

- \mathbf{DR}_a for $\lambda r_a (r_a = u_{0,a} \vee \dots \vee r_a = u_{99,a})$
- $i[\delta_a]j$ for $\forall r_a (\mathbf{DR}_a r_a \wedge r_a \neq \delta_a \rightarrow r_a j = r_a i)$
- $i[\delta_{1,a1} \dots \delta_{n,an}]j$ for $\exists k (i[\delta_{1,a1}]k \wedge k[\delta_{2,a2} \dots \delta_{n,an}]j)$
- u_n, u, u' for $u_{n,e}, u_{99,e}, u_{98,e}$
- T_n, T for $u_{n,\tau}, u_{99,\tau}$
- W_n, W for $u_{n,\omega}, u_{99,\omega}$
- E_n for $u_{n,\varepsilon}$
- $(\delta_{sa} = \delta'_{sa})$ for $\lambda i (\delta i = \delta' i)$ $a \in \{e, \varepsilon, \omega, \tau\}$
- $(\delta_{sa} \neq \delta'_{sa})$ for $\lambda i \neg (\delta i = \delta' i)$
- $(E_n : \alpha_{eet} \delta_{se})$ for $\lambda i \alpha E_n i \delta i$ [e.g., $(E_0 : cmh_{eet} u_0)$]
- $(E_n : \delta_{se} \alpha_{eet} \delta'_{se})$ for $\lambda i \alpha E_n i \delta' i \delta i$ [e.g., $(E_0 : u_0 see_{eet} u_1)$]
- $(E_n : \delta_{se} \alpha_{eet} \delta'_{se} \delta''_{se})$ for $\lambda i \alpha E_n i \delta'' i \delta' i \delta i$ [e.g., $(E_0 : u_0 give_{eet} u_1 u_2)$]
- $(\vartheta \delta_{se} < T)$ for $\lambda i (\vartheta \delta i < T i)$
- $(\delta_{se} \bullet < \delta'_{se})$ for $\lambda i (\vartheta \delta i \bullet < \vartheta \delta' i \wedge \exists w (\delta i \text{ in } w \wedge \delta' i \text{ in } w))$
- $(\delta_{se} \subseteq \delta'_{se})$ for $\lambda i (\vartheta \delta i \subseteq \vartheta \delta' i \wedge \exists w (\delta i \text{ in } w \wedge \delta' i \text{ in } w))$
- $(J_{sst} \Rightarrow K_{sst})$ for $\lambda i \forall j (Jij \rightarrow \exists k Kjk)$
- $(\mathbf{not} J_{sst})$ for $\lambda i \neg \exists j Jij$
- $[\delta_1 \dots \delta_n | \chi_1, \dots, \chi_m]$ for $\lambda ij (i[\delta_1 \dots \delta_n]j \wedge \chi_{1j} \wedge \dots \wedge \chi_{mj})$
- $[| \chi_1, \dots, \chi_m]$ for $\lambda ij (i = j \wedge \chi_{1j} \wedge \dots \wedge \chi_{mj})$
- $(J_{sst} ; K_{sst})$ for $\lambda ij \exists k (J_{sst} ik \wedge K_{sst} kj)$

D3 (Truth and equivalence).

- A t -term ϕ is *true* in M under θ , written $\models_{M,\theta} \phi$, iff $\llbracket \phi \rrbracket^{M,\theta} = 1$
- A t -term ϕ is *true* in M , written $\models_M \phi$, iff $\forall \theta: \models_{M,\theta} \phi$
- An *sst*-term K is *true* in M , written $\models_M K$, iff $\models_M \exists ij Kij$ [Kamp 1981 def. of truth]
- α is *equivalent* to β , written $\alpha \equiv \beta$, iff $\forall M, \theta: \llbracket \alpha \rrbracket^{M,\theta} = \llbracket \beta \rrbracket^{M,\theta}$

3 TRANSLATION RULES FOR ENGLISH (a la Muskens 1995, 1996)

			<i>Type</i>	<i>'Category'</i>
• \mathbf{B}_{Eng} . Base rule for <i>English</i>				
¹ John _{<i>n</i>}	\rightsquigarrow $[u_n = john_e^\circ, u_n \neq u_0]$	$n > 2$	\square	n
ⁿ \top	\rightsquigarrow $[u_0 u_n = u_0]$	$n > 2$	\square	n \ N (subject)
ⁿ \perp	\rightsquigarrow $[u_n = u_1]$	$n > 2$	\square	(object)
^{1'} I	\rightsquigarrow $[u_0 u_0 = u]$		\square	N
me	\rightsquigarrow $[u \neq u_0]; [u = u_1]$		\square	
he _{<i>n</i>}	\rightsquigarrow $[male_{et} u_n]; [u_0 u_n = u_0]$		\square	
him _{<i>n</i>}	\rightsquigarrow $[male_{et} u_n, u_n \neq u_0]; [u_n = u_1]$	$n > 2$	\square	
² go#out-	\rightsquigarrow $[E_0: go.out_{et} u_0]; [E_1 E_0 = E_1]; [E_0 E_1 \bullet < E_0]$		\square	event-v
^{2'} be- (A)	\rightsquigarrow $[E_1 E_0 \subseteq E_1]$		\square	A/state-v
^{2'} be- (PRG)	\rightsquigarrow $[E_2 E_2 = E_0]; [E_0 E_2 \subseteq E_0]$		\square	V/state-v (asp)
have- (PRF)	\rightsquigarrow $[E_2 E_2 = E_0]; [E_0 E_0 < E_2]$		\square	V/state-v (asp)
^{2''} asleep	\rightsquigarrow $[E_1: asleep_{et} u_0]$		\square	A
-PRG, -PRF	\rightsquigarrow $[E_0 E_0 = E_2]$		\square	v \ V (asp)
³ -PRS	\rightsquigarrow $\lambda p ([(\exists E_0 = T), (E_0 \text{ in } W)]; p)$		$\square \square$	v \ V (tns)
-PST	\rightsquigarrow $\lambda p ([(\exists E_0 < T), (E_0 \text{ in } W)]; p)$		$\square \square$	
⁴ when	\rightsquigarrow $[E_1 E_1 = E_0]; [E_0 E_1 < E_0]$		\square	C

- *Application rule (A)*
If A and B are sisters, $A \rightsquigarrow \alpha$, and $B \rightsquigarrow \beta$,
then $[A, B] \rightsquigarrow \alpha\beta$, provided that this is an LCh₅-term.
- *Sequencing rule (;)*
If A and B are sisters, $A \rightsquigarrow \alpha$, $B \rightsquigarrow \beta$,
then $[A B] \rightsquigarrow (\alpha; \beta)$, provided that this is an LCh₅-term.
- *Property sequencing rule (;')*
If A and B are sisters, $A \rightsquigarrow \alpha$, $B \rightsquigarrow \beta$,
then $[A B] \rightsquigarrow \lambda r(\alpha r; \beta r)$, provided that this is an LCh₅-term.
- *Reduction rule (R)*
If $A \rightsquigarrow \alpha$, and $\alpha \equiv \beta$, then $A \rightsquigarrow \beta$.

4 FROM ENGLISH TO KALAALLISUT (Shaer 2003)

Shaer 2003 ('Toward a tenseless analysis of a tenseless language', *SULA* 3) proposes to extend Muskens 1995 analysis to tenseless Kalaallisut as follows: Kalaallisut is just like English minus the temporal tests of the English tenses. The following base rule \mathbf{B}_{Ksl} implements this idea:

Abbreviations: $(t = t' \cup t'')$ for $t' \subseteq t \wedge t'' \subseteq t \wedge \forall t''' (t' \subseteq t''' \wedge t'' \subseteq t''' \rightarrow t \subseteq t''')$
 $(\vartheta\delta_{se} = \vartheta\delta_{se}' \cup \vartheta\delta_{se}'')$ for $\lambda i(\vartheta\delta i = \vartheta\delta' i \cup \vartheta\delta'' i)$

• B_{Kal} . Base rule for <i>Kalaallisut</i>		Type	'Category'
¹ John-	$\rightsquigarrow \lambda r[r = john_e^\circ, r \neq u_0]$	[se]	cn (name)
-sg _n	$\rightsquigarrow \lambda P(Pu_n; [u_0 u_0 = u_n])$	[se] []	cn\CN (subj.)
-sg _n [⊥]	$\rightsquigarrow \lambda P(Pu_n; [u_1 u_1 = u_n])$	[se] []	cn\CN (obj.)
² go.out-	$\rightsquigarrow [E_0: go.out_{\varepsilon et} u_0]; [E_1 E_0 = E_1]; [E_0 E_1 \bullet < E_0]$	[]	event-iv
	$\rightsquigarrow [u_0 E_0: go.out_{\varepsilon et} u_0]; [E_1 E_0 = E_1]; [E_0 E_1 \bullet < E_0]$	[]	
be.asleep-	$\rightsquigarrow [E_1 (E_1: asleep_{\varepsilon et} u_0), (E_0 \subseteq E_1)]$	[]	state-iv
	$\rightsquigarrow [u_0 E_1 (E_1: asleep_{\varepsilon et} u_0), (E_0 \subseteq E_1)]$	[]	
^{2'} -prf	$\rightsquigarrow \lambda p([E_2 E_2 = E_0]; [E_0 E_0 < E_2]; p; [E_0 E_0 = E_2])$	[] []	iv\iv (asp)
-begin	$\rightsquigarrow ([E_2 (E_0 \bullet < E_2), (\exists E_1 = \exists E_0 \cup \exists E_2)];$ $[E_1 E_1 = E_0]; [E_0 E_1 \bullet < E_0])$	[]	
³ -IND.IV	$\rightsquigarrow [E_0 \text{ in } W]$	[]	iv\IV' (mood)
-FCT _T	$\rightsquigarrow [E_1 E_1 = E_0]; [E_0 (E_1 < E_0), (E_0 \text{ in } W)];$	[]	
-FCT _⊥	$\rightsquigarrow [E_1 E_1 = E_0]; [E_0 (E_1 < E_0), (E_0 \text{ in } W)]; [u_1 u_1 = u_0]$	[]	
^{3'} -1s	$\rightsquigarrow [u_0 = u]$	[]	IV'/IV (1s agr)
-3s _(T)	$\rightsquigarrow [u_0 \neq u, u_0 \neq u']$	[]	(3s _T agr)
-3s _⊥	$\rightsquigarrow [u_1 \neq u, u_1 \neq u']$	[]	(3s _⊥ agr)

5 TEMPORAL ANAPHORA IN KALAALLISUT

S speaking to H in w_0 at t_0 (about Ann and John)	$s_0 = \{\langle e, \langle \dots, J, A, H, S \rangle \rangle,$	<i>entity-drefs</i>
Locate initial reference point, e_0 , in w_0 at t_0	$\langle \omega, \langle \dots, w_0 \rangle \rangle,$	<i>world-drefs</i>
(null assumption)	$\langle \tau, \langle \dots, t_0 \rangle \rangle,$	<i>time-drefs</i>
	$\langle \varepsilon, \langle e_0, \dots \rangle \rangle \}$	<i>evt-drefs</i>

speech world w_0 | t_0 : *speech time*
• e_0 : $\exists e_0 = t_0$

Kalaallisut (1K) and (2K) are *translation equivalents* of English (1) and (2), while (3K) and (4K) illustrate some (simple or recursive) *aspectual shifts*:

(1K) ¹ Aani	<i>ani-vu-q.</i>	² Juuna	<i>sinip-pu-q.</i>
Ann-sg ₉₇	go.out-IND.IV-3s	John-sg ₉₆	be.asleep-IND.IV-3s
(2K) ¹ Angirlar-a-ma	² Aani	<i>ani-vu-q.</i>	³ Juuna
come.home-FCT _T -1s	Ann-sg ₉₇	go.out-IND.IV-3s	John-sg ₉₆
			be.asleep-IND.IV-3s
(3K) ¹ Aani	<i>ani-sima-vu-q.</i>	² Juuna	<i>sini-lir-pu-q.</i>
Ann-sg ₉₅	go.out-prf-IND.IV-3s	Juuna-sg ₉₇	be.asleep-IND.IV-3s
(4K) ¹ Juuna	<i>sini-lir-sima-lir-(m)m-at</i>	² Aani	<i>ani-vu-q.</i>
John-sg ₉₅	be.asleep-begin-prf-begin-FCT _⊥ -3s _⊥	Ann-sg ₉₇	go.out-IND.IV-3s

Kalaallisut (1K)

• Predicted interpretation

(1K) ¹Ann.sg₉₇

[$u_{97} = ann, u_{97} \neq u_0$]; [$u_0 | u_0 = u_{97}$]

$s_{11} = \{ \langle e, \langle A, \dots J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_0, \dots \rangle \rangle \}$

go.out-IND.IV-3s

[$E_0: go.out_{\varepsilon et} u_0$]; [$E_1 | E_0 = E_1$]; [$E_0 | E_1 \bullet < E_0$]; [$s_{12} = \{ \langle e, \langle A, \dots J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_1, e_0, \dots \rangle \rangle \}$]; [$E_0 in W$]; [$u_0 \neq u, u_0 \neq u'$]

²John.sg₉₆

[$u_{96} = john_e, u_{96} \neq u_0$]; [$u_0 | u_0 = u_{96}$]

$s_{21} = \{ \langle e, \langle J, \dots J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_1, e_0, \dots \rangle \rangle \}$

be.asleep-IND.IV-3s

[$E_1 | (E_1: asleep_{\varepsilon et} u_0), (E_0 \subseteq E_1)$]; [$s_{22} = \{ \langle e, \langle J, \dots J, A, H, S \rangle \rangle, \langle \omega, \langle \dots, w_0 \rangle \rangle, \langle \tau, \langle \dots, t_0 \rangle \rangle, \langle \varepsilon, \langle e_1, e_2, \dots \rangle \rangle \}$]; [$E_0 in W$]; [$u_0 \neq u, u_0 \neq u'$]

speech world w_0

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t_0 : *speech time*

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$e_0: \vartheta e_0 = t_0, \quad \text{Ann goes out}$

•

$e_1: \vartheta e_0 \bullet < \vartheta e_1$

●

^{w0?} $(e_2: \vartheta e_1 \subseteq \vartheta e_2, \quad \text{John is asleep})$

• Actual interpretation

(1K) ¹Aani ani-vu-q.

²Juuna sinip-pu-q.

\approx Eng. (1)

Ann-sg₉₇ go.out-IND.IV-3s John-sg₉₆ be.asleep-IND.IV-3s

Ann has gone out. John is asleep.

Kalaallisut (2K)

• Predicted interpretation

(2K) come.home-FCT_T-1s

$$[u_0 | E_0: cmh_{e_0} u_0]; [E_1 | E_0 = E_1]; [E_0 | E_1 \bullet < E_0]; \mathbf{s}_{12} = \{ \langle e, \langle S, \dots, J, A, H, S \rangle \rangle, \\ [E_1 | E_1 = E_0]; [E_0 | (E_1 < E_0), (E_0 \text{ in } W)]; \langle \omega, \langle \dots, w_0 \rangle \rangle, \\ [| u_0 = u] \langle \tau, \langle \dots, t_0 \rangle \rangle, \\ \langle \varepsilon, \langle e_2, e_1, \dots \rangle \rangle \}$$

²Ann.sg₉₇

$$[| u_{97} = ann, u_{97} \neq u_0]; [u_0 | u_0 = u_{97}] \mathbf{s}_{21} = \{ \langle e, \langle A, \dots, J, A, H, S \rangle \rangle, \\ \langle \omega, \langle \dots, w_0 \rangle \rangle, \\ \langle \tau, \langle \dots, t_0 \rangle \rangle, \\ \langle \varepsilon, \langle e_2, e_1, \dots \rangle \rangle \}$$

go.out-IND.IV-3s

$$[| E_0: go.out u_0]; [E_1 | E_0 = E_1]; [E_0 | E_1 \bullet < E_0]; \mathbf{s}_{22} = \{ \langle e, \langle A, \dots, J, A, H, S \rangle \rangle, \\ [| E_0 \text{ in } W]; [| u_0 \neq u, u_0 \neq u'] \langle \omega, \langle \dots, w_0 \rangle \rangle, \\ \langle \tau, \langle \dots, t_0 \rangle \rangle, \\ \langle \varepsilon, \langle e_3, e_2, \dots \rangle \rangle \}$$

³John.sg₉₆

$$[| u_{96} = john_e, u_{96} \neq u_0]; [u_0 | u_0 = u_{96}] \mathbf{s}_{31} = \{ \langle e, \langle J, \dots, J, A, H, S \rangle \rangle, \\ \langle \omega, \langle \dots, w_0 \rangle \rangle, \\ \langle \tau, \langle \dots, t_0 \rangle \rangle, \\ \langle \varepsilon, \langle e_3, e_2, \dots \rangle \rangle \}$$

be.asleep-IND.IV-3s_T

$$[E_1 | (E_1: asleep u_0), (E_0 \subseteq E_1)]; \mathbf{s}_{32} = \{ \langle e, \langle J, \dots, J, A, H, S \rangle \rangle, \\ [| E_0 \text{ in } W]; [| u_0 \neq u, u_0 \neq u'] \langle \omega, \langle \dots, w_0 \rangle \rangle, \\ \langle \tau, \langle \dots, t_0 \rangle \rangle, \\ \langle \varepsilon, \langle e_3, e_4, \dots \rangle \rangle \}$$

speech world w_0

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t_0 : *speech time*

e_0 : $\vartheta e_0 = t_0$, speaker comes home

^{w0?} $(e_1: \vartheta e_0 \bullet < \vartheta e_1)$

e_2 : $\vartheta e_1 < \vartheta e_2$, Ann goes out

e_3 : $\vartheta e_2 \bullet < \vartheta e_3$

^{w0?} $(e_4: \vartheta e_3 \subseteq \vartheta e_3, \text{ John is asleep})$

• Actual interpretation

(2K) ¹Angirlar-a-ma ²Aani ani-vu-q. ³Juuna sinip-pu-q. ≈ Eng (2)

come.home-FCT_T-1s Ann-sg₉₇ go.out-IND.IV-3s John-sg₉₅ be.asleep-IND.IV-3s

When I came home, Ann went out. John was asleep.

Kalaallisut (3K) & (4K)

See analysis at <http://www.rci.rutgers.edu/~mbittner/lct.html>, Semantics 2, top 6b, pp.15–16.

6 CONCLUSION

Wrong approach. A tenseless language cannot be analyzed like a tensed language minus tense.