

Lectures 17-18
CROSS-LINGUISTIC COMPOSITION

1. Semantic convergence across structural diversity

ENGLISH (E)	KALAALLISUT (K)
<ul style="list-style-type: none"> • pronoun (Ole is home today.) 	<ul style="list-style-type: none"> • -PRN
(1 _E) He is sick. he be.TNS sick [<i>sick</i> _{τ_ω} τδ]; [p p = τω{}]	(1 _K) Naparsimavuq naparsima-pu-q sick-DEC.IV-3SG [<i>sick</i> _{τ_ω} τδ]; [p p = τω{}]
<ul style="list-style-type: none"> • name 	<ul style="list-style-type: none"> • name
(2 _E) Ole is sick. Ole be.TNS sick [x (x = ole) ^o]; [<i>sick</i> _{τ_ω} τδ]; [p p = τω{}]	(2 _K) Ole naparsimavuq. Ole naparsima-pu-q Ole sick-DEC.IV-3SG [x (x = ole) ^o]; [<i>sick</i> _{τ_ω} τδ]; [p p = τω{}]
<ul style="list-style-type: none"> • indef. object of tv 	<ul style="list-style-type: none"> • cn-base of iv\cn-suffix
(3 _E) Ole has bought a dog. Ole hv.TNS buy.ASP A dog [x (x = ole) ^o]; [y <i>dog</i> _{τ_ω} y, <i>buy</i> _{τ_ω} ⟨τδ, y⟩]; [p p = τω{}]	(3 _K) Ole qimmisivuq. Ole <i>qimmiq-si-pu-q</i> Ole dog-get-DEC.IV-3SG [x (x = ole) ^o]; [y <i>buy</i> _{τ_ω} ⟨τδ, y⟩, <i>dog</i> _{τ_ω} y]; [p p = τω{}]
<ul style="list-style-type: none"> • n-modifier 	<ul style="list-style-type: none"> • s-modifier (MOD) + [_s cn-iv\cn-...]
(4 _E) Ole has bought a sick dog. Ole hv.TNS buy.ASP A sick dog [x (x = ole) ^o]; [y <i>dog</i> _{τ_ω} y, <i>sick</i> _{τ_ω} y, <i>buy</i> _{τ_ω} ⟨τδ, y⟩]; [p p = τω{}]	(4 _K) Ole naparsimasumik qimmisivuq. Ole <i>naparsima-tuq-mik</i> qimmiq-si-pu-q Ole sick-n\vp-MOD dog-get-DEC.IV-3SG [x (x = ole) ^o]; [y <i>sick</i> _{τ_ω} y]; [X X = ⊥δ{ τ _ω }]; [<i>buy</i> _{τ_ω} ⟨τδ, ⊥δ⟩, <i>dog</i> _{τ_ω} ⊥δ]; [p p = τω{}]
<ul style="list-style-type: none"> • negation 	<ul style="list-style-type: none"> • negation
(5 _E) Ole hasn't bought a sick dog. Ole hv.TNS=N'T buy.ASP A sick dog rdg1 [x (x = ole) ^o]; [w w ∈ τω{}]; [y <i>dog</i> _{⊥ω} y, <i>sick</i> _{⊥ω} y, <i>buy</i> _{⊥ω} ⟨τδ, y⟩]; [p p = ⊥ω{}]; [τω ∉ ⊥Ω]; [p p = τω{}]	(5 _K) Ole naparsimasumik qimmisinngilaq. Ole naparsima-tuq-mik qimmi-si-nngit-la-q Ole sick-n\vp-MOD dog-get-not-DEC.NG-3SG [x (x = ole) ^o]; [w w ∈ τω{}]; [y <i>sick</i> _{⊥ω} y]; [X X = ⊥δ{ ⊥ω}]; [<i>get</i> _{⊥ω} ⟨τδ, ⊥δ⟩, <i>dog</i> _{⊥ω} ⊥δ]; [p p = ⊥ω{}]; [τω ∉ ⊥Ω]; [p p = τω{}]
rdg2 [x (x = ole) ^o]; [y <i>dog</i> _{τ_ω} y, <i>sick</i> _{τ_ω} y]; [w w ∈ τω{}]; [<i>buy</i> _{⊥ω} ⟨τδ, ⊥δ⟩]; [p p = ⊥ω{}]; [τω ∉ ⊥Ω]; [p p = τω{}]	

Homework 7 (to be assigned M 4/20, due Th 4/23)

2. Update with Centering (UC2) and presupposition

DEFINITION 1.1 The set of UC2 types is the smallest set **Typ** such that:

- $t, \delta, \omega, s \in \mathbf{Typ}$
- $(ab) \in \mathbf{Typ}$, if $a, b \in \mathbf{Typ}$

The set of discourse object types is the subset $\mathbf{DTyp} = \{\delta, \omega, \omega t\} \cup \{(at): a \in \{\delta, \omega, \omega t\}\}$

ABBREVIATIONS 1

$W := s\omega$	$\Omega := \omega t$	$a_1 \dots a_n b := (a_1 \dots (a_n b))$
$D := s\delta$	$[] := (st)st$	$a_1 \dots a_n [] := [a_1 \dots a_n]$

DEFINITION 1.2 (Basic UC2 terms). For each type $a \in \mathbf{Typ}$, there is a set of a -constants, \mathbf{Con}_a , a set of plain a -variables, ${}^{\perp}\mathbf{Var}_a$, and a set of topic-setting a -variables, ${}^{\top}\mathbf{Var}_a$, including:

$a \in \mathbf{Typ}$	${}^{\perp}\mathbf{Var}_a$	${}^{\top}\mathbf{Var}_a$	\mathbf{Con}_a	Name of objects
δ	x, y	\mathbf{x}, \mathbf{y}	<i>ole</i>	(ordinary) individuals
ω	w, v	\mathbf{w}, \mathbf{v}		worlds
Ω	p, q	\mathbf{p}, \mathbf{q}		propositions
δt	X, Y	\mathbf{X}, \mathbf{Y}		sets of individuals
Ωt	Q	\mathbf{Q}		sets of propositions
$\omega \delta t$			<i>dog, sick, ...</i>	$\omega \delta$ -relations
s	i, j, k, h			indices (<i>aka</i> topic-background lists)
st	I, J			infotention states
$[]$	K			updates
D	$\underline{x}, \underline{y}$			dynamic δ -concepts
W	$\underline{w}, \underline{v}$			dynamic ω -concepts
$[W]$	$\underline{W}, \underline{V}$			dynamic propositions
$[DW]$	\underline{P}			dynamic verbal properties
$[WD]$	\underline{N}			dynamic nominal properties

DEFINITION 1.3 (UC2 syntax). For each type $a \in \mathbf{Typ}$ the set of a -terms is defined as follows:

b	$A \in \mathbf{Term}_a$	if $A \in \mathbf{Con}_a \cup {}^{\top}\mathbf{Var}_a \cup {}^{\perp}\mathbf{Var}_a$
a	$\top a_n, \perp a_n \in \mathbf{Term}_{sa}$	if $a \in \mathbf{DTyp}$ & $n \in \{1, 2, \dots\}$
$\{ \}$	$A \{B\} \in \mathbf{Term}_{at}$	if $a \in \mathbf{DTyp}$, $A \in \mathbf{Term}_{sa}$ & $B \in \mathbf{Term}_{st}$
\cdot	$(u_a \cdot B) \in \mathbf{Term}_s$	if $a \in \mathbf{DTyp}$, $u_a \in {}^{\top}\mathbf{Var}_a \cup {}^{\perp}\mathbf{Var}_a$ & $B \in \mathbf{Term}_s$
$;$	$(A; B), (A^{\top}; B), (A^{\perp}; B) \in \mathbf{Term}_{(st)st}$	if $A, B \in \mathbf{Term}_{(st)st}$
λ	$\lambda u_a(B) \in \mathbf{Term}_{ab}$	if $u_a \in {}^{\top}\mathbf{Var}_a \cup {}^{\perp}\mathbf{Var}_a$ & $B \in \mathbf{Term}_b$
A	$AB \in \mathbf{Term}_a$	if $A \in \mathbf{Term}_{ba}$ & $B \in \mathbf{Term}_b$
C	$\neg A, (A \rightarrow B), (A \wedge B), (A \vee B) \in \mathbf{Term}_t$	if $A, B \in \mathbf{Term}_t$
Q	$\forall u_a B, \exists u_a B \in \mathbf{Term}_t$	if $u_a \in {}^{\top}\mathbf{Var}_a \cup {}^{\perp}\mathbf{Var}_a$ & $B \in \mathbf{Term}_t$
$=$	$(A = B) \in \mathbf{Term}_t$	if $A, B \in \mathbf{Term}_a$

DEFINITION 2.1 (UC2 frames). A UC2 frame is a set of sets $\{D_a\}_{a \in \mathbf{Typ}}$ such that:

- $D_t = \{1, 0\}$, D_{δ} , and D_{ω} are non-empty pairwise disjoint sets
- $D_s = \bigcup_{n, m \geq 0} (D^n \times D^m)$, where $D = \bigcup_{a \in \mathbf{DTyp}} D_a$
- $D_{ab} = \{f \mid \emptyset \subset \text{Dom } f \subseteq D_a \text{ \& } \text{Ran } f \subseteq D_b\}$

ABBREVIATIONS 2 (lists, projections, dot-extensions). Let $D = \bigcup_{a \in \mathbf{DTyp}} D_a$.

- For $\dot{i} = \langle \dot{i}_1, \dot{i}_2 \rangle \in D^n \times D^m$, $\top \dot{i} = \dot{i}_1$ is the *topic list* & $\perp \dot{i} = \dot{i}_2$ is the *background list* of \dot{i}
- For $\mathbf{x} \in D^{n+m}$, $(\mathbf{x})_n$ is \mathbf{x}_n (n th coordinate) & $(\mathbf{x})_a$ is the subsequence of type a coordinates
- For $\mathbf{d} \in D$, $\mathbf{x} \in D^n$: $\langle \mathbf{d} \cdot \langle \rangle \rangle = \langle \mathbf{d} \rangle$ & $\langle \mathbf{d} \cdot \mathbf{x} \rangle := \langle \mathbf{d}, \mathbf{x}_1, \dots, \mathbf{x}_n \rangle$ (*d-extensions*)
- For $\mathbf{x} \in D^n$, $\mathbf{y} \in D^{n+m}$: \mathbf{y} *dot-extends* \mathbf{x} , $\mathbf{y} \cdot \mathbf{x}$, iff $\exists \mathbf{d}_1 \dots \mathbf{d}_m \in D$: $\mathbf{y} = (\mathbf{d}_1 \cdot \dots \cdot (\mathbf{d}_m \cdot \mathbf{x}))$

DEFINITION 2.2 (UC2-models and assignments)

- A UC2-model is a pair, $M = \langle \{D_a\}_{a \in \text{Typ}}, \llbracket \cdot \rrbracket \rangle$, s.t. (i) $\{D_a\}_{a \in \text{Typ}}$ is a standard UC2-frame, (ii) $\llbracket \cdot \rrbracket$ is an *interpretation function* that assigns to each $\alpha \in \mathbf{Con}_a$ a denotation $\llbracket \alpha \rrbracket \in D_a$.
- An M -assignment is a function g that assigns $g(u) \in D_a$ to each $u \in {}^T\mathbf{Var}_a \cup {}^\perp\mathbf{Var}_a$. If $d \in D_a$ then $g[u/d]$ is the M -assignment s.t. (i) $g[u/d](u') = g(u')$ for all $u' \neq u$, (ii) $g[u/d](u) = d$.

ABBREVIATIONS 3 (Functions & sets). For $f \in D_{a_1 \dots a_n t}$ and $\langle a_1, \dots, a_n \rangle \in D_{a_1} \times \dots \times D_{a_n}$:

- $f(a_1, \dots, a_n) := f(a_1) \dots (a_n)$
- $\overset{\circ}{\cup}(f) = \{\langle a_1, \dots, a_n \rangle : f(a_1, \dots, a_n) = 1\}$ (set characterized by function f)
- ${}^\times(A)$ (characteristic function of set A) is the function f such that $\overset{\circ}{\cup}(f) = A$

DEFINITION 2.3 (UC2 semantics). The value $\llbracket A \rrbracket^g$ of a term A on a model $M = \langle \{D_a\}_{a \in \text{Typ}}, \llbracket \cdot \rrbracket \rangle$ under an assignment g is defined as follows (Note: we write (i) $c f$ for $f(c)$, (ii) ' $X \doteq Y$ ' for ' X is Y , if Y is defined, else X is undefined', (iii) ' $X[Y/Z]$ ' for the result of replacing every occurrence of Y in X with Z , and (iv) use the Von Neumann definition of 0 and 1, so $0 = \emptyset$ and $1 = \{\emptyset\}$)

b	$\llbracket A \rrbracket^g$	$= \llbracket A \rrbracket$	if $A \in \mathbf{Con}_a$
		$= g(A)$	if $A \in {}^T\mathbf{Var}_a \cup {}^\perp\mathbf{Var}_a$
a	$\llbracket \top a_n \rrbracket^g(i)$	$\doteq ((\top i)_{a_n})$	for any $i \in D_s$
	$\llbracket \perp a_n \rrbracket^g(i)$	$\doteq ((\perp i)_{a_n})$	for any $i \in D_s$
$\{$	$\llbracket A \{B\} \rrbracket^g$	$\doteq {}^\times\{\llbracket A \rrbracket^g(i) \mid i \in \overset{\circ}{\cup}\llbracket B \rrbracket^g\}$	
\cdot	$\llbracket (u_a \cdot B) \rrbracket^g$	$\doteq \langle (g(u_a) \cdot \top \llbracket B \rrbracket^g), \perp \llbracket B \rrbracket^g \rangle$	for any $u_a \in {}^T\mathbf{Var}_a$
		$\doteq \langle \top \llbracket B \rrbracket^g, (g(u_a) \cdot \perp \llbracket B \rrbracket^g) \rangle$	for any $u_a \in {}^\perp\mathbf{Var}_a$
$;$	$c \llbracket (A; B) \rrbracket^g$	$\doteq c \llbracket A \rrbracket^g \llbracket B \rrbracket^g$	for any $c \in D_{st}$
	$c \llbracket A^\top; B \rrbracket^g$	$\doteq \{k \in c \llbracket (A; B) \rrbracket^g \mid \exists a (\exists j \in c \llbracket A \rrbracket^g \exists i \in c : \top j \cdot \top i \ \& \ (\top j)_1 \in D_a \ \& \ \exists n : \llbracket B \rrbracket^g \neq \llbracket B[\top a_1 / \top a_n] \rrbracket^g)\}$	
	$c \llbracket A^\perp; B \rrbracket^g$	$\doteq \{k \in c \llbracket (A; B) \rrbracket^g \mid \exists a (\exists j \in c \llbracket A \rrbracket^g \exists i \in c : \perp j \cdot \perp i \ \& \ (\perp j)_1 \in D_a \ \& \ \exists n : \llbracket B \rrbracket^g \neq \llbracket B[\perp a_1 / \perp a_n] \rrbracket^g)\}$	
λ	$\llbracket \lambda u_a (B) \rrbracket^g(d)$	$\doteq \llbracket B \rrbracket^{g[u/d]}$	for any $d \in D_a$
A	$\llbracket AB \rrbracket^g$	$\doteq \llbracket A \rrbracket^g (\llbracket B \rrbracket^g)$	
C	$\llbracket \neg A \rrbracket^g$	$\doteq (1 - \llbracket A \rrbracket^g)$	
	$\llbracket (A \rightarrow B) \rrbracket^g$	$\doteq (1 - (\llbracket A \rrbracket^g - \llbracket B \rrbracket^g))$	
	$\llbracket (A \wedge B) \rrbracket^g$	$\doteq (\llbracket A \rrbracket^g \cap \llbracket B \rrbracket^g)$	
	$\llbracket (A \vee B) \rrbracket^g$	$\doteq (\llbracket A \rrbracket^g \cup \llbracket B \rrbracket^g)$	
Q	$\llbracket \forall u_a A \rrbracket^g$	$= 1$	if $\forall d \in D_a : \llbracket A \rrbracket^{g[u/d]} = 1$
		$= 0$	otherwise
	$\llbracket \exists u_a A \rrbracket^g$	$= 1$	if $\exists d \in D_a : \llbracket A \rrbracket^{g[u/d]} = 1$
		$= 0$	otherwise
$=$	$\llbracket (A = B) \rrbracket^g$	$= 1$	if $\llbracket A \rrbracket^g, \llbracket B \rrbracket^g \in D_a \ \& \ \llbracket A \rrbracket^g = \llbracket B \rrbracket^g$
		$= 0$	if $\llbracket A \rrbracket^g, \llbracket B \rrbracket^g \in D_a \ \& \ \llbracket A \rrbracket^g \neq \llbracket B \rrbracket^g$

DEFINITION 3.1 (Initial context & default state). An *initial context* is a proposition $p_0 \in D_\Omega$ s.t.

$\overset{\circ}{\cup} p_0 \neq \emptyset$. This induces the *default infotention state* $*p_0 = {}^\times\{\langle w, p_0 \rangle, \langle \rangle \mid w \in \overset{\circ}{\cup} p_0\}$.

DEFINITION 3.2 (Topic, assertion, truth value). For a model M , initial context p_0 , and (st) st -term K :

- $\tau_{M, p_0} K := \{(\top i)_1 \mid \forall g : i \in \overset{\circ}{\cup} (*p_0 \llbracket K \rrbracket^g)\}$ (set of *primary topics* introduced by K wrt M and p_0)
- Given M and p_0 , K can be used to *assert* a proposition $q \in D_\Omega$, iff $\tau_{M, p_0} K = \{q\}$
- K is *true* in world w wrt M and p_0 , iff $\exists q \in D_\Omega : \tau_{M, p_0} K = \{q\} \ \& \ w \in \overset{\circ}{\cup} q$
 K is *false* in world w wrt M and p_0 , iff $\exists q \in D_\Omega : \tau_{M, p_0} K = \{q\} \ \& \ w \notin \overset{\circ}{\cup} q$

ABBREVIATIONS 4 (Syntactic sugar)

i. Set theory

$$\begin{aligned} A_a \in B_{at} &:= BA \\ A_a \notin B_{at} &:= \neg BA \end{aligned}$$

ii. Dynamic terms

$$\begin{aligned} \top a, \perp a &:= \top a_1, \perp a_1 && \text{if } a \in \mathbf{DTyp} \\ A_a^\circ &:= \lambda i(A) && \text{if } a \in \mathbf{DTyp} \\ &:= A && \text{if } a \in \{sb \mid b \in \mathbf{DTyp}\} \end{aligned}$$

iii. (Local) conditions

$$\begin{aligned} (B_W A) &:= \lambda i(B W^\circ i A^\circ i) \\ B_W \langle A_1, \dots, A_n \rangle &:= \lambda i(B W^\circ i A_1^\circ i \dots A_n^\circ i) \\ (A = B)^\circ &:= \lambda i(A^\circ i = B^\circ i) \\ (C_1, C_2) &:= \lambda i(C_1 i \wedge C_2 i) \end{aligned}$$

iv. Local drt-boxes

$$\begin{aligned} [C] &:= \lambda I \lambda j (I j \wedge C j) \\ [u_1 \dots u_n] &:= \lambda I \lambda j \exists u_1 \dots u_n \exists i (j = (u_1 \cdot \dots \cdot (u_n \cdot i)) \wedge I i) \\ [u_1 \dots u_n | C] &:= \lambda I \lambda j \exists u_1 \dots u_n \exists i (j = (u_1 \cdot \dots \cdot (u_n \cdot i)) \wedge I i \wedge C i) \end{aligned}$$

v. Substates & global drt-boxes

$$\begin{aligned} I_{B:i} &:= \lambda j (I j \wedge B j = B i) \\ [A' R A \{\}] &:= \lambda I \lambda j (I j \wedge A' R A \{I\}) && R \in \{=, \in, \notin\} \\ [A' R A \{B\}] &:= \lambda I \lambda j (I j \wedge A' R A \{I_{B:j}\}) && R \in \{=, \in, \notin\} \\ [u | u R A \{\}] &:= \lambda I \lambda j \exists u \exists i (j = (u \cdot i) \wedge I i \wedge u R A \{I\}) && R \in \{=, \in, \notin\} \\ [u | u R A \{B\}] &:= \lambda I \lambda j \exists u \exists i (j = (u \cdot i) \wedge I i \wedge u R A \{I_{B:i}\}) && R \in \{=, \in, \notin\} \\ (\partial K) &:= \lambda I \lambda j (I j \wedge KI = I) && \text{(Beaver presupposition)} \\ \mathbf{0} &:= \lambda I \lambda j \neg(j = j) && \text{(absurd update)} \end{aligned}$$

3. English and Kalaallisut in CCG + UC2

• UNIVERSAL RULES

Forward application

$$(>) \quad X/Y: \beta_{ab} \quad Y: \alpha_a \quad \Rightarrow > \quad X: \beta\alpha$$

Backward application

$$(<) \quad Y: \alpha_a \quad X \setminus Y: \beta_{ab} \quad \Rightarrow < \quad X: \beta\alpha$$

Forward composition

$$(>\mathbf{B}) \quad X/Y: \beta_{ab} \quad Y/Z: \alpha_{ca} \quad \Rightarrow >\mathbf{B} \quad X/Z: \lambda v_c(\beta\alpha v)$$

Backward composition

$$(<\mathbf{B}) \quad YZ: \alpha_{ca} \quad X \setminus Y: \beta_{ab} \quad \Rightarrow <\mathbf{B} \quad X \setminus Z: \lambda v_c(\beta\alpha v)$$

• CATEGORY-TO-TYPE CORRESPONDENCE, **tp**:

$$\begin{aligned} \mathbf{tp}(s) &= [] \\ \mathbf{tp}(\underline{s}) &= [W] \end{aligned}$$

$$\begin{aligned} \mathbf{tp}(np) &= D \\ \mathbf{tp}(n) &= [WD] \\ \mathbf{tp}(vp) &= [DW] \end{aligned}$$

$$\begin{aligned} \mathbf{tp}(X/Y) &= \mathbf{tp}(Y)\mathbf{tp}(X) \\ \mathbf{tp}(X \setminus Y) &= \mathbf{tp}(Y)\mathbf{tp}(X) \\ \mathbf{tp}(X') &= \mathbf{tp}(X) \end{aligned}$$

- ENGLISH LEXICON (fragment 3)
 - i. English categories: AB-categories based on $\{s, \underline{s}, np, n, n'\}$.
E-abbreviations: $vp := \underline{s} \backslash np_N$
 $tv := vp/np_A$
 - ii. Eng. item | Category X: UC2-translation of type $tp(X)$
 - Lexical cat's: (pro)nouns, adjectives & verbs

he		$np_N: \tau \delta$
		$np_N: \perp \delta$
Ole		$\underline{s}/vp: \lambda \underline{P} \lambda \underline{w}([\underline{x} (\underline{x} = ole)^\circ]; \underline{P} \tau \delta \underline{w})$
		$vp \backslash (vp/np_A): \lambda \underline{R} \lambda \underline{x} \lambda \underline{w}([y (y = ole)^\circ]; \underline{R} \perp \delta \underline{x} \underline{w})$
dog		$n: \lambda \underline{w} \lambda \underline{x} [dog_w \underline{x}]$
sick		$n': \lambda \underline{w} \lambda \underline{x} [sick_w \underline{x}]$
		$n/n: \lambda \underline{N} \lambda \underline{w} \lambda \underline{x} ([sick_w \underline{x}]; \underline{N} \underline{w} \underline{x})$
buy.ASP		$vp/np_A: \lambda \underline{y} \lambda \underline{x} \lambda \underline{w} [buy_w \langle \underline{x}, \underline{y} \rangle]$
 - Grammatical cat's: determiners, auxiliaries, prosody

A		$n'/n: \lambda \underline{N} (\underline{N})$
		$\underline{s}/vp/n: \lambda \underline{N} \lambda \underline{P} \lambda \underline{w}([\underline{x}]; \underline{N} \underline{w} \tau \delta; \underline{P} \tau \delta \underline{w})$
		$vp \backslash (vp/np_A)/n: \lambda \underline{N} \lambda \underline{R} \lambda \underline{x} \lambda \underline{w}([y]; \underline{N} \underline{w} \perp \delta; \underline{R} \perp \delta \underline{x} \underline{w})$
HV.TNS		$vp/vp: \lambda \underline{P} (\underline{P})$
BE.TNS		$vp/n': \lambda \underline{N} \lambda \underline{x} \lambda \underline{w} (\underline{N} \underline{w} \underline{x})$
HV.TNS=N'T		$vp/vp: \lambda \underline{P} \lambda \underline{x} \lambda \underline{w} ([v v \in \underline{w} \{\}]; \underline{P} \underline{x} \perp \omega; [p p = \perp \omega \{\}]; [\underline{w} \notin \perp \Omega])$
BE.TNS=N'T		$vp/n': \lambda \underline{N} \lambda \underline{x} \lambda \underline{w} ([v v \in \underline{w} \{\}]; \underline{N} \perp \omega \underline{x}; [p p = \perp \omega \{\}]; [\underline{w} \notin \perp \Omega])$
.(prosody)		$\underline{s}/s: \lambda \underline{W} (\underline{W} \tau \omega; [\underline{p} \underline{p} = \tau \omega \{\}])$
 - KALAALLISUT LEXICON (fragment 1)
 - i. Kalaallisut categories: AB-categories based on $\{s, vp, np, n\}$
 - ii. Kal. item | Category X: UC2-translation of type $tp(X)$
 - Lexical cat's: roots & derivation

Ole-		$n: \lambda \underline{w} \lambda \underline{x} [(x = ole)^\circ]$
dog-		$n: \lambda \underline{w} \lambda \underline{x} [dog_w \underline{x}]$
sick-		$vp: \lambda \underline{x} \lambda \underline{w} [sick_w \underline{x}]$
-get		$vp \backslash n: \lambda \underline{N} \lambda \underline{x} \lambda \underline{w} ([y buy_w \langle \underline{x}, y \rangle]; \underline{N} \underline{w} \perp \delta)$
		$vp \backslash n: \lambda \underline{N} \lambda \underline{x} \lambda \underline{w} (\partial [\perp \delta \tau = \perp \delta \{\underline{w}\}]; [buy_w \langle \underline{x}, \perp \delta \rangle]; \underline{N} \underline{w} \perp \delta)$
-not		$vp \backslash vp: \lambda \underline{P} \lambda \underline{x} \lambda \underline{w} ([v v \in \underline{w} \{\}]; \underline{P} \underline{x} \perp \omega; [p p = \perp \omega \{\}]; [\underline{w} \notin \perp \Omega])$
		$vp \backslash vp: \lambda \underline{P} \lambda \underline{x} \lambda \underline{w} (\partial [\perp \omega \in \underline{w} \{\}]; \underline{P} \underline{x} \perp \omega; [p p = \perp \omega \{\}]; [\underline{w} \notin \perp \Omega])$
-n\vp		$n \backslash vp: \lambda \underline{P} \lambda \underline{w} \lambda \underline{x} (\underline{P} \underline{x} \underline{w})$
 - Grammatical cat's: inflection

-DEC.IV		$s \backslash vp/np: \lambda \underline{x} \lambda \underline{P} (\underline{P} \underline{x} \tau \omega; [\underline{p} \underline{p} = \tau \omega \{\}])$
-DEC.NG		$s \backslash vp/np: \lambda \underline{x} \lambda \underline{P} (\underline{P} \underline{x} \tau \omega^\perp; (\partial [\tau \omega \notin \perp \Omega]; [\underline{p} \underline{p} = \tau \omega \{\}]))$
-3SG		$np: \tau \delta$
-Ø		$s/s/n: \lambda \underline{N} \lambda K (([\underline{x}]; \underline{N} \tau \omega \tau \delta)^\top; K)$
		$s/s/n: \lambda \underline{N} \lambda K (([y]; \underline{N} \tau \omega \perp \delta)^\perp; K)$
-MOD		$s/s/n: \lambda \underline{N} \lambda K (([y]; \underline{N} \tau \omega \perp \delta; [X X = \perp \delta \{\tau \omega\}])^\perp; K)$
		$s/s/n: \lambda \underline{N} \lambda K (([w w \in \tau \omega \{\}]; [y]; \underline{N} \perp \omega \perp \delta; [X X = \perp \delta \{\perp \omega\}])^\perp; K)$

4. Compositional analyses in CCG + UC2

• ENGLISH

(1'E)	he	BE.TNS	sick	.
	np _N :	vp/n':	n':	s\':
	$\tau\delta$	$\lambda N \lambda x \lambda w (N w x)$	$\lambda w \lambda x [sick_w x]$	$\lambda W (W \tau\omega; [p] p = \tau\omega\{\})$
		$\underline{vp} (:= s \backslash np_N): \lambda x \lambda w [sick_w x]$		
	\underline{s} :	$\lambda w ([sick_w \tau\delta])$		
	$s: ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$			

(2'E)	Ole	BE.TNS	sick	.
	s/vp:	vp:	s\':	
	$\lambda P \lambda w ([x (x = ole)^\circ]; P \tau\delta w)$	$\lambda x \lambda w [sick_w x]$	$\lambda W (W \tau\omega; [p] p = \tau\omega)$	
	$\underline{s}: \lambda w ([x (x = ole)^\circ]; [sick_w \tau\delta])$			
	$s: ([x (x = ole)^\circ]; [sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$			

cat \rightarrow tp

$s \rightarrow []$
$\underline{s} \rightarrow [W]$
$np \rightarrow D$
$n \rightarrow [WD], n' \rightarrow [WD]$
$vp (:= s \backslash np_N) \rightarrow [DW]$
$tv (:= vp / np_A) \rightarrow [DDW]$
where
$W := s\omega \quad \underline{w} \in {}^+Var_w$
$D := s\delta \quad \underline{x} \in {}^+Var_D$
$[] := (st)st \quad K \in {}^+Var_{[]}$
$[ab] := a(b []) \quad \underline{W} \in {}^+Var_{[w]}$
$\underline{P} \in {}^+Var_{[DW]}$
$\underline{N} \in {}^+Var_{[wD]}$
\vdots

• KALAALLISUT

(1'K) **Morphology 1:** s-word with $\tau\delta$

sick-	-DEC.IV	-3SG
vp:	s \backslash vp / np:	np:
$\lambda x \lambda w [sick_w x]$	$\lambda x \lambda P (P x \tau\omega; [p] p = \tau\omega\{\})$	$\tau\delta$
	$\underline{s \backslash vp}: \lambda P (P \tau\delta \tau\omega; [p] p = \tau\omega\{\})$	
$s: ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$		

(2'K) **Morphology 2:** s-modifier (unmarked)

τ -rdg (sub): add $\tau\delta$ & look for comment

Ole-	\emptyset
------	-------------

n:	s/s \backslash n:
$\lambda w \lambda x [(x = ole)^\circ]$	$\lambda N \lambda K (([x]; N \tau\omega \tau\delta)^\tau; K)$
$\underline{s/s}: \lambda K (([x]; [(\tau\delta = ole)^\circ])^\tau; K)$	
$\lambda K ([x (x = ole)^\circ]^\tau; K)$	

Syntax

✓ topic-comment sequence ($A^\tau; B$)

(2''K) Ole- \emptyset sick-DEC.IV-3SG

$s/s: \lambda K ([x (x = ole)^\circ]^\tau; K)$	$s: ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$
$s: [x (x = ole)^\circ]^\tau; ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$	
$[x (x = ole)^\circ]; ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$	

\perp -rdg (obj): add $\perp\delta$ & look for elaboration

Ole-	\emptyset
n:	s/s \backslash n:
$\lambda w \lambda x [(x = ole)^\circ]$	$\lambda N \lambda K (([y]; N \tau\omega \perp\delta)^\perp; K)$
$\underline{s/s}: \lambda K (([y]; [(\perp\delta = ole)^\circ])^\perp; K)$	
$\lambda K ([y (y = ole)^\circ]^\perp; K)$	

* background-elaboration sequence ($A^\perp; B$)

Ole- \emptyset sick-DEC.IV-3SG

$s/s: \lambda K ([y (y = ole)^\circ]^\perp; K)$	$s: ([sick_{\tau\omega} \tau\delta]; \dots)$
$s: [y (y = ole)^\circ]^\perp; ([sick_{\tau\omega} \tau\delta]; [p] p = \tau\omega\{\})$	
$\mathbf{0} (:= \lambda I \lambda j \neg(j = j))$ (ABSURD UPDATE)	

• ENGLISH (3_E)(3'_E) **Syntax 1:** indefinite object

A	dog
vp\tv/n: $\lambda N \lambda R \lambda x \lambda w ([y]; \underline{N} \ w \ \perp \ \delta; \underline{R} \ \perp \ \delta \ x \ w)$	n: $\lambda w \lambda x [dog_w \ x]$
>	
vp\tv: $\lambda R \lambda x \lambda w ([y]; [dog_w \ \perp \ \delta]; \underline{R} \ \perp \ \delta \ x \ w)$	
$\lambda R \lambda x \lambda w ([y] \ dog_w \ y]; \underline{R} \ \perp \ \delta \ x \ w)$	

Syntax 2: transitive sentence

Ole	HV.TNS buy.ASP	A dog	.
\underline{s}/vp : $\lambda P \lambda w ([x (x = ole)^\circ]; \underline{P} \ \tau \ \delta \ w)$	vp/vp: tv (:= vp/np _A) $\lambda P(P) \ \lambda y \lambda x \lambda w [buy_w \langle x, y \rangle]$	vp\tv: $\lambda R \lambda x \lambda w ([y] \ dog_w \ y]; \underline{R} \ \perp \ \delta \ w)$	\underline{s}/s : $\lambda W (W \ \tau \ \omega; [p p = \tau \ \omega \{\}])$
> B			
tv: $\lambda y \lambda x \lambda w [buy_w \langle x, y \rangle]$			
>			
vp: $\lambda x \lambda w ([y] \ dog_w \ y]; [buy_w \langle x, \perp \ \delta \rangle])$			
$\lambda x \lambda w [y] \ dog_w \ y, buy_w \langle x, y \rangle]$			
<			
\underline{s} : $\lambda w ([x (x = ole)^\circ]; [y] \ dog_w \ y, buy_w \langle \tau \ \delta, y \rangle])$			
<			
s: ($[x (x = ole)^\circ]; [y] \ dog_{\tau \ \omega} \ y, buy_{\tau \ \omega} \langle \tau \ \delta, y \rangle]; [p p = \tau \ \omega \{\}])$			

KALAALLISUT (3_K)(3'_K) **Morphology:** n- + -vp\n(cf. English *theory* + *-ize* → *theorize*)

dog-	-get	-DEC.IV	-3SG
n: $\lambda w \lambda x [dog_w \ x]$	n\vp $\lambda N \lambda x \lambda w ([y] \ buy_w \langle x, y \rangle]; \underline{N} \ w \ \perp \ \delta)$	s\vp/np: $\lambda x \underline{P}(P \ x \ \tau \ \omega; [p p = \tau \ \omega \{\}])$	np: $\tau \ \delta$
<		>	
vp: $\lambda x \lambda w ([y] \ buy_w \langle x, y \rangle]; [dog_w \ \perp \ \delta])$		s\vp: $\lambda P(P \ \tau \ \delta \ \tau \ \omega; [p p = \tau \ \omega \{\}])$	
$\lambda x \lambda w [y] \ buy_w \langle x, y \rangle, dog_w \ y]$		<	
s: ($[y] \ buy_{\tau \ \omega} \langle \tau \ \delta, y \rangle, dog_{\tau \ \omega} \ y]; [p p = \tau \ \omega \{\}])$			

Syntax

Ole-Ø	dog-get-DEC.IV-3SG
s/s: $\lambda K ([x (x = ole)^\circ]^\tau; K)$	s: ($[y] \ buy_{\tau \ \omega} \langle \tau \ \omega, y \rangle, dog_{\tau \ \omega} \ y]; [p p = \tau \ \omega \{\}])$
>	
s: ($[x (x = ole)^\circ]^\tau; ([y] \ buy_{\tau \ \omega} \langle \tau \ \delta, y \rangle, dog_{\tau \ \omega} \ y]; [p p = \tau \ \omega \{\}])$)	
$\underbrace{\hspace{10em}}_{\text{-----}}$ $([x (x = ole)^\circ]; [y] \ buy_{\tau \ \omega} \langle \tau \ \delta, y \rangle, dog_{\tau \ \omega} \ y]; [p p = \tau \ \omega \{\}])$	

Updates for (3_E) ≡ (3_K)

1	$*p_0[[x (x = ole)]]^g$	
=	$\lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle \rangle \rangle w \in {}^0 p_0 \}$	=: c ₁
2	$c_1[[y] \ dog_{\tau \ \omega} \ y, buy_{\tau \ \omega} \langle \tau \ \delta, y \rangle]]^g$	
=	$\lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle d \rangle \rangle w \in {}^0 p_0 \ \& \ d \in {}^0 [[dog]](w) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](w) \}$	=: c ₂
3	$c_2[[p p = \tau \ \omega \{\}]]^g$	
=	$\lambda \{ \langle \langle p_1, [ole], w, p_0 \rangle, \langle d \rangle \rangle w \in {}^0 p_0 \ \& \ d \in {}^0 [[dog]](w) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](w) \ \& \ p_1 = \lambda \{ v \in {}^0 p_0 \exists d' : d' \in {}^0 [[dog]](v) \ \& \ \langle [ole], d' \rangle \in {}^0 [[buy]](v) \} \}$	=: c ₃

Example: Let $p_0 = \{w_0, w_1, w_2\}$,

and $\{d \mid d \in \text{dom}[\text{dog}] \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}]\} = \{\}$
 $\{d \mid d \in \text{dom}[\text{dog}] \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}]\} = \{d_1\}$
 $\{d \mid d \in \text{dom}[\text{dog}] \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}]\} = \{d_1, d_2\}$

Then:

$*p_0[[[x \mid (x = ole)^\circ]]]^g$ $c_1[[[y \mid \text{dog}_{\tau\omega} y, \text{buy}_{\tau\omega}(\tau\delta, y)]]]^g$ $c_2[[[p \mid p = \tau\omega\{\}\]]]^g$
 $\{ \langle \langle \text{ole}, w_0, p_0, \langle \rangle \rangle, \langle \langle \text{ole}, w_1, p_0, \langle \rangle \rangle, \langle \langle \text{ole}, w_2, p_0, \langle \rangle \rangle \rangle \}$
 $\{ \langle \langle \text{ole}, w_1, p_0, \langle d_1 \rangle \rangle, \langle \langle \text{ole}, w_2, p_0, \langle d_1 \rangle \rangle, \langle \langle \text{ole}, w_2, p_0, \langle d_2 \rangle \rangle \rangle \}$
 $\{ \langle \langle \{w_1, w_2\}, \text{ole}, w_1, p_0, \langle d_1 \rangle \rangle, \langle \langle \{w_1, w_2\}, \text{ole}, w_2, p_0, \langle d_1 \rangle \rangle, \langle \langle \{w_1, w_2\}, \text{ole}, w_2, p_0, \langle d_2 \rangle \rangle \rangle \}$

ENGLISH (4_E)

(4'_E) **Syntax 1:** n/n + n

A	<i>sick</i>	<i>dog</i>
$vp \setminus tv / n: \lambda N \lambda R \lambda x \lambda w ([y; N w \perp \delta; R \perp \delta x w])$	$n / n: \lambda N \lambda w \lambda x ([sick_w x]; N w x)$	$n: \lambda w \lambda x [dog_w x]$
	\longrightarrow	
	$n: \lambda w \lambda x [sick_w x, dog_w x]$	
	\longrightarrow	
$vp \setminus tv: \lambda R \lambda x \lambda w ([y \mid sick_w y, dog_w y]; R \perp \delta x w)$		

Syntax 2: sentence

Ole	HV.TNS buy.ASP	A <i>sick dog</i>	.
	\longrightarrow B		
$\underline{s} / vp:$ $\lambda P \lambda w ([x \mid (x = ole)^\circ]; P \tau \delta w)$	$tv:$ $\lambda y \lambda x \lambda w [buy_w(x, y)]$	$vp \setminus tv:$ $\lambda R \lambda x \lambda w ([y \mid sick_w y, dog_w y]; R \perp \delta w)$	$\underline{s} / s:$ $\lambda W (W \tau \omega; [p \mid p = \tau \omega \{\}])$
	\longrightarrow		
	$vp: \lambda x \lambda w [y \mid sick_w y, dog_w y, buy_w(x, y)]$		
	\longleftarrow		
$\underline{s}: \lambda w ([x \mid (x = ole)^\circ]; [y \mid sick_w y, dog_w y, buy_w(\tau\delta, y)])$			
	\longleftarrow		
$s: ([x \mid (x = ole)^\circ]; [y \mid sick_{\tau\omega} y, dog_{\tau\omega} y, buy_{\tau\omega}(\tau\delta, y)]; [p \mid p = \tau\omega\{\}])$			

Updates

1 $*p_0[[[x \mid (x = ole)^\circ]]]^g$
 $= \{ \langle \langle \text{ole}, w, p_0, \langle \rangle \rangle \mid w \in \text{dom}[\text{ole}] \rangle \} \quad =: c_1$

2 $c_1[[[y \mid sick_{\tau\omega} y, dog_{\tau\omega} y, buy_{\tau\omega}(\tau\delta, y)]]]^g$
 $= \{ \langle \langle \langle \text{ole}, w, p_0, \langle d \rangle \rangle \mid w \in \text{dom}[\text{ole}] \& d \in \text{dom}[\text{sick}](w) \& d \in \text{dom}[\text{dog}](w) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w) \rangle \} \quad =: c_2$

3 $c_2[[[p \mid p = \tau\omega\{\}]]]^g$
 $= \{ \langle \langle \langle p_1, \text{ole}, w, p_0, \langle d \rangle \rangle \mid w \in \text{dom}[\text{ole}] \& d \in \text{dom}[\text{sick}](w) \& d \in \text{dom}[\text{dog}](w) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w) \& p_1 = \{v \in \text{dom}[\text{buy}]\} \exists d': \& d' \in \text{dom}[\text{sick}](v) \& d' \in \text{dom}[\text{dog}](v) \& \langle \text{ole}, d' \rangle \in \text{dom}[\text{buy}](v) \rangle \rangle \} \quad =: c_3$

Example: Let $p_0 = \{w_0, w_1, w_2, w_3\}$

and $\{d \mid d \in \text{dom}[\text{sick}](w_0) \& d \in \text{dom}[\text{dog}](w_0) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w_0)\} = \{\}$
 $\{d \mid d \in \text{dom}[\text{sick}](w_1) \& d \in \text{dom}[\text{dog}](w_1) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w_1)\} = \{\}$
 $\{d \mid d \in \text{dom}[\text{sick}](w_2) \& d \in \text{dom}[\text{dog}](w_2) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w_2)\} = \{d_2\}$
 $\{d \mid d \in \text{dom}[\text{sick}](w_3) \& d \in \text{dom}[\text{dog}](w_3) \& \langle \text{ole}, d \rangle \in \text{dom}[\text{buy}](w_3)\} = \{d_1\}$

Then:

$*p_0[[[x \mid (x = ole)^\circ]]]^g$ $c_1[[[y \mid sick_{\tau\omega} y, dog_{\tau\omega} y, buy_{\tau\omega}(\tau\delta, y)]]]^g$ $c_2[[[p \mid p = \tau\omega\{\}]]]^g$
 $\{ \langle \langle \text{ole}, w_0, p_0, \langle \rangle \rangle, \langle \langle \text{ole}, w_1, p_0, \langle \rangle \rangle, \langle \langle \text{ole}, w_2, p_0, \langle \rangle \rangle, \langle \langle \text{ole}, w_3, p_0, \langle \rangle \rangle \rangle \}$
 $\{ \langle \langle \text{ole}, w_2, p_0, \langle d_2 \rangle \rangle, \langle \langle \text{ole}, w_3, p_0, \langle d_1 \rangle \rangle \rangle \}$
 $\{ \langle \langle \{w_2, w_3\}, \text{ole}, w_2, p_0, \langle d_2 \rangle \rangle, \langle \langle \{w_2, w_3\}, \text{ole}, w_3, p_0, \langle d_1 \rangle \rangle \rangle \}$

KALAALLISUT (4_K)(4'_K) **Morphology 1.** MOD(ifier) s/s

sick-	-n\vp	-MOD
vp: $\lambda x \lambda w [sick_w x]$	n\vp $\lambda P \lambda w \lambda x (P x w)$	s/s\n: $\lambda N \lambda K ([y]; \underline{N} \tau \omega \perp \delta; [X] X = \perp \delta \{\tau \omega\})^\perp; K)$
<		
n: $\lambda w \lambda x ([sick_w x])$		
<		
s/s: $\lambda K (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; K)$		

Morphology 2. s with anaphoric -get

dog-	-get	-DEC.IV	-3SG
n:	vp\n	s\vp/np:	np:
$\lambda w \lambda x [dog_w x]$	$\lambda N \lambda x \lambda w (\partial [\perp \delta t = \perp \delta \{\tau \omega\}]; [buy_w \langle x, \perp \delta \rangle]; \underline{N} w \perp \delta)$	$\lambda x P (P x \tau \omega; [p] p = \tau \omega \{\})$	$\tau \delta$
<		>	
vp: $\lambda x \lambda w (\partial [\perp \delta t = \perp \delta \{\tau \omega\}]; [buy_w \langle x, \perp \delta \rangle], dog_w \perp \delta)$		s\vp: $\lambda P (P \tau \delta \tau \omega; [p] p = \tau \omega \{\})$	
<			
s: $(\partial [\perp \delta t = \perp \delta \{\tau \omega\}]; [buy_{\tau \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\tau \omega} \perp \delta; [p] p = \tau \omega \{\})$			

✓ **Syntax** w/anaphoric -get

Ole	sick-n\vp-MOD	dog-get-DEC.IV-3SG
<		<
s/s: $\lambda K ([x] (x = ole)^\circ)^\top; K)$	s/s: $\lambda K (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; K)$	s: $\partial [\perp \delta t = \perp \delta \{\tau \omega\}]; \dots [p] p = \tau \omega \{\}$
>B		>
s/s: $\lambda K (([x] (x = ole)^\circ)^\top; (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; K)$		
>		
s: $([x] (x = ole)^\circ)^\top; (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; (\partial [\perp \delta t = \perp \delta \{\tau \omega\}]; [buy_{\tau \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\tau \omega} \perp \delta; [p] p = \tau \omega \{\})$		
----- ----- -----		
([x] (x = ole) ^o ; [y] sick _{τω} y; [X] X = ⊥δ{τω} }; ∂[⊥δt = ⊥δ{τω}]; [buy _{τω} ⟨τδ, ⊥δ⟩], dog _{τω} ⊥δ; [p] p = τω{ })		
([x] (x = ole) ^o ; [y] sick _{τω} y; [X] X = ⊥δ{τω} }; [buy _{τω} ⟨τδ, ⊥δ⟩], dog _{τω} ⊥δ; [p] p = τω{ })		

cf. * **Syntax** w/non-anaphoric -get

Ole	sick-n\vp-MOD	dog-get-DEC.IV-3SG
<		<
s/s: $\lambda K ([x] (x = ole)^\circ)^\top; K)$	s/s: $\lambda K (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; K)$	s: $[y] buy_{\tau \omega} \langle \tau \delta, y \rangle, \dots; [p] p = \tau \omega \{\}$
>B		>
s/s: $\lambda K (([x] (x = ole)^\circ)^\top; (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; K)$		
>		
s: $([x] (x = ole)^\circ)^\top; (([y] sick_{\tau \omega} y); [X] X = \perp \delta \{\tau \omega\})^\perp; ([y] buy_{\tau \omega} \langle \tau \delta, y \rangle, dog_{\tau \omega} y; [p] p = \tau \omega \{\})$		
----- ----- -----		
0 (= λλj ¬(j = j))		

• **Updates**

1	$*p_0 [[x] (x = ole)]^g$	
=	$\lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle \rangle \rangle \mid w \in {}^i p_0 \}$	=: c ₁
2	$c_1 [[y] sick_{\tau \omega} y]^g$	
=	$\lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle d \rangle \rangle \mid w \in {}^i p_0 \ \& \ d \in {}^i [sick](w) \}$	=: c ₂
3	$c_2 [[X] X = \perp \delta \{\tau \omega\}]^g$	
=	$\lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle x, d \rangle \rangle \mid w \in {}^i p_0 \ \& \ d \in {}^i [sick](w) \ \& \ x = \lambda \{ d' \mid d' \in {}^i [sick](w) \} \}$	=: c ₃
∂3	$c_3 [\partial [\perp \delta t = \perp \delta \{\tau \omega\}]]^g$	
=	$\lambda \{ j \in c_3 \mid c_3 [[\perp \delta t = \perp \delta \{\tau \omega\}]]^g = c_3 \}$	= c ₃

$$\begin{aligned}
&^4 \quad c_3[[buy_{\tau\omega}\langle\tau\delta, \perp\delta\rangle, dog_{\tau\omega}\perp\delta]]^g \\
= &^x\{\langle\langle[ole], w, p_0\rangle, \langle x, d\rangle\rangle \mid w \in {}^{\text{O}}p_0 \ \& \ d \in {}^{\text{O}}[[sick]](w) \ \& \ x = {}^x\{d' \mid d' \in {}^{\text{O}}[[sick]](w)\} \\
&\quad \& \ d \in {}^{\text{O}}[[dog]](w) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w)\} \quad =: c_4 \\
&^5 \quad c_4[[p \mid p = \tau\omega\{\}\]]^g \\
= &^x\{\langle\langle p_1, [ole], w, p_0\rangle, \langle x, d\rangle\rangle \mid w \in {}^{\text{O}}p_0 \ \& \ d \in {}^{\text{O}}[[sick]](w) \ \& \ x = {}^x\{d' \mid d' \in {}^{\text{O}}[[sick]](w)\} \\
&\quad \& \ d \in {}^{\text{O}}[[dog]](w) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w) \\
&\quad \& \ p_1 = {}^x\{v \in {}^{\text{O}}p_0 \mid \exists d': d' \in {}^{\text{O}}[[sick]](v) \ \& \ d' \in {}^{\text{O}}[[dog]](v) \ \& \ \langle[ole], d'\rangle \in {}^{\text{O}}[[buy]](v)\}\} \quad =: c_5
\end{aligned}$$

Example: Suppose $p_0 = {}^x\{w_0, w_1, w_2, w_3\}$

and $[ole] = d_0$

$$\begin{aligned}
&\{d \mid d \in {}^{\text{O}}[[sick]](w_0)\} = \{\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_1)\} = \{d_2\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_2)\} = \{d_1, d_2\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_3)\} = \{d_0, d_1\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_0) \ \& \ d \in {}^{\text{O}}[[dog]](w_0) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w_0)\} = \{\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_1) \ \& \ d \in {}^{\text{O}}[[dog]](w_1) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w_1)\} = \{\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_2) \ \& \ d \in {}^{\text{O}}[[dog]](w_2) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w_2)\} = \{d_2\} \\
&\{d \mid d \in {}^{\text{O}}[[sick]](w_3) \ \& \ d \in {}^{\text{O}}[[dog]](w_3) \ \& \ \langle[ole], d\rangle \in {}^{\text{O}}[[buy]](w_3)\} = \{d_1\}
\end{aligned}$$

then

$$\begin{array}{ccc}
{}^*p_0[[[x \mid (x = ole)]]]^g & c_1[[y \mid sick_{\tau\omega}y]]^g & c_2[[X \mid X = \perp\delta\{\}_{\tau\omega}\{\}]]^g \\
{}^x\{\langle\langle d_0, w_0, p_0\rangle, \langle \ \rangle\rangle, \\
\quad \langle\langle d_0, w_1, p_0\rangle, \langle \ \rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle \ \rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle \ \rangle\rangle\} & {}^x\{\langle\langle d_0, w_1, p_0\rangle, \langle d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle d_1\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle d_0\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle d_1\rangle\rangle\} & {}^x\{\langle\langle d_0, w_1, p_0\rangle, \langle {}^x\{d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_1\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_0\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_1\rangle\rangle\} \\
c_3[[\partial[\perp\delta\tau = \perp\delta\{\}_{\tau\omega}\{\}]]]^g & c_3[[buy_{\tau\omega}\langle\tau\delta, \perp\delta\rangle, dog_{\tau\omega}\perp\delta]]^g & c_4[[[p \mid p = \tau\omega\{\}]]]^g \\
{}^x\{\langle\langle d_0, w_1, p_0\rangle, \langle {}^x\{d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_1\rangle\rangle, \\
\quad \langle\langle d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_0\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_1\rangle\rangle\} & \langle\langle d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_1\rangle\rangle\} & {}^x\{\langle\langle {}^x\{w_2, w_3\}, d_0, w_2, p_0\rangle, \langle {}^x\{d_1, d_2\}, d_2\rangle\rangle, \\
\quad \langle\langle {}^x\{w_2, w_3\}, d_0, w_3, p_0\rangle, \langle {}^x\{d_0, d_1\}, d_1\rangle\rangle\}
\end{array}$$

Homework 7

Complete the following analysis of Kalaallisut (5_K).

NOTE: Given the lexical ambiguity of ‘-MOD’, ‘-get’, and ‘-not’ in Kalaallisut fragment 1, there are seven other analyses to consider. For now, don't worry about other lexical entries (unless you want to)—we'll come back to them when we discuss the solution to this homework.

• KALAALLISUT (5_K)

(5_K) Ole naparsimasumik qimmisinngilaq.
 Ole naparsima-tuq-mik qimmiq-si-nngit-la-q
 Ole sick-n\vp-MOD dog-get-not-DEC.NG-3SG

(5'_K) **Morphology 1:** (MOD)ifier s for negated s-word

sick- -n\vp -MOD
 ⋮

Morphology 2: negated vp-base

dog- -get -not -DEC.NG -3SG
 ⋮

✓ **Syntax**

Ole sick-n\vp-MOD dog-get-not-DEC.NG-3SG
 ⋮

s:

$([x | (x = ole)^o]^T; (([v | v \in \tau\omega\{\}]; [y | sick_{\perp\omega} y]; [X | X = \perp\delta\{\}_{\perp\omega}])^{\perp}; (\partial[\perp\omega \in \tau\omega\{\}]; \partial[\perp\delta = \perp\delta\{\}_{\perp\omega}]); [buy_{\perp\omega} \langle \tau\delta, \perp\delta \rangle,$

$dog_{\perp\omega} \perp\delta]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}]))$

$([x | (x = ole)^o]; [v | v \in \tau\omega\{\}]; [y | sick_{\perp\omega} y]; [X | X = \perp\delta\{\}_{\perp\omega}]; \partial[\perp\omega \in \tau\omega\{\}]; \partial[\perp\delta = \perp\delta\{\}_{\perp\omega}]; [buy_{\perp\omega} \langle \tau\delta, \perp\delta \rangle,$
 $dog_{\perp\omega} \perp\delta]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}])$

Updates

¹ * $p_0[[x | (x = ole)]]^g$
 = $\lambda\langle\langle[[ole]], w, p_0\rangle, \langle\rangle\rangle | w \in {}^0p_0$ =: c₁
 ⋮
⁸ $c_7[[p | p = \tau\omega\{\}]]^g$
 = $\lambda\langle\langle p_2, [[ole]], w, p_0\rangle, \langle p_1, X, d, v \rangle \rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\}$
 $\ \& \ \langle [[ole]], d \rangle \in {}^0[[buy]](v) \ \& \ d \in {}^0[[dog]](v)$
 $\ \& \ p_1 = \lambda\{u \in {}^0p_0 | \exists d': d' \in {}^0[[sick]](u) \ \& \ d' \in {}^0[[dog]](u) \ \& \ \langle [[ole]], d' \rangle \in {}^0[[buy]](u)\}$
 $\ \& \ w \notin {}^0p_1 \ \& \ p_2 = \lambda\{w' \in {}^0p_0 | w' \notin {}^0p_1\}$ =: c₈

Example: Let $p_0 = \lambda\{w_0, w_1, w_2, w_3\}$,

$[[ole]] = d_0$

$\{d | d \in {}^0[[sick]](w_0)\} = \{\}$

$\{d | d \in {}^0[[sick]](w_1)\} = \{d_2\}$

$\{d | d \in {}^0[[sick]](w_2)\} = \{d_1, d_2\}$

$\{d | d \in {}^0[[sick]](w_3)\} = \{d_0, d_1\}$

$\{d | d \in {}^0[[sick]](w_0) \ \& \ d \in {}^0[[dog]](w_0) \ \& \ \langle [[ole]], d \rangle \in {}^0[[buy]](w_0)\} = \{\}$

$\{d | d \in {}^0[[sick]](w_1) \ \& \ d \in {}^0[[dog]](w_1) \ \& \ \langle [[ole]], d \rangle \in {}^0[[buy]](w_1)\} = \{\}$

$\{d | d \in {}^0[[sick]](w_2) \ \& \ d \in {}^0[[dog]](w_2) \ \& \ \langle [[ole]], d \rangle \in {}^0[[buy]](w_2)\} = \{d_2\}$

$\{d | d \in {}^0[[sick]](w_3) \ \& \ d \in {}^0[[dog]](w_3) \ \& \ \langle [[ole]], d \rangle \in {}^0[[buy]](w_3)\} = \{d_1\}$

Then:

⋮

Solution to homework 7 (and related issues)

KALAALLISUT (5_K)(5'_K) **Morphology 1:** (MOD)ifier s for negated s-word

sick-	-n\vp	-MOD
vp: $\lambda x \lambda w [sick_{w,x}]$	n\vp $\lambda P \lambda w \lambda x (P x w)$	s/s\n: $\lambda N \lambda K ([v] v \in \tau \omega \{\}; [y]; N \perp \omega \perp \delta; [X] X = \perp \delta \{\perp \omega\})^{\perp}; K)$
n: $\lambda w \lambda x ([sick_{w,x}]$		
s/s: $\lambda K ([v] v \in \tau \omega \{\}; [y]; [sick_{\perp \omega} \perp \delta]; [X] X = \perp \delta \{\perp \omega\})^{\perp}; K)$ $\lambda K ([v] v \in \tau \omega \{\}; [y] sick_{\perp \omega} y]; [X] X = \perp \delta \{\perp \omega\})^{\perp}; K)$		

Morphology 2a: negated vp-base

dog-get-	-not
vp: $\lambda x \lambda w (\partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{w,x}, \perp \delta], dog_{w,x} \perp \delta)$	vp\vp $\lambda P \lambda x \lambda w (\partial[\perp \omega \in w \{\}]; P x \perp \omega; [p] p = \perp \omega \{\}; [w] w \notin \perp \Omega)$
vp: $\lambda x \lambda w (\partial[\perp \omega \in w \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle x, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [w] w \notin \perp \Omega)$	

Morphology 2b: negated s-word

dog-get-not-	-DEC.NG	-3SG
vp: $\lambda x \lambda w (\partial[\perp \omega \in w \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle x, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [w] w \notin \perp \Omega)$	s\vp/np: $\lambda x P (P x \tau \omega^{\perp}; (\partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\})) \tau \delta$	np: $\tau \delta$
s\vp: $\lambda P (P \tau \delta \tau \omega^{\perp}; (\partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\}))$		
s: $(\partial[\perp \omega \in \tau \omega \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [\tau \omega \notin \perp \Omega]; \partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\})$		

✓ **Syntax**

Ole sick-n\vp-MOD	>B	dog-get-not-DEC.NG-3SG
s/s: $\lambda K ([x] (x = ole)^{\circ \top}; (([v] v \in \tau \omega \{\}); [y] sick_{\perp \omega} y]; [X] X = \perp \delta \{\perp \omega\})^{\perp}; K)$		s: $(\partial[\perp \omega \in \tau \omega \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [\tau \omega \notin \perp \Omega]; \partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\})$
s: $([x] (x = ole)^{\circ \top}; (([v] v \in \tau \omega \{\}); [y] sick_{\perp \omega} y]; [X] X = \perp \delta \{\perp \omega\})^{\perp}; (\partial[\perp \omega \in \tau \omega \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [\tau \omega \notin \perp \Omega]; \partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\})))$		
$([x] (x = ole)^{\circ \top}; [v] v \in \tau \omega \{\}; [y] sick_{\perp \omega} y; [X] X = \perp \delta \{\perp \omega\}; \partial[\perp \omega \in \tau \omega \{\}]; \partial[\perp \delta t = \perp \delta \{\perp \omega\}]; [buy_{\perp \omega} \langle \tau \delta, \perp \delta \rangle], dog_{\perp \omega} \perp \delta; [p] p = \perp \omega \{\}; [\tau \omega \notin \perp \Omega]; \partial[\tau \omega \notin \perp \Omega]; [p] p = \tau \omega \{\})$		

Updates

$$\begin{aligned}
1 \quad & *_{p_0}[[x | (x = ole)]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle\rangle\rangle | w \in {}^0p_0\} \quad =: c_1 \\
2 \quad & c_1[[v | v \in \tau\omega\{\}]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle v\rangle\rangle | w, v \in {}^0p_0\} \quad =: c_2 \\
3 \quad & c_2[[y | sick_{\perp\omega}y]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v)\} \quad =: c_3 \\
4 \quad & c_3[[X | X = \perp\delta\{\perp\omega\}]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\}\} \quad =: c_4 \\
\partial 2 \quad & c_4[[\perp\omega \in \tau\omega\{\}]]^g \\
= & \lambda\{j \in c_4 | ((\perp j)_{\omega})_1 \in \{((\tau i)_{\omega})_1 | i \in {}^0c_4\}\} \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ v \in \{w' | \exists i \in {}^0c_4: w' = ((\tau i)_{\omega})_1\}\} \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \ \& \ v \in {}^0p_0\} \quad = c_4 \\
& c_4[[\partial[\perp\omega \in \tau\omega\{\}]]]^g \\
= & \lambda\{j \in c_4 | c_4[[\perp\omega \in \tau\omega\{\}]]^g = c_4\} \\
= & \lambda\{j \in c_4 | c_4 = c_4\} \quad = c_4 \\
\partial 4 \quad & c_4[[\perp\delta\tau = \perp\delta\{\perp\omega\}]]^g \\
= & \lambda\{j \in c_4 | ((\perp j)_{\delta\tau})_1 = \lambda\{((\perp i)_{\delta})_1 | i \in {}^0c_4 \ \& \ ((\perp i)_{\omega})_1 = ((\perp j)_{\omega})_1\}\} \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ X = \lambda\{d'' | \exists i \in {}^0c_4: ((\perp i)_{\omega})_1 = v \ \& \ d'' = ((\perp i)_{\delta})_1\}\} \quad = c_4 \\
& c_4[[\partial[\perp\delta\tau = \perp\delta\{\perp\omega\}]]]^g \\
= & \lambda\{j \in c_4 | c_4[[\perp\delta\tau = \perp\delta\{\perp\omega\}]]^g = c_4\} \\
= & \lambda\{j \in c_4 | c_4 = c_4\} \quad = c_4 \\
5 \quad & c_4[[buy_{\perp\omega}(\tau\delta, \perp\delta), dog_{\perp\omega} \perp\delta]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ \langle[ole], d\rangle \in {}^0[[buy]](v) \ \& \ d \in {}^0[[dog]](v)\} \quad =: c_5 \\
6 \quad & c_5[[p | p = \perp\omega\{\}]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle p_1, X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ \langle[ole], d\rangle \in {}^0[[buy]](v) \ \& \ d \in {}^0[[dog]](v) \\
& \ \& \ p_1 = \lambda\{u \in {}^0p_0 | \exists d': d' \in {}^0[[sick]](u) \ \& \ d' \in {}^0[[dog]](u) \ \& \ \langle[ole], d'\rangle \in {}^0[[buy]](u)\}\} \quad =: c_6 \\
7 \quad & c_6[[\tau\omega \notin \perp\Omega]]^g \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle p_1, X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ \langle[ole], d\rangle \in {}^0[[buy]](v) \ \& \ d \in {}^0[[dog]](v) \\
& \ \& \ p_1 = \lambda\{u \in {}^0p_0 | \exists d': d' \in {}^0[[sick]](u) \ \& \ d' \in {}^0[[dog]](u) \ \& \ \langle[ole], d'\rangle \in {}^0[[buy]](u)\} \\
& \ \& \ w \notin {}^0p_1\} \quad =: c_7 \\
\partial 7 \quad & c_7[[\tau\omega \notin \perp\Omega]]^g \\
= & \lambda\{j \in c_7 | ((\tau j)_{\omega})_1 \notin {}^0((\perp j)_{\Omega})_1\} \\
= & \lambda\{\langle\langle[ole], w, p_0\rangle, \langle p_1, X, d, v\rangle\rangle | w, v \in {}^0p_0 \ \& \ d \in {}^0[[sick]](v) \ \& \ X = \lambda\{d' | d' \in {}^0[[sick]](v)\} \\
& \ \& \ \langle[ole], d\rangle \in {}^0[[buy]](v) \ \& \ d \in {}^0[[dog]](v) \\
& \ \& \ p_1 = \lambda\{u \in {}^0p_0 | \exists d': d' \in {}^0[[sick]](u) \ \& \ d' \in {}^0[[dog]](u) \ \& \ \langle[ole], d'\rangle \in {}^0[[buy]](u)\} \\
& \ \& \ w \notin {}^0p_1 \ \& \ w \notin {}^0p_1\} \quad = c_7 \\
& c_7[[\partial[\tau\omega \notin \perp\Omega]]]^g \\
= & \lambda\{j \in c_7 | c_7[[\tau\omega \notin \perp\Omega]]^g = c_7\} \quad = c_7
\end{aligned}$$

$$\begin{aligned}
& c_7[[\mathbf{p} \mid \mathbf{p} = \tau\omega\{\}\]]^g \\
= & \lambda\{\langle\langle p_2, [\mathit{ole}], w, p_0 \rangle, \langle p_1, x, d, v \rangle \rangle \mid w, v \in \mathbb{U}_{p_0} \ \& \ d \in \mathbb{U}[[\mathit{sick}]](v) \ \& \ x = \lambda\{d' \mid d' \in \mathbb{U}[[\mathit{sick}]](v)\} \\
& \ \& \ \langle [\mathit{ole}], d \rangle \in \mathbb{U}[[\mathit{buy}]](v) \ \& \ d \in \mathbb{U}[[\mathit{dog}]](v) \\
& \ \& \ p_1 = \lambda\{u \in \mathbb{U}_{p_0} \mid \exists d': d' \in \mathbb{U}[[\mathit{sick}]](u) \ \& \ d' \in \mathbb{U}[[\mathit{dog}]](u) \ \& \ \langle [\mathit{ole}], d' \rangle \in \mathbb{U}[[\mathit{buy}]](u)\} \\
& \ \& \ w \notin \mathbb{U}_{p_1} \ \& \ p_2 = \lambda\{w' \in \mathbb{U}_{p_0} \mid w' \notin \mathbb{U}_{p_1}\} \qquad \qquad \qquad =: c_8
\end{aligned}$$

Example: Let $p_0 = \lambda\{w_0, w_1, w_2, w_3\}$,

$$[[\mathit{ole}]] = d_0$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_0)\} = \{\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_1)\} = \{d_2\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_2)\} = \{d_1, d_2\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_3)\} = \{d_0, d_1\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_0) \ \& \ d \in \mathbb{U}[[\mathit{dog}]](w_0) \ \& \ \langle [\mathit{ole}], d \rangle \in \mathbb{U}[[\mathit{buy}]](w_0)\} = \{\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_1) \ \& \ d \in \mathbb{U}[[\mathit{dog}]](w_1) \ \& \ \langle [\mathit{ole}], d \rangle \in \mathbb{U}[[\mathit{buy}]](w_1)\} = \{\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_2) \ \& \ d \in \mathbb{U}[[\mathit{dog}]](w_2) \ \& \ \langle [\mathit{ole}], d \rangle \in \mathbb{U}[[\mathit{buy}]](w_2)\} = \{d_2\}$$

$$\{d \mid d \in \mathbb{U}[[\mathit{sick}]](w_3) \ \& \ d \in \mathbb{U}[[\mathit{dog}]](w_3) \ \& \ \langle [\mathit{ole}], d \rangle \in \mathbb{U}[[\mathit{buy}]](w_3)\} = \{d_1\}$$

Then:

$$\begin{array}{cccc}
*_{p_0}[[\mathbf{x} \mid (\mathbf{x} = \mathit{ole})^o]]^g & c_1[[\mathbf{v} \mid v \in \tau\omega\{\}\]]^g & c_2[[\mathbf{y} \mid \mathit{sick}_{\perp\omega} \mathbf{y}]]^g & c_3[[\mathbf{X} \mid X = \perp\delta\{\perp\omega\}\]]^g \\
\lambda\{\langle\langle d_0, w_0, p_0 \rangle, \langle \ \ \rangle\ \rangle, & \lambda\{\langle\langle d_0, w_0, p_0 \rangle, \langle w_0 \rangle\ \rangle, & \lambda\{\langle\langle d_0, w_0, p_0 \rangle, \langle d_2, w_1 \rangle\ \rangle, & \lambda\{\langle\langle d_0, w_0, p_0 \rangle, \langle \lambda\{d_2\}, \ d_2, w_1 \rangle\ \rangle, \\
& \langle\langle d_0, w_0, p_0 \rangle, \langle w_1 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle d_1, w_2 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_1, w_2 \rangle\ \rangle, \\
& \langle\langle d_0, w_0, p_0 \rangle, \langle w_2 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle d_2, w_2 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_2, w_2 \rangle\ \rangle, \\
& \langle\langle d_0, w_0, p_0 \rangle, \langle w_3 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle d_0, w_3 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_0, w_3 \rangle\ \rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle \ \ \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle w_0 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle d_1, w_3 \rangle\ \rangle, & \langle\langle d_0, w_0, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_1, w_3 \rangle\ \rangle, \\
& \langle\langle d_0, w_1, p_0 \rangle, \langle w_1 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle d_2, w_1 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle \lambda\{d_2\}, \ d_2, w_1 \rangle\ \rangle, \\
& \langle\langle d_0, w_1, p_0 \rangle, \langle w_2 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle d_1, w_2 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_1, w_2 \rangle\ \rangle, \\
& \langle\langle d_0, w_1, p_0 \rangle, \langle w_3 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle d_2, w_2 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_2, w_2 \rangle\ \rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle \ \ \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle w_0 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle d_0, w_3 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_0, w_3 \rangle\ \rangle, \\
& \langle\langle d_0, w_2, p_0 \rangle, \langle w_1 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle d_1, w_3 \rangle\ \rangle, & \langle\langle d_0, w_1, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_1, w_3 \rangle\ \rangle, \\
& \langle\langle d_0, w_2, p_0 \rangle, \langle w_2 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle d_2, w_1 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle \lambda\{d_2\}, \ d_2, w_1 \rangle\ \rangle, \\
& \langle\langle d_0, w_2, p_0 \rangle, \langle w_3 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle d_1, w_2 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_1, w_2 \rangle\ \rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle \ \ \rangle\ \rangle & \langle\langle d_0, w_3, p_0 \rangle, \langle w_0 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle d_2, w_2 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_2, w_2 \rangle\ \rangle, \\
& \langle\langle d_0, w_3, p_0 \rangle, \langle w_1 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle d_0, w_3 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_0, w_3 \rangle\ \rangle, \\
& \langle\langle d_0, w_3, p_0 \rangle, \langle w_2 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle d_1, w_3 \rangle\ \rangle, & \langle\langle d_0, w_2, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_1, w_3 \rangle\ \rangle, \\
& \langle\langle d_0, w_3, p_0 \rangle, \langle w_3 \rangle\ \rangle & \langle\langle d_0, w_3, p_0 \rangle, \langle d_2, w_1 \rangle\ \rangle, & \langle\langle d_0, w_3, p_0 \rangle, \langle \lambda\{d_2\}, \ d_2, w_1 \rangle\ \rangle, \\
& & \langle\langle d_0, w_3, p_0 \rangle, \langle d_1, w_2 \rangle\ \rangle, & \langle\langle d_0, w_3, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_1, w_2 \rangle\ \rangle, \\
& & \langle\langle d_0, w_3, p_0 \rangle, \langle d_2, w_2 \rangle\ \rangle, & \langle\langle d_0, w_3, p_0 \rangle, \langle \lambda\{d_1, d_2\}, \ d_2, w_2 \rangle\ \rangle, \\
& & \langle\langle d_0, w_3, p_0 \rangle, \langle d_0, w_3 \rangle\ \rangle, & \langle\langle d_0, w_3, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_0, w_3 \rangle\ \rangle, \\
& & \langle\langle d_0, w_3, p_0 \rangle, \langle d_1, w_3 \rangle\ \rangle & \langle\langle d_0, w_3, p_0 \rangle, \langle \lambda\{d_0, d_1\}, \ d_1, w_3 \rangle\ \rangle
\end{array}$$

$$\begin{array}{l}
c_4[[\partial[\perp\omega \in \tau\omega\{\}]]]; \partial[\perp\delta\tau = \perp\delta\{\perp\omega\}]]^g \\
\begin{array}{l}
x\{\langle\langle d_0, w_0, p_0 \rangle, \langle x\{d_2\}, d_2, w_1 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{d_1, d_2\}, d_1, w_2 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{d_0, d_1\}, d_0, w_3 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{d_2\}, d_2, w_1 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{d_1, d_2\}, d_1, w_2 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{d_0, d_1\}, d_0, w_3 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{d_2\}, d_2, w_1 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{d_1, d_2\}, d_1, w_2 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{d_0, d_1\}, d_0, w_3 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{d_2\}, d_2, w_1 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{d_1, d_2\}, d_1, w_2 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{d_0, d_1\}, d_0, w_3 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{d_0, d_1\}, d_1, w_3 \rangle\rangle\}
\end{array} \\
c_5[[p|p = \perp\omega\{\}]]^g \\
\begin{array}{l}
x\{\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_2, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_3, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle\}
\end{array} \\
c_6[[\tau\omega \notin \perp\Omega]]^g \\
\begin{array}{l}
x\{\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle\}
\end{array} \\
c_7[[\partial[\tau\omega \notin \perp\Omega]]]^g \\
\begin{array}{l}
x\{\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle\}
\end{array} \\
c_7[[\mathbf{p}|p = \tau\omega\{\}]]^g \\
\begin{array}{l}
x\{\langle\langle x\{w_0, w_1\}, d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle x\{w_0, w_1\}, d_0, w_0, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle, \\
\langle\langle x\{w_0, w_1\}, d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_1, d_2\}, d_2, w_2 \rangle\rangle, \\
\langle\langle x\{w_0, w_1\}, d_0, w_1, p_0 \rangle, \langle x\{w_2, w_3\}, x\{d_0, d_1\}, d_1, w_3 \rangle\rangle\}
\end{array}
\end{array}$$

Absurd (or equivalent) updates:

$[x| (x = ole)^\circ]^\top; (([w| w \in \tau\omega\{\}]; [y| sick_{\perp\omega}y]; [X| X = \perp\delta\{\perp\omega\}])^\perp; (\partial[\perp\omega \in \tau\omega\{\}]; [y| get_{\perp\omega}\langle\tau\delta, y\rangle, dog_{\perp\omega}y]; [p| p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [\mathbf{p}| p = \tau\omega\{\}]))$

$\equiv \mathbf{0}$ (no $\perp\delta\tau$ in the box after $^\perp$.)

$[x| (x = ole)^\circ]^\top; (([w| w \in \tau\omega\{\}]; [y| sick_{\perp\omega}y]; [X| X = \perp\delta\{\perp\omega\}])^\perp; ([v| v \in \tau\omega\{\}]; \partial[\perp\delta\tau = \perp\delta\{\perp\omega\}]; [get_{\perp\omega}\langle\tau\delta, \perp\delta\rangle, dog_{\perp\omega}\perp\delta]; [p| p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [\mathbf{p}| p = \tau\omega\{\}]))$

$\equiv \mathbf{0}$ (presupposition failure due to $\perp\omega$ reset by $[v| v \in \tau\omega\{\}]$)

$[x| (x = ole)^\circ]^\top; (([w| w \in \tau\omega\{\}]; [y| sick_{\perp\omega}y]; [X| X = \perp\delta\{\perp\omega\}])^\perp; ([v| v \in \tau\omega\{\}]; [y| get_{\perp\omega}\langle\tau\delta, y\rangle, dog_{\perp\omega}y]; [p| p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [\mathbf{p}| p = \tau\omega\{\}]))$

$\equiv \mathbf{0}$ (no $\perp\delta\tau$ in the box after $^\perp$.)

$[x| (x = ole)^\circ]^\top; (([y| sick_{\tau\omega}y]; [X| X = \perp\delta\{\tau\omega\}])^\perp; (\partial[\perp\omega \in \tau\omega\{\}]; \partial[\perp\delta\tau = \perp\delta\{\perp\omega\}]; [get_{\perp\omega}\langle\tau\delta, \perp\delta\rangle, dog_{\perp\omega}\perp\delta]; [p| p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [\mathbf{p}| p = \tau\omega\{\}]))$

$\equiv \mathbf{0}$ (presupposition failure for any realistic $[[\cdot]]$, due to $\tau\omega$ in $[X| X = \perp\delta\{\tau\omega\}]$ vs. $\perp\omega$ in $\partial[\perp\delta\tau = \perp\delta\{\perp\omega\}]$)

$[x | (x = ole)^{\circ}]^{\top}$; $(([y | sick_{\tau\omega} y]; [X | X = \perp\delta\{\{\tau\omega\}\}]^{\perp}); (\partial[\perp\omega \in \tau\omega\{\}]; [y | get_{\perp\omega}(\tau\delta, y), dog_{\perp\omega} y]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}]))$
 $\equiv \mathbf{0}$ (no $\perp\delta t$ in the box after $^{\perp}$;))

$[x | (x = ole)^{\circ}]^{\top}$; $(([y | sick_{\tau\omega} y]; [X | X = \perp\delta\{\{\tau\omega\}\}]^{\perp}); ([v | v \in \tau\omega\{\}]; \partial[\perp\delta t = \perp\delta\{\{\perp\omega\}\}]; [get_{\perp\omega}(\tau\delta, \perp\delta), dog_{\perp\omega} \perp\delta]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}]))$
 $\equiv \mathbf{0}$ (presupposition failure for any realistic $[\cdot]$, due to $\tau\omega$ in $[X | X = \perp\delta\{\{\tau\omega\}\}]^{\perp}$ vs. $\perp\omega$ in $\partial[\perp\delta t = \perp\delta\{\{\perp\omega\}\}]$)

$[x | (x = ole)^{\circ}]^{\top}$; $(([y | sick_{\tau\omega} y]; [X | X = \perp\delta\{\{\tau\omega\}\}]^{\perp}); ([v | v \in \tau\omega\{\}]; [y | get_{\perp\omega}(\tau\delta, y), dog_{\perp\omega} y]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; \partial[\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}]))$
 $\equiv \mathbf{0}$ (no $\perp\delta t$ in the box after $^{\perp}$;))

cf. ENGLISH (5_E)

(5'_E) • **Syntax 1:** negated vp

HV.TNS=N'T

buy.ASP A sick dog

vp/vp:

$\lambda P \lambda x \lambda w ([v | v \in w\{\}]; P x \perp\omega; [p | p = \perp\omega\{\}]; [w \notin \perp\Omega])$

tv:

vp: $\lambda x \lambda w [y | sick_w y, dog_w y, buy_w(x, y)]$

vp:

$\lambda x \lambda w ([v | v \in w\{\}]; [y | sick_{\perp\omega} y, dog_{\perp\omega} y, buy_{\perp\omega}(x, y)]; [p | p = \perp\omega\{\}]; [w \notin \perp\Omega])$

• **Syntax 2:** sentence

Ole

HV.TNS=N'T buy.ASP A sick dog

.

\underline{s}/vp :

$\lambda P \lambda w ([x | (x = ole)^{\circ}]; P \tau\delta w)$

vp:

$\lambda x \lambda w ([v | v \in w\{\}]; \dots; [p | p = \perp\omega\{\}]; [w \notin \perp\Omega])$

\underline{s}/s :

$\lambda W ([W \tau\omega; [p | p = \tau\omega\{\}])$

\underline{s} : $\lambda w ([x | (x = ole)^{\circ}]; [v | v \in w\{\}]; [y | sick_{\perp\omega} y, dog_{\perp\omega} y, buy_{\perp\omega}(x, y)]; [p | p = \perp\omega\{\}]; [w \notin \perp\Omega])$

s : $([x | (x = ole)^{\circ}]; [v | v \in \tau\omega\{\}]; [y | sick_{\perp\omega} y, dog_{\perp\omega} y, buy_{\perp\omega}(x, y)]; [p | p = \perp\omega\{\}]; [\tau\omega \notin \perp\Omega]; [p | p = \tau\omega\{\}])$

• **Updates**

- 1 $*p_0 [[x | (x = ole)]]^g$
 $= \lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle \rangle \rangle \mid w \in {}^0 p_0 \}$ $=: c_1$
- 2 $c_1 [[v | v \in \tau\omega\{\}]]^g$
 $= \lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle v \rangle \rangle \mid w, v \in {}^0 p_0 \}$ $=: c_2$
- 3 $c_2 [[y | sick_{\perp\omega} y, dog_{\perp\omega} y, buy_{\perp\omega}(\tau\delta, y)]]^g$
 $= \lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle d, v \rangle \rangle \mid w, v \in {}^0 p_0 \ \& \ d \in {}^0 [[sick]](v) \ \& \ d \in {}^0 [[dog]](v) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](v) \}$ $=: c_3$
- 4 $c_3 [[p | p = \perp\omega\{\}]]^g$
 $= \lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle p_1, d, v \rangle \rangle \mid w, v \in {}^0 p_0 \ \& \ d \in {}^0 [[sick]](v) \ \& \ d \in {}^0 [[dog]](v) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](v) \ \& \ p_1 = \lambda \{ u \in {}^0 p_0 \mid \exists d': d' \in {}^0 [[sick]](u) \ \& \ d' \in {}^0 [[dog]](u) \ \& \ \langle [ole], d' \rangle \in {}^0 [[buy]](u) \} \}$ $=: c_4$
- 5 $c_4 [[\tau\omega \notin \perp\Omega]]^g$
 $= \lambda \{ \langle \langle [ole], w, p_0 \rangle, \langle p_1, d, v \rangle \rangle \mid w, v \in {}^0 p_0 \ \& \ d \in {}^0 [[sick]](v) \ \& \ d \in {}^0 [[dog]](v) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](v) \ \& \ p_1 = \lambda \{ u \in {}^0 p_0 \mid \exists d': d' \in {}^0 [[sick]](u) \ \& \ d' \in {}^0 [[dog]](u) \ \& \ \langle [ole], d' \rangle \in {}^0 [[buy]](u) \} \ \& \ w \notin {}^0 p_1 \}$ $=: c_5$
- 6 $c_5 [[p | p = \tau\omega\{\}]]^g$
 $= \lambda \{ \langle \langle p_2, [ole], w, p_0 \rangle, \langle p_1, d, v \rangle \rangle \mid w, v \in {}^0 p_0 \ \& \ d \in {}^0 [[sick]](v) \ \& \ d \in {}^0 [[dog]](v) \ \& \ \langle [ole], d \rangle \in {}^0 [[buy]](v) \ \& \ p_1 = \lambda \{ u \in {}^0 p_0 \mid \exists d': d' \in {}^0 [[sick]](u) \ \& \ d' \in {}^0 [[dog]](u) \ \& \ \langle [ole], d' \rangle \in {}^0 [[buy]](u) \} \ \& \ w \notin {}^0 p_1 \ \& \ p_2 = \lambda \{ w' \in {}^0 p_0 \mid w' \notin {}^0 p_1 \} \}$ $=: c_6$

