

QUANTIFIED NP'S AND DONKEY ANAPHORA

1. SOME OBSERVATIONS

- Obs. 1: Quantified np's (*every np*, *most np*, ...) selectively quantify over the head n-referent (not unselectively over information states).

(1) *Most_x men^x who own a^y donkey beat it_y.*

e.g. $\neq_{M,g} (1)$

if $[[man]] = \{m_0, \dots, m_9\}$

& m_0 owns & beats donkey d_0, \dots, d_9

& m_1 owns & beats donkeys d_{10}, \dots, d_{19}

& m_2 owns donkey d_{20} (only) but doesn't beat d_{20}

⋮

& m_9 owns donkey d_{90} (only) but doesn't beat d_{90}

- Obs. 2: Any indefinite in the matrix clause of a quantified np has existential force.
- Obs. 3: In most donkey sentences, indefinites in the restriction of a quantified np have universal force ('strong reading').

(2) *Most_x men^x who own a^y donkey beat it_y, with a^z stick.*

= for most people x , for **every** donkey y owned by x ,
there is **a** stick z s.t. person x beats y with z

(3) *Most_x people^x that owned a^y slave also owned his_y offspring.*

(Heim 1990)

= for most **people** x , for **every** slave y owned by x & **every** child z of y ,
person x also owned z .

- Obs. 4: In some donkey sentences, indefinites in the restriction of a quantified np have existential force ('weak reading').

(4) *Most_x people^x who had a^y credit card paid their bill with it.*

(R. Cooper)

= for most people x s.t. x had **a** credit card y ,
person x paid x 's bill with credit card y .

(5) *(In those days) most_x people^x that owned a^y slave were considered respectable.*

= for most people x s.t. x owned **a** slave y ,
person x was considered respectable.

- Obs. 5: When the restriction of a quantified np contains multiple indefinites it may be possible to read some of them as strong and some as weak ('mixed reading').

(6) *Most_x people^x who have a^y credit card and buy a^z computer use it_y to pay for it_z.*

(Brasoveanu 2007)

= for most people x s.t. x has **a** credit card y , for **every** computer z s.t. x buys y ,
 x uses one of x 's credit cards to pay for z .

2. REPRESENTATION IN DPL^Q (df. in APPENDIX)

Notation:

$\text{all}(X, Y)$ iff $X \subseteq Y$

$\text{most}(X, Y)$ iff $|X \cap Y| > |X \setminus Y|$

- Obs. 1: Quantified np's (*every np, most np, ...*) selectively quantify over the head n-referent (not unselectively over information states).
- Obs. 2: Any indefinites in the matrix clause of a quantified np have existential force.
- Obs. 3: Most indefinites in the restriction of a quantified np have universal force ('strong reading').

(2) $\text{Most}_x \text{men}^x$ who own a^y donkey beat it_y with a^z stick.

(2') $g[\text{MOST}_x^v([x \ y \ \text{man } x, \ \text{dnk } y, \ \text{own } xy], [z \ \text{stk } z, \ \text{btw } xyz])]h$
iff $g = h$

& $\text{most}(\{k(x): h[[x \ y \ \text{man } x, \ \text{dnk } y, \ \text{own } xy]]k\},$
 $\{k(x): h[[x \ y \ \text{man } x, \ \text{dnk } y, \ \text{own } xy]]k$
 & $\forall j(h[[x \ y \ \text{man } x, \ \text{dnk } y, \ \text{own } xy]]j \ \& \ j(x) = k(x) \Rightarrow \exists i: j[[z \ \text{stk } z, \ \text{btw } xyz]]i\})$

iff $g = h$

& $\text{most}(\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])\},$
 $\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])$
 & $\forall b \in D(a \in [\text{man}] \ \& \ b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}]$
 $\Rightarrow \exists c \in D(c \in [\text{stk}] \ \& \ \langle a, b, c \rangle \in [\text{btw}]))$

iff $g = h$

& $\text{most}(\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])\},$
 $\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])$
 & $\forall b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}] \Rightarrow \exists c \in D(c \in [\text{stk}] \ \& \ \langle a, b, c \rangle \in [\text{btw}]))$

(2'') $\models_{M, g} \text{MOST}_x^v([x \ y \ \text{man } x, \ \text{dnk } y, \ \text{own } xy], [z \ \text{stk } z, \ \text{btw } xyz])$

iff $\text{most}(\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])\},$
 $\{a \in D \mid a \in [\text{man}] \ \& \ \exists b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])$
 & $\forall b \in D(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}] \Rightarrow \exists c \in D(c \in [\text{stk}] \ \& \ \langle a, b, c \rangle \in [\text{btw}]))$

e.g. If $[\text{man}] = \{m_0, \dots, m_9\}$

& m_0 owns & beats donkey d_0, \dots, d_9 with a stick

& m_1 owns & beats donkeys d_{10}, \dots, d_{19} with a stick

& m_2 owns donkey d_{20} (only) but doesn't beat d_{20} with any stick

⋮

& m_9 owns donkey d_{90} (only) but doesn't beat d_{90} with any stick

then $\{a \mid a \in [\text{man}] \ \& \ \exists b(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])\}$

$= \{m_0, \dots, m_9\}$

$\{a \mid a \in [\text{man}] \ \& \ \exists b(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}])\}$

& $\forall b(b \in [\text{dnk}] \ \& \ \langle a, b \rangle \in [\text{own}] \Rightarrow \exists c(c \in [\text{stk}] \ \& \ \langle a, b, c \rangle \in [\text{btw}]))$

$= \{m_0, m_1\}$

Hence $\text{most}(\{m_0, \dots, m_9\}, \{m_0, m_1\})$ is **false**,

since $|\{m_0, m_1\}| = 2 > |\{m_2, \dots, m_9\}| = 8$ if **false**

- (3) *Most_x people^x that owned a^y slave also owned his_y offspring.* (Heim 1990)
 = for most **people** x , for **every** slave y owned by x & **every** child z of y ,
 person x owned z .

- (3') $g[\text{MOST}^{\forall}_x([x\ y\ z\ | \text{prs } x, \text{slv } y, \text{own } xy, \text{chl } zy], [\text{own } xz])]h$ (local presup.
 iff $g = h$ accommodation)
 & $\text{most}(\{k(x): h[[x\ y\ z\ | \text{prs } x, \text{slv } y, \text{own } xy, \text{chl } zy]]k\},$
 $\{k(x): h[[x\ y\ z\ | \text{prs } x, \text{slv } y, \text{own } xy, \text{chl } zy]]k$
 & $\forall j(h[[x\ y\ z\ | \text{prs } x, \text{slv } y, \text{own } xy, \text{chl } zy]]j \ \& \ j(x) = k(x) \Rightarrow \exists i: j[[\text{own } xz]]i\})$
 iff $g = h$
 & $\text{most}(\{a \in D \mid a \in [\text{prs}] \ \& \ \exists b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}])\},$
 $\{a \in D \mid a \in [\text{prs}] \ \& \ \exists b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}])$
 & $\forall b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}] \Rightarrow \langle a, c \rangle \in [\text{own}])\})$

- (3'') $\models_{M, g} \text{MOST}^{\forall}_x([x\ y\ z\ | \text{prs } x, \text{slv } y, \text{own } xy, \text{chl } zy], [\text{own } xz])$
 iff $\text{most}(\{a \in D \mid a \in [\text{prs}] \ \& \ \exists b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}])\},$
 $\{a \in D \mid a \in [\text{prs}] \ \& \ \exists b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}])$
 & $\forall b, c \in D(b \in [\text{slv}] \ \& \ \langle a, b \rangle \in [\text{own}] \ \& \ \langle c, b \rangle \in [\text{chl}] \Rightarrow \langle a, c \rangle \in [\text{own}])\})$

- Obs. 4: Some indefinites in the restriction of a quantified np have existential force ('weak reading').

- (4) *Most_x people^x who had a^y credit card paid their bill with it.* (R. Cooper)
 = for most people x s.t. x had **a** credit card y ,
 person x paid x 's bill with credit card y .

- (4') $g[\text{MOST}^{\exists}_x(\underline{\hspace{10em}}, \underline{\hspace{10em}})]h$
 iff $g = h$
 & $\underline{\hspace{10em}}$

- (5) *(In those days) most_x people^x that owned a^y slave were considered respectable.*
 = for most people x s.t. x owned **a** slave y ,
 person x was considered respectable.

- (5') $\underline{\hspace{10em}}$
 $\underline{\hspace{10em}}$

3. NO REPRESENTATION IN DPL^Q

- Obs. 5: When the restriction of a quantified np contains multiple indefinites it may be possible to read some of them as strong and some as weak ('mixed reading').

- (6) *Most_x people^x who have a^y credit card and buy a^z computer* (Brasoveanu 2007)
use it_y to pay for it_z.
 = for most people x s.t. x has **a** credit card y , for **every** computer z s.t. x buys y ,
 x uses one of x 's credit cards to pay for z .

APPENDIX: DPL⁺ WITH QUANTIFIERS (DPL^Q, K&R 93, Brasoveanu 07)D1.1 (DPL^Q basic expressions)

- **Var** = $\{x, y, z, \dots, z_1, z_2, \dots\}$
- **Con** = $\{jim, ann, \dots\}$
- **Prd¹** = $\{SG, PL\} \cup \{man, dnk, \dots, cm, sit, \dots\}$
- **Prd²** = $\{hsb, wife, see, spk, \dots\}$
- **Prd³** = $\{gvt, \dots\}$

D1.2 (DPL^Q syntax)

- | | | | | |
|-------|---------------------------------------|----|---|---------------------|
| 0. | $t \in \mathbf{Trm}$ | if | $t \in \mathbf{Con} \cup \mathbf{Var}$ | |
| i. | $(t_1 + t_2) \in \mathbf{Trm}$ | if | $t, t_1, \dots, t_n \in \mathbf{Trm}$ | summation |
| ii. | $[u] \in \mathbf{For}$ | if | $u \in \mathbf{Var}$ | |
| iii. | $(\phi, \psi) \in \mathbf{For}$ | if | $\phi, \psi \in \mathbf{For}$ | |
| iv. | $\neg\phi \in \mathbf{For}$ | if | $\phi \in \mathbf{For}$ | |
| v. | $Rt_1 \dots t_n \in \mathbf{For}$ | if | $R \in \mathbf{Prd}^n$ & $t_1, \dots, t_n \in \mathbf{Trm}$ | |
| vi. | $t_1 \leq t_2 \in \mathbf{For}$ | if | $t_1, t_2 \in \mathbf{Trm}$ | part-of |
| vii. | $ t = n \in \mathbf{For}$ | if | $t \in \mathbf{Trm}$ and $n \in \{1, 2, \dots\}$ | cardinality |
| viii. | $t = \Sigma u. \phi \in \mathbf{For}$ | if | $t \in \mathbf{Trm}, u \in \mathbf{Var}$ and $\phi \in \mathbf{For}$ | abstraction |
| ix. | $Q^v_u(\phi, \psi) \in \mathbf{For}$ | if | $\phi, \psi \in \mathbf{For}, u \in \mathbf{Var}$ and $Q \in \{\text{ALL, MOST}, \dots\}$ | selective str. qnt. |
| | $Q^3_u(\phi, \psi) \in \mathbf{For}$ | if | $\phi, \psi \in \mathbf{For}, u \in \mathbf{Var}$ and $Q \in \{\text{ALL, MOST}, \dots\}$ | selective wk. qnt. |

D2.1 (DPL^Q models and assignments)

- A DPL^Q *model* is a pair $M = \langle D, \llbracket \cdot \rrbracket \rangle$ such that (i) $D = \mathcal{P}(A) \setminus \{\emptyset\}$ for some set $A \neq \emptyset$
- (ii) $\llbracket \cdot \rrbracket$ assigns $\llbracket A \rrbracket \in \{d \in D : |d| = 1\}$ to each $A \in \mathbf{Con}$, and $\llbracket B \rrbracket \subseteq (D)^n$ to each $B \in \mathbf{Prd}^n$.
- (iii) $\llbracket SG \rrbracket = \{d \in D : |d| = 1\}$
- $\llbracket PL \rrbracket = \{d \in D : |d| > 1\}$
- $G := \{g \mid g: \mathbf{Var} \rightarrow D\}$ is the set of all *M*-assignments
- for all $g, h \in G$: $g[u_1 \dots u_n]h$ iff $\forall v \in \mathbf{Var} \setminus \{u_1, \dots, u_n\} : g(v) = h(v)$

D2.2 (DPL^Q semantics). For any DPL-model $M = \langle D, \llbracket \cdot \rrbracket \rangle$ and *M*-assignments g and h :

- | | | | | |
|-------|--|-----|--|-------------------------|
| 0. | $\llbracket t \rrbracket^h$ | = | $\llbracket t \rrbracket$ | if $t \in \mathbf{Con}$ |
| | | = | $h(t)$ | if $t \in \mathbf{Var}$ |
| i. | $\llbracket t_1 + t_2 \rrbracket^h$ | = | $(\llbracket t_1 \rrbracket^h \cup \llbracket t_2 \rrbracket^h)$ | |
| ii. | $g[\llbracket u \rrbracket]h$ | iff | $g[u]h$ | |
| iii. | $g[\llbracket \phi; \psi \rrbracket]h$ | iff | $\exists k(g[\llbracket \phi \rrbracket]k \ \& \ k[\llbracket \psi \rrbracket]h)$ | |
| iv. | $g[\llbracket \neg\phi \rrbracket]h$ | iff | $g = h \ \& \ \neg \exists k: h[\llbracket \phi \rrbracket]k$ | |
| v. | $g[\llbracket Rt_1 \dots t_n \rrbracket]h$ | iff | $g = h \ \& \ \langle \llbracket t_1 \rrbracket^h, \dots, \llbracket t_n \rrbracket^h \rangle \in \llbracket R \rrbracket$ | |
| vi. | $g[\llbracket t_1 \leq t_2 \rrbracket]h$ | iff | $g = h \ \& \ \llbracket t_1 \rrbracket^h \subseteq \llbracket t_2 \rrbracket^h$ | |
| vii. | $g[\llbracket t = n \rrbracket]h$ | iff | $g = h \ \& \ \llbracket t \rrbracket^h = n$ | |
| viii. | $g[\llbracket t = \Sigma u. \phi \rrbracket]h$ | iff | $g = h \ \& \ \llbracket t \rrbracket^h = \cup \{k(u) : h[\llbracket \phi \rrbracket]k\}$ | |
| ix. | $g[\llbracket Q^3_u(\phi, \psi) \rrbracket]h$ | iff | $g = h \ \& \ Q(\{k(u) : h[\llbracket \phi \rrbracket]k\}, \{k(u) : \exists j(h[\llbracket \phi \rrbracket]j \ \& \ j[\llbracket \psi \rrbracket]k)\})$ | |
| | $g[\llbracket Q^v_u(\phi, \psi) \rrbracket]h$ | iff | $g = h \ \& \ Q(\{k(u) : h[\llbracket \phi \rrbracket]k\}, \{k(u) : h[\llbracket \phi \rrbracket]k \ \& \ \forall j(h[\llbracket \phi \rrbracket]j \ \& \ j(u) = k(u) \Rightarrow \exists i: j[\llbracket \psi \rrbracket]i)\})$ | |

D3 (Truth, entailment, and equivalence).

- i. ϕ is true in M given input g , written $\models_{M,g} \phi$, iff $\exists h: g \llbracket \phi \rrbracket h$
- ii. ϕ (DPL^Q-)entails ψ , written $\phi \models \psi$, iff $\forall M \forall g \forall h (g \llbracket \phi \rrbracket h \Rightarrow \exists k: h \llbracket \psi \rrbracket k)$
- iii. ϕ is (DPL^Q-)equivalent to ψ , written $\phi \simeq \psi$, iff $\forall M \forall g \forall h (g \llbracket \phi \rrbracket h \text{ iff } g \llbracket \psi \rrbracket h)$

ABBREVIATIONS

- $\exists u \phi$:= $([u]; \phi)$
- $\forall u \phi$:= $\neg \exists u \neg \phi$
- $(\phi \rightarrow \psi)$:= $\neg(\phi; \neg \psi)$
- $[\phi]$ = ϕ if ϕ is a test (i.e. $\forall M \forall g \forall h: g \llbracket \phi \rrbracket h \Rightarrow g = h$)
- $[\phi]$ = undefined otherwise
- (ϕ, ψ) := $([\phi]; [\psi])$
- $[u_1 \dots u_n | \phi]$:= $([u_1]; \dots ([u_n]; [\phi]))$
- $(t_1 = t_2)$:= $(t_1 \leq t_2, t_2 \leq t_1)$
- $(t_1 \in t_2)$:= (SG $t_1, t_1 \leq t_2$)
- $(t_1 \emptyset t_2)$:= $\neg([x]; [x \leq t_1, x \leq t_2])$

USEFUL FACTS

- $F\exists$: $g \llbracket \exists u \phi \rrbracket h$
iff $\exists k (g[u]k \ \& \ k \llbracket \phi \rrbracket h)$
- $F\forall$: $g \llbracket \forall u \phi \rrbracket h$
iff $g = h \ \& \ \forall k (g[u]k \Rightarrow \exists i: k \llbracket \phi \rrbracket i)$
- $F\rightarrow$: $g \llbracket \phi \rightarrow \psi \rrbracket h$
iff $g = h \ \& \ \forall k (g \llbracket \phi \rrbracket k \Rightarrow \exists i: k \llbracket \psi \rrbracket i)$
- $F,$: $g \llbracket (\phi, \psi) \rrbracket h$
iff $g = h \ \& \ h \llbracket \phi \rrbracket h \ \& \ h \llbracket \psi \rrbracket h$
- $F|$: $g \llbracket [u | \phi_1, \dots, \phi_n] \rrbracket h$
iff $g[u]h \ \& \ h \llbracket \phi_1 \rrbracket h \ \& \ \dots \ h \llbracket \phi_n \rrbracket h$
- $g \llbracket [u_1 u_2 | \phi_1, \dots, \phi_n] \rrbracket h$
iff $g[u_1 u_2]h \ \& \ h \llbracket \phi_1 \rrbracket h \ \& \ \dots \ h \llbracket \phi_n \rrbracket h$
- $g \llbracket [u_1 u_2 u_3 | \phi_1, \dots, \phi_n] \rrbracket h$
iff $g[u_1 u_2 u_3]h \ \& \ h \llbracket \phi_1 \rrbracket h \ \& \ \dots \ h \llbracket \phi_n \rrbracket h$
- $F=$: $g \llbracket (t_1 = t_2) \rrbracket h$
iff $g = h \ \& \ \llbracket t_1 \rrbracket^h = \llbracket t_2 \rrbracket^h$
- $F\in$: $g \llbracket (t_1 \in t_2) \rrbracket h$
iff $g = h \ \& \ |\llbracket t_1 \rrbracket^h| = 1 \ \& \ \llbracket t_1 \rrbracket^h \subseteq \llbracket t_2 \rrbracket^h$
- $F\emptyset$: $g \llbracket (t_1 \emptyset t_2) \rrbracket h$
iff $g = h \ \& \ \llbracket t_1 \rrbracket^h \cap \llbracket t_2 \rrbracket^h = \emptyset$

DISCOURSE ANAPHORA TO QUANTIFIED NP'S

1. OBSERVATIONS

- Obs. 1 (E-type anaphora) In a configuration [...a/qnt [...][...]]^u (and) ... **ana_u**, the anaphor *ana_u* is usu. anaphoric to the entire set of entities that satisfy the antecedent description.
 - (1) [*There is a doctor in London*]^x. ?**He_x** is Welsh. (Evans 1977)
he_x = the (unique) doctor in London
 \neq *There is a doctor in London who is Welsh.*
 - (2) [*Few/*no MPs came to the party*]^x but **they_x** had a marvelous time. (Evans 1977)
[*One/*no woman_x drank champagne*]^y. **She_y** got very drunk.
they_x = the (few) MPs who came to the party
 \neq *There are few/no MPs that came to the party and had a marvelous time.*
she_y = the (one) female MP who came to the party & drank champagne.
 \neq *There was one female MP who (came to the party,) drank champagne and got very drunk.*
 - (3) [*Yesterday Ann^x put exactly ten balls in the bag*]^y. [*Now two of **them_y** are missing*]^z.
They_z are under the sofa. (Partee, cf. (7))
them_y = the (ten) balls that Ann put in the bag yesterday
they_z = the (two) balls that Ann put in the bag yesterday & that are now missing
- Obs. 2: (domain set ana.). Most quantifiers license anaphora to the domain of quantification.
 - (4) [*Many/at least ten/most [MPs]^x came to the party*]^y. **A few_x** preferred a quiet evening at home.
a few_x = a few of the aforementioned MPs
 - (5) [*Few [MPs]^x came to the party*]^y. **Most_x** preferred a quiet evening at home.
most_x = most of the aforementioned MPs
 - (6) (*In this Republican district*) [*few [parents]^x agreed to let their children^y listen to Obama's speech*].
Most_x were afraid he would indoctrinate them_y.
most_x = most of the aforementioned parents
- Obs. 3 (complement set ana.). Most quantifiers do not support anaphora to the complement set (see #(7³), #(8²)), but some, e.g. *few*, do (see (8), (9)).
 - (7) [*Ann put exactly ten balls in the bag*]^y. [*Only eight of **them_y** are in the bag now*]^z.
#*They_{y|z} are under the sofa.* (Partee cf. (3))
they_{y|z} = the (two) balls that Ann put in the bag yesterday & that are not in the bag now
 - (8) [*Many/at least ten/most [MPs]^x came to the party*]^y. #**They_{x|y}** preferred a quiet evening at home.
they_{x|y} = the (minority of) MPs who didn't come to the party
 - (9) [*Few [MPs]^x came to the party*]^y. **They_{x|y}** preferred a quiet evening at home.
they_{x|y} = the (majority of) MPs who didn't come to the party
 - (10) (*In this Republican district*) [*few [parents]^x agreed to let their children^y listen to Obama's speech*]^z.
They_{x|z} were afraid he would indoctrinate them_y.
they_{x|z} = most of the aforementioned parents

2. E-TYPE ANAPHORA IN DPL^Q

- Obs. 1 (E-type anaphora) In a configuration [...*a/qnt* [...][...]]^u (*and*) ... ***ana_u***, the anaphor *ana_u* is usu. anaphoric to the entire set of entities that satisfy the antecedent description.

cf. *There is a doctor in London who is Welsh.*

g[_____
iff _____

(1¹) [*There is a doctor in London*]^x.

g[_____
iff _____

(1²) ?***He_x*** *is Welsh.*

g[_____
iff _____

(2¹) [*Few MPs came to the party*]^x ...

g[_____
iff _____

... ***they_x*** *had a marvelous time.*

g[_____
iff _____

(2²) [*One woman_x drank champagne*]^y.

g[_____
iff _____

(2³) ***She_y*** *got very drunk.*

g[_____
iff _____

cf. # [*No MPs came to the party*]^x *but* ***they_x*** *had a marvelous time.*

g[_____
iff _____

3. DOMAIN SET ANAPHORA IN DPL^Q

- Obs. 2: (domain set ana.). Most quantifiers license anaphora to the domain of quantification.

(4¹) [*Most* [MPs]^x came to the party]^y.

g[_____
 iff _____

(4²) **A few**_x preferred a quiet evening at home.

g[_____
 iff _____

(5¹) [*Few* [MPs]^x came to the party]^y.

g[_____
 iff _____

(5²) **Most**_x preferred a quiet evening at home.

g[_____
 iff _____

(6¹) *Few* [parents]^x agreed to let their children^y listen to Obama's speech.

g[_____
 iff _____

(6²) **Most**_x were afraid he would indoctrinate them_y.

g[_____
 iff _____

4. COMPLEMENT SET ANAPHORA IN DPL^Q

- Obs. 3 (complement set ana.). Most quantifiers do not support anaphora to the complement set (see #(7³), #(8²)), but some, e.g. *few*, do (see (8), (9)).

(8¹) [*Most* [MPs]^x came to the party]^y.

g[_____

iff _____

(9¹) [*Few* [MPs]^x came to the party]^y.

g[_____

iff _____

(...²) **They**_{x,y} preferred a quiet evening at home.

g[_____

iff _____

(10¹) [*Few* [parents]^x agreed to let their children^z listen to Obama's speech]^y.

g[_____

iff _____

(10²) **They**_{x,y} were afraid he would indoctrinate them_z.

g[_____

iff _____

APPENDIX 1: E-TYPE ANAPHORA IN DPL^Q

- Obs. 1 (E-type anaphora) In a configuration [...*a*/qnt [...][...]]^u (*and*) ... **ana_u**, the anaphor *ana_u* is usu. anaphoric to the entire set of entities that satisfy the antecedent description.
- cf. *There is a doctor in London who is Welsh.*
 $g[[x| doc\ x, in.london\ x, welsh\ x]]k$
iff $g[x]k \ \& \ k(x) \in [[doc]] \ \& \ k(x) \in [[in.london]] \ \& \ k(x) \in [[welsh]]$
- $|=_{M,g} [x| doc\ x, in.london\ x, welsh\ x]]h$
iff $\exists a(a \in [[doc]] \ \& \ a \in [[in.london]] \ \& \ a \in [[welsh]])$
- (1) [*There is a doctor in London*]^x. **He_x** *is Welsh.*
 $g[[x| doc\ x, in.london\ x]; [y| SG\ y, y = \Sigma x.[x| doc\ x, in.london\ x], welsh\ y]]k$
iff $\exists h(g[x]h \ \& \ h(x) \in [[doc]] \ \& \ h(x) \in [[in.london]]$
 $h[y]k \ \& \ |k(y)| = 1 \ \& \ k(y) = \cup \{a| a \in [[doc]], a \in [[in.london]]\} \ \& \ k(y) \in [[welsh]])$
iff $g[x\ y]k \ \& \ \{k(x)\} = \{a| a \in [[doc]] \ \& \ a \in [[in.london]]\}$
 $\ \& \ |k(y)| = 1 \ \& \ k(y) = k(x) \ \& \ k(y) \in [[welsh]]$
- $|=_{M,g} [x| doc\ x, in.london\ x]; [y| SG\ y, y = \Sigma x.[x| doc\ x, in.london\ x], welsh\ y]]h$
iff $\exists a(|a| = 1 \ \& \ \{a\} = \{b| b \in [[doc]] \ \& \ b \in [[in.london]]\} \ \& \ a \in [[welsh]])$
- (2¹) [*Few MPs came to the party*]^x ...
 $g[[FEW^3_x([x| mp\ x], [cm.prt\ x]]]h$
iff $g = h \ \& \ FEW(\{a| a \in [[mp]]\}, \{a| a \in [[mp]] \ \& \ a \in [[cm.prt]]\})$
- (2²) ... **they_x** *had a marvelous time.*
 $g[[x| PL\ x, x = \Sigma x_1.[x_1| mp\ x_1, cm.prt\ x_1], hv.mrv.tm\ x]]h$
iff $g[x]h \ \& \ |h(x)| > 1 \ \& \ h(x) = \cup \{b| b \in [[mp]] \ \& \ b \in [[cm.prt]]\} \ \& \ h(x) \in [[hv.mrv.tm]]$
- (2³) [*One woman_x drank champagne*]^y.
 $h[[ONE^3_y([y| wmn\ y, y \in x], [drnk.ch\ y]]]k$
iff $h = k \ \& \ |\{a| a \in [[wmn]] \ \& \ |a| = 1 \ \& \ a \subseteq k(x) \ \& \ a \in [[drnk.ch]]\}| = 1$
- (2⁴) **She_y** *got very drunk.*
 $h[[z| SG\ z, z = \Sigma y.[y| wmn\ y, y \in x, drnk.ch\ y], gt.drnk\ z]]k$
iff $h[z]k \ \& \ \{k(z)\} = \{a| a \in [[wmn]] \ \& \ |a| = 1 \ \& \ a \subseteq k(x) \ \& \ a \in [[drnk.ch]]\} \ \& \ k(z) \in [[gt.drnk]]$
- cf. # [*No MPs came to the party*]^x *but they_x had a marvelous time.*
 $g[[NO^3_x([x| mp\ x], [cm.prt\ x]); [y| PL\ y, y = \Sigma x.[x| mp\ x, cm.prt\ x], hv.mrv.tm\ y]]h$
iff $g[y]h \ \& \ \{a| a \in [[mp]] \ \& \ a \in [[cm.prt]]\} = \emptyset$
 $\ \& \ |h(y)| > 1 \ \& \ h(y) = \cup \{a| a \in [[mp]] \ \& \ a \in [[cm.prt]]\} \ \& \ h(y) \in [[hv.mrv.tm]]$
iff $g[y]h \ \& \ \{a| a \in [[mp]] \ \& \ a \in [[cm.prt]]\} = \emptyset$
 $\ \& \ |h(y)| > 1 \ \& \ h(y) = \emptyset \ \& \ h(y) \in [[hv.mrv.tm]]$
(i.e. no output *h* for any input *g*)

APPENDIX 2: DOMAIN SET ANAPHORA IN DPL^Q

- Obs. 2: (domain set ana.). Most quantifiers license anaphora to the domain of quantification.

(4¹) *Most [MPs]^x came to the party.*

$g[[\text{MOST}_x^3([x|mp\ x], [cm.prt\ x])]]h$

iff $g = h \ \& \ \text{MOST}(\{a \mid a \in [[mp]]\}, \{a \mid a \in [[mp]] \ \& \ a \in [[cm.prt]]\})$

(4²) *A few_x preferred a quiet evening at home.*

$g[[[x\ y \mid x = \Sigma x_1.[x_1 \mid mp\ x_1], y \leq x,$

$\text{FEW}_z^3([z \mid z \in x], [z \in y]),$

$\text{ALL}_z^3([z \mid z \in y], [prfr.qu\ z])]]h$

iff $g[x\ y]h$

$\& \ h(x) = \cup \{a \mid a \in [[mp]]\} \ \& \ h(y) \subseteq h(x)$

$\& \ \text{FEW}(\{b \mid |b| = 1 \ \& \ b \subseteq h(x)\}, \{b \mid |b| = 1 \ \& \ b \subseteq h(x) \ \& \ b \subseteq h(y)\})$

$\& \ \text{ALL}(\{b \mid |b| = 1 \ \& \ b \subseteq h(y)\}, \{b \mid |b| = 1 \ \& \ b \subseteq h(y) \ \& \ b \in [[prfr.qu]]\})$

(5¹) *Few [MPs]^x came to the party.*

$g[[\text{FEW}_x^3([x|mp\ x], [cm.prt\ x])]]h$

iff $g = h$

$\& \ \text{FEW}(\{a \mid a \in [[mp]]\}, \{a \mid a \in [[mp]] \ \& \ a \in [[cm.prt]]\})$

(5²) *Most_x preferred a quiet evening at home.*

$g[[\text{MOST}_x^3([x|mp\ x], [prfr.qu\ x])]]h$

iff $g = h \ \& \ \text{MOST}(\{b \mid b \in [[mp]]\}, \{b \mid b \in [[mp]] \ \& \ b \in [[prfr.qu]]\})$

(6¹) *Few [parents]^x agreed to let their children^y listen to Obama's speech.*

$g[[[z \mid z = obm, \text{FEW}_x^3([x\ y \mid y = \Sigma y_1.[y_1 \mid prnt\ xy_1, chld\ y_1]], [agr.lstn\ xyz])]]h$

iff $g[z]h \ \& \ h(z) = [[obm]]$

$\& \ \text{FEW}(\{a \mid \exists b \in D(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \ \& \ c = [[chld]]\})\},$

$\{a \mid \exists b \in D(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \ \& \ c \in [[chld]]\} \ \& \ \langle a, b, h(z) \rangle \in [[agr.lstn]]\})$

(6²) *Most_x were afraid he would indoctrinate them_y.*

$h[[\text{MOST}_x^3([x\ y \mid y = \Sigma y_1.[y_1 \mid prnt\ xy_1, chld\ y_1]], [afrd.ind\ xzy])]]k$

iff $h = k$

$\& \ \text{MOST}(\{a \mid \exists b \in D(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \ \& \ c = [[chld]]\})\},$

$\{a \mid \exists b \in D(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \ \& \ c = [[chld]]\} \ \& \ \langle a, h(z), b \rangle \in [[afrd.ind]]\})$

APPENDIX 3: COMPLEMENT SET ANAPHORA IN DPL^Q

- Obs. 3 (complement set ana.). Most quantifiers do not support anaphora to the complement set (see #(7³), #(8²)), but some, e.g. *few*, do (see (8), (9)).

(8¹) [*Most [MPs]^x came to the party*]^y.

$g[[\text{MOST}^3_x([x|mp\ x], [cm.prt\ x])]]h$

iff $g = h$

$\& \text{MOST}(\{a \mid a \in [[mp]]\}, \{a \mid a \in [[mp]] \& a \in [[cm.prt]]\})$

(9¹) [*Few [MPs]^x came to the party*]^y.

$g[[\text{FEW}^3_x([x|mp\ x], [cm.prt\ x])]]h$

iff $g = h$

$\& \text{FEW}(\{a \mid a \in [[mp]]\}, \{a \mid a \in [[mp]] \& a \in [[cm.prt]]\})$

(...²) *They_{x,y} preferred a quiet evening at home.*

$g[[[x] \text{ PL } x, x = \Sigma x_1.[x_1|mp\ x_1, \neg[cm.prt\ x_1]], \text{ALL}^3_z([z|z \in x], [prfr.qu\ z])]]h$

iff $g[x]h$

$\& |h(x)| > 1 \& h(x) = \cup \{a \mid a \in [[mp]] \& a \notin [[cm.prt]]\}$

$\& \text{ALL}(\{b \mid |b| = 1 \& b \subseteq h(x)\}, \{b \mid |b| = 1 \& b \subseteq h(x) \& b \in [[prfr.qu]]\})$

Problem:

No asymmetry between *few* and other quantifiers. Hence no explanation why discourse (9¹); (...²) is more acceptable e.g. than the clearly infelicitous (8¹); (...²)

(10¹) [*Few [parents]^x agreed to let their children^z listen to Obama's speech*]^y.

$g[[[z] z = obm, \text{FEW}^3_x([x\ y|y = \Sigma y_1.[y_1|prnt\ xy_1, chld\ y_1]], [agr.lstn\ xyz])]]h$

iff $g[z]h \& h(z) = [[obm]]$

$\& \text{FEW}(\{a \mid \exists b \in D(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \& c = [[chld]]\})\},$

$\{a \mid \exists b \in D(b = \cup \{b' \mid \langle a, c \rangle \in [[prnt]] \& c \in [[chld]] \& \langle a, b, h(z) \rangle \in [[agr.lstn]]\})\})$

(10²) *They_{x,y} were afraid he would indoctrinate them_z.*

$h[[[x] \text{ PL } x, x = \Sigma x_1.[x_1\ y|y = \Sigma y_1.[y_1|prnt\ xy_1, chld\ y_1], \neg[agr.lstn\ xyz]],$

$\text{ALL}_{x_1}^3([x_1\ y|y = \Sigma y_1.[y_1|prnt\ xy_1, chld\ y_1], \neg[agr.lstn\ xyz]], [afrd.ind\ xzy])]]k$

iff $h[x]k$

$\& k(x) = \cup \{a \mid \exists b(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \& c = [[chld]]\}) \& \langle a, b, h(z) \rangle \notin [[agr.lstn]]\}$

$\& \text{ALL}(\{a \mid \exists b(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \& c = [[chld]]\}) \& \langle a, b, h(z) \rangle \notin [[agr.lstn]]\},$

$\{a \mid \exists b(b = \cup \{c \mid \langle a, c \rangle \in [[prnt]] \& c = [[chld]]\}) \& \langle a, b, h(z) \rangle \notin [[agr.lstn]]$

$\& \langle a, h(z), b \rangle \in [[afrd.ind]]\})$